

Lecture 5

Network Science

Strong and Weak

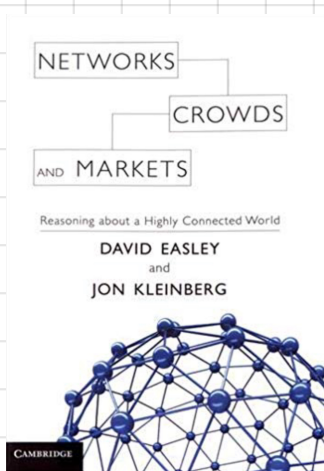
ties:

Partitions and

Betweenness

Today's Topics

- Closure, Structural Holes, and Social Capital
- Betweenness Measures and Graph Partitioning



Chapter 3

"Strong and Weak Ties"

Heterogeneity

tightly connected groups

vs

Weak ties (local bridges)



• different roles

• measures

focus on edges that

span

across different regions

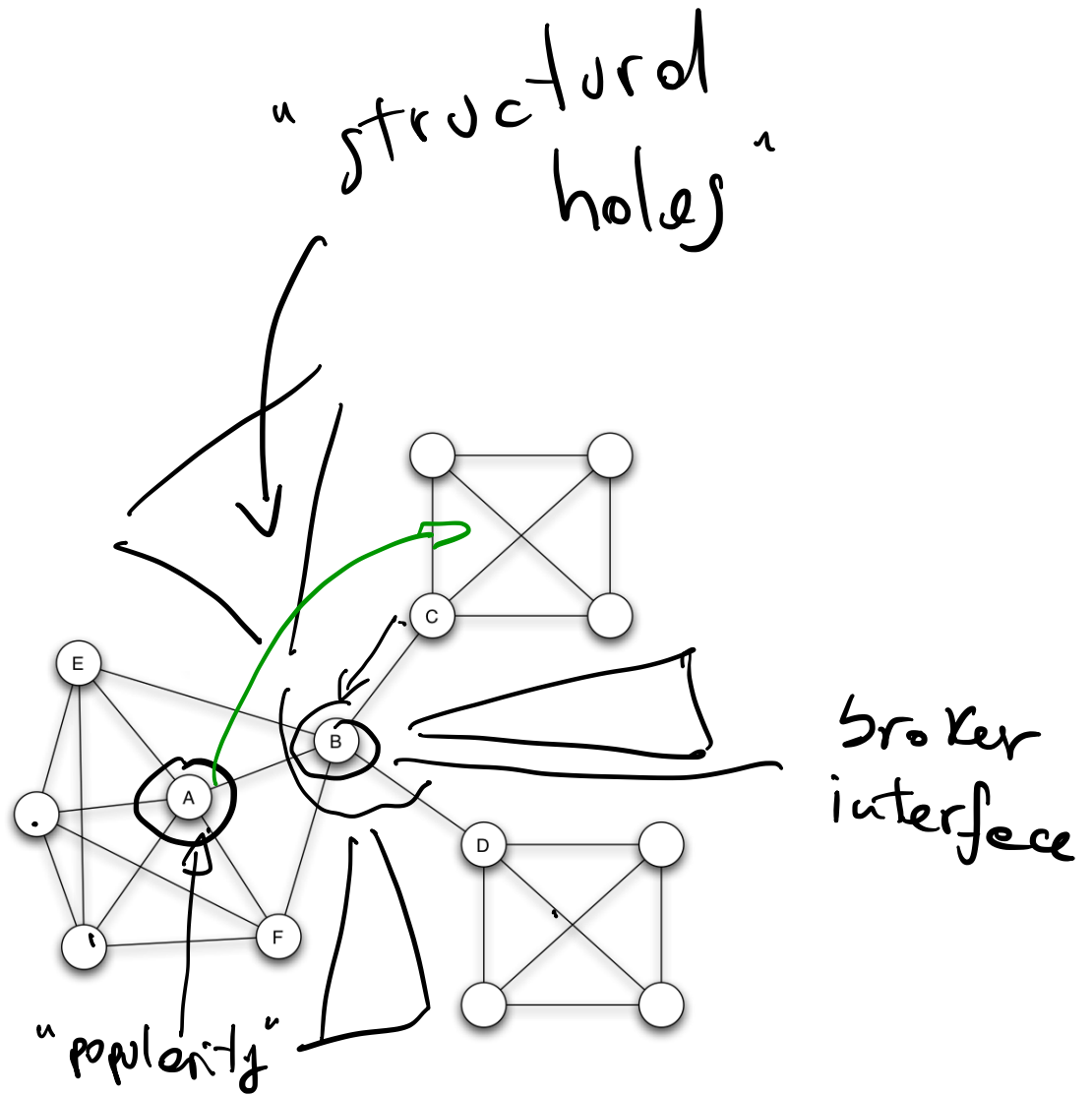


Figure 3.11: The contrast between densely-knit groups and boundary-spanning links is reflected in the different positions of nodes *A* and *B* in the underlying social network.

Ronald Burt, 2000
 "organizations"

Embeddedness

$$NO_{AB} = \frac{\# \text{ common neigh. of } A \& B}{\# \text{ neigh. of } A \& B}$$

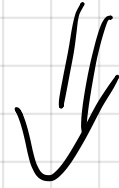
Embeddedness = numerator of
NO

"similar" the clustering coefficient

these measures that correlate
"a lot"

When the interaction is
"on display"

embeddedness \Leftrightarrow misbehavior



Social consequences

Weak ties have a kind
of privilege:

not the same consequences
on misbehavior

Structural Holes

Good

Structural holes

advantage in organizations

• "influence"

• informational advantage

• "creativity"

Bad things:

• manipulation

• "gate keeper"

latent
power

Social Capital

"the ability of actors to secure benefits by virtue of membership in social networks or other social structures"
(A. Portes)

individuals vs organizations

Social Capital and Triadic Closure

Burt : " social capital
is a tension between
closure and brokerage "

bounding

vs

bridging

Take home message

- social structures are facilitators of actions by individual and groups
- networks are at the heart of such discussions

Betweenness Measures and Graph Partitioning

Networks are made of **tightly-knit regions** connected by means of **sparser interconnections**

We do not have introduced so far a formal definition of such regions

although we have some useful measures and definitions:

- clustering coefficient
- local bridges
- neighborhood overlap
- triadic closure ...

the problem of finding denser regions in a network is called: **Graph Partitioning** or also **Community Detection**

Example 1: Co-authorship network



S_1 and S_2 are network scientists that co-authored at least one scientific paper

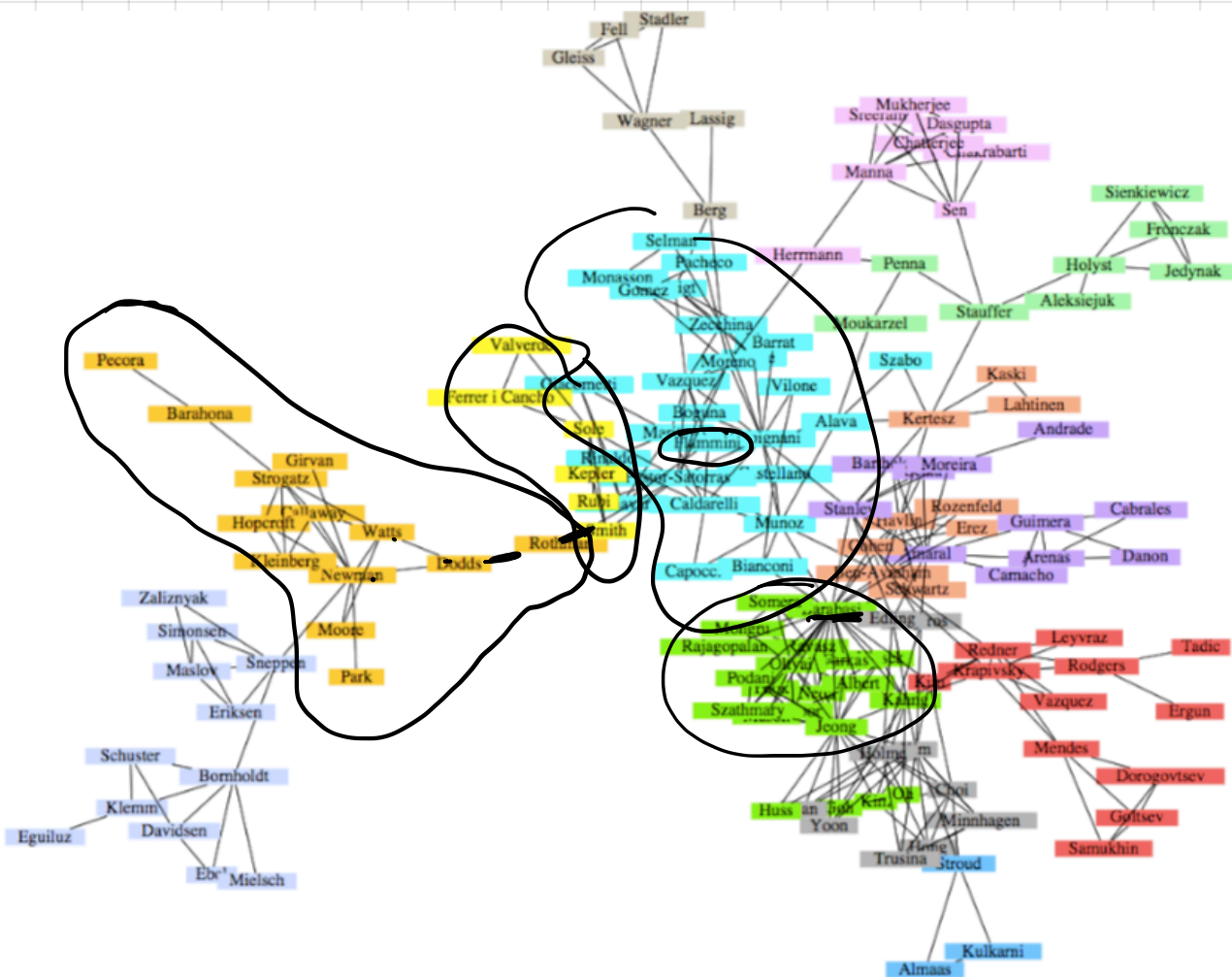


Figure 3.12: A co-authorship network of physicists and applied mathematicians working on networks [322]. Within this professional community, more tightly-knit subgroups are evident

Example 2: Zachary's Karate Club

can I use structure
itself to predict
the "fault line"

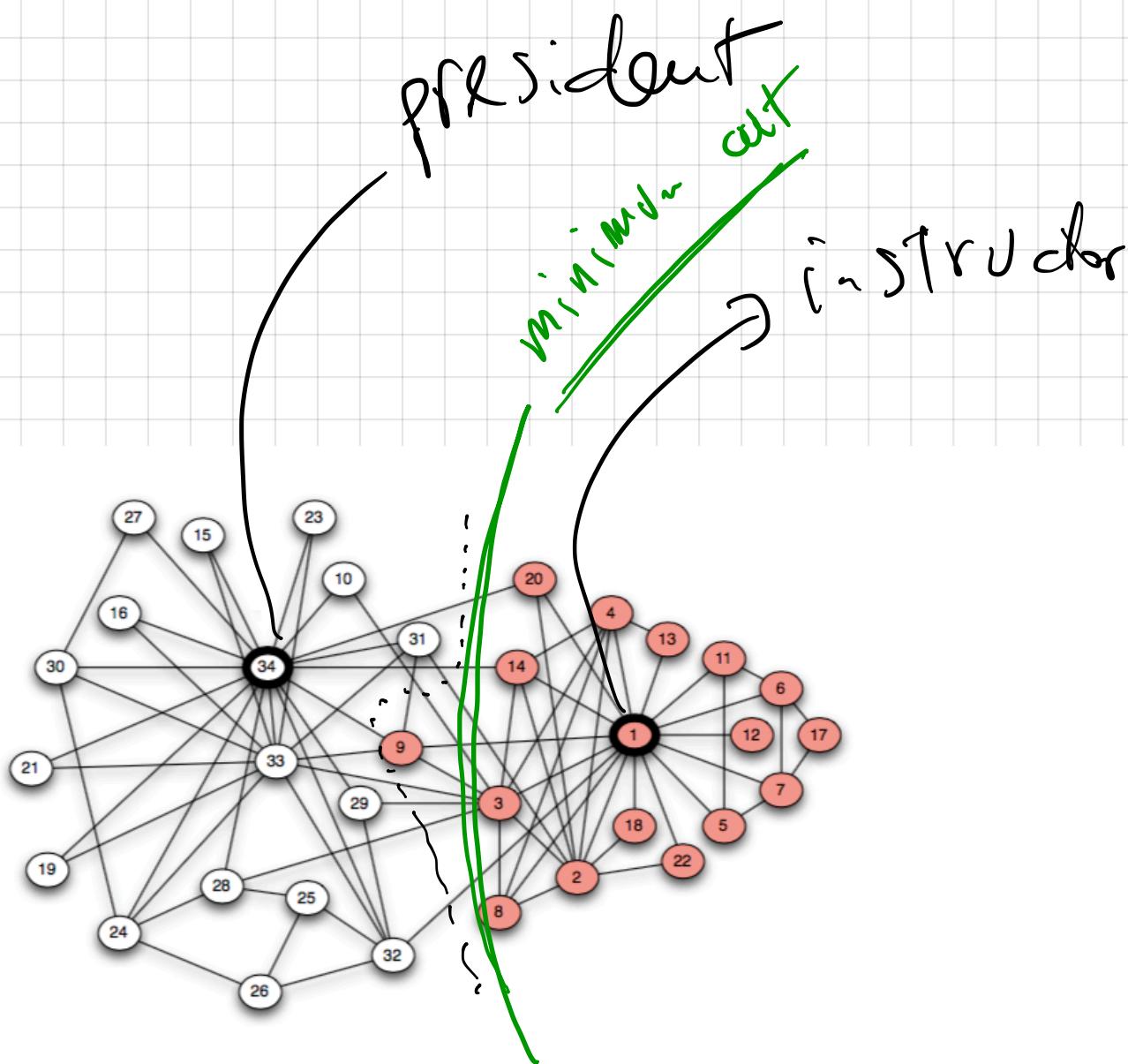
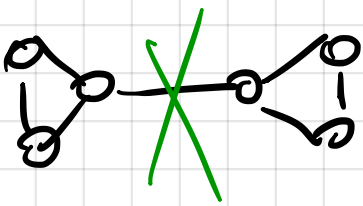


Figure 3.13: A karate club studied by Wayne Zachary [421] — a dispute during the course of the study caused it to split into two clubs. Could the boundaries of the two clubs be predicted from the network structure?

General Approaches to Graph Partitioning

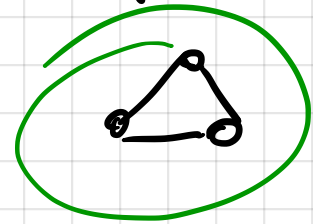
divisive
top-down

"Spanning links"
goal: remove

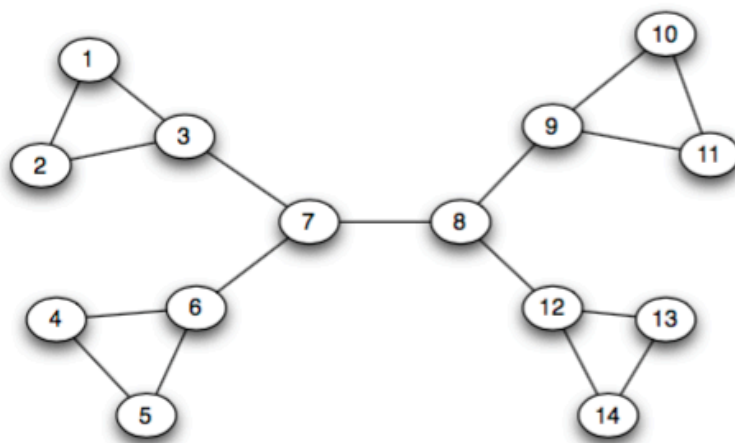


agglomerative
bottom-up

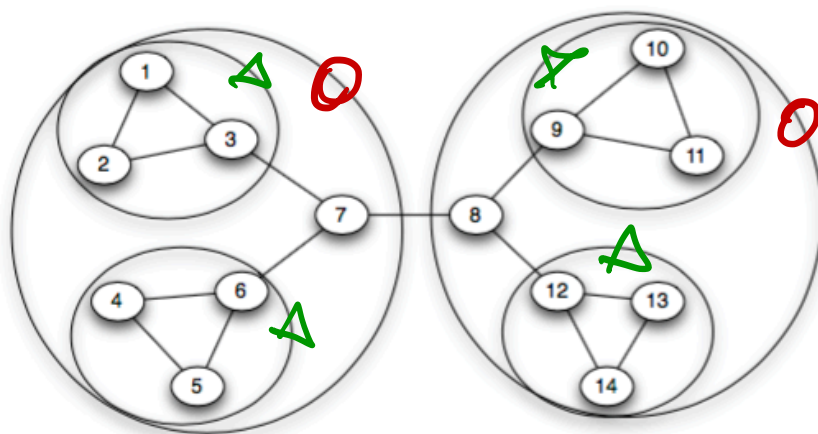
nodes merged
with other
nodes to form
groups



Nested structure



(a) A sample network

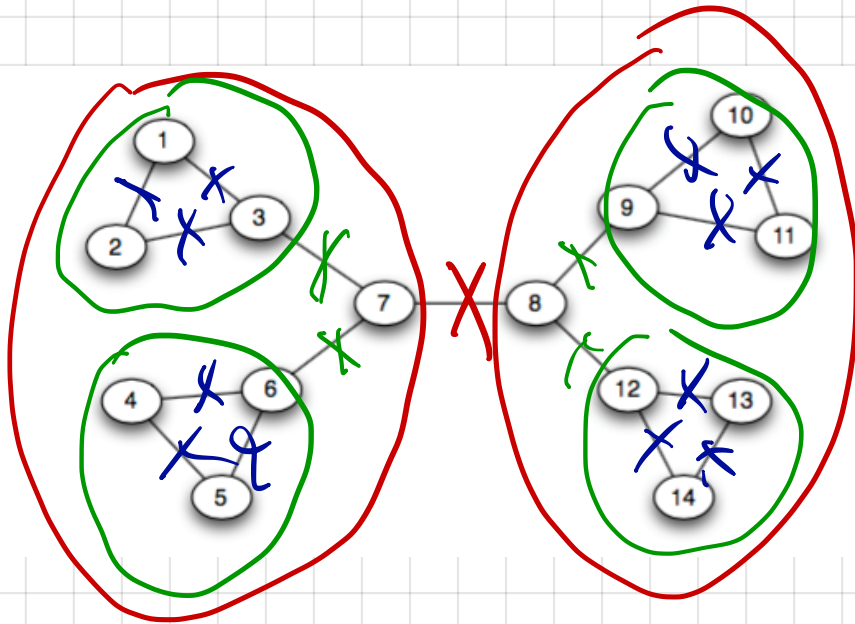
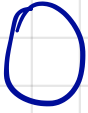


(b) Tightly-knit regions and their nested structure

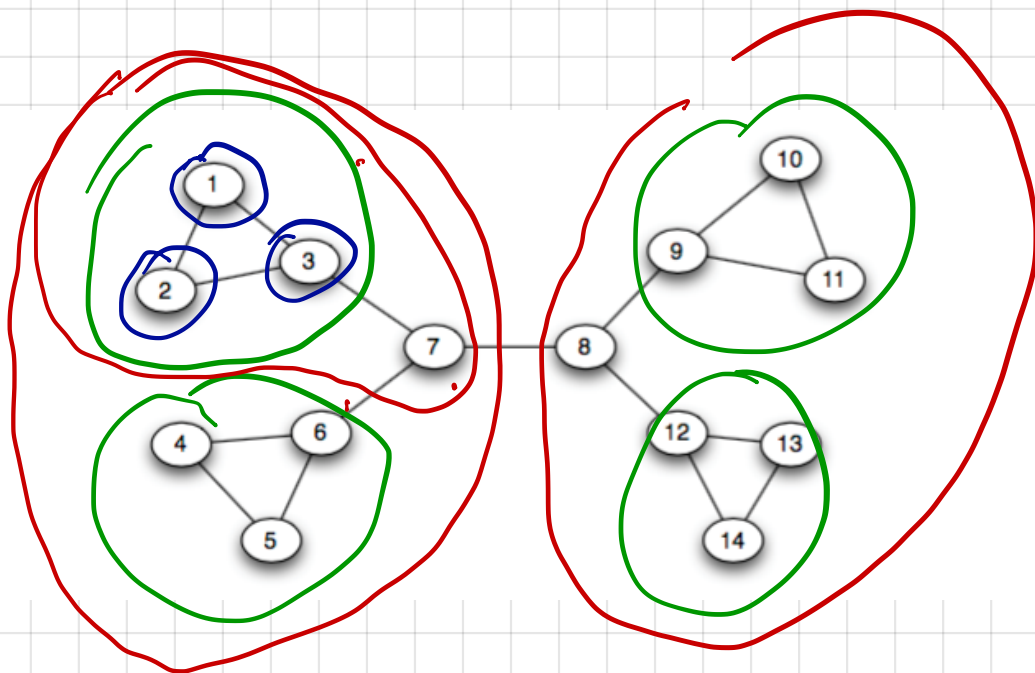
Figure 3.14: In many networks, there are tightly-knit regions that are intuitively apparent, and they can even display a *nested* structure, with smaller regions nesting inside larger ones.

Given - Newman

Divisive method



Agglomerative Method



Which edge to remove first
in *divisive* methods?

removing local bridges...
if disconnected: distances
will be
larger
($\rightarrow \infty$)

how many shortest paths
will be affected if
edge (x, y) is
removed?

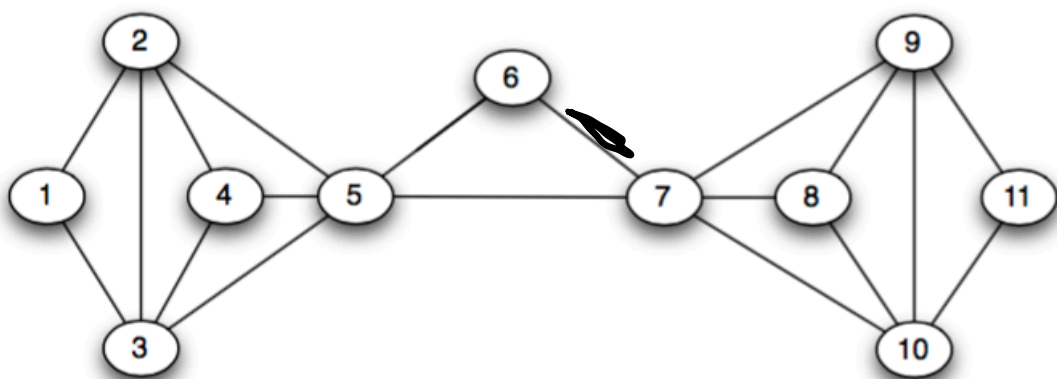
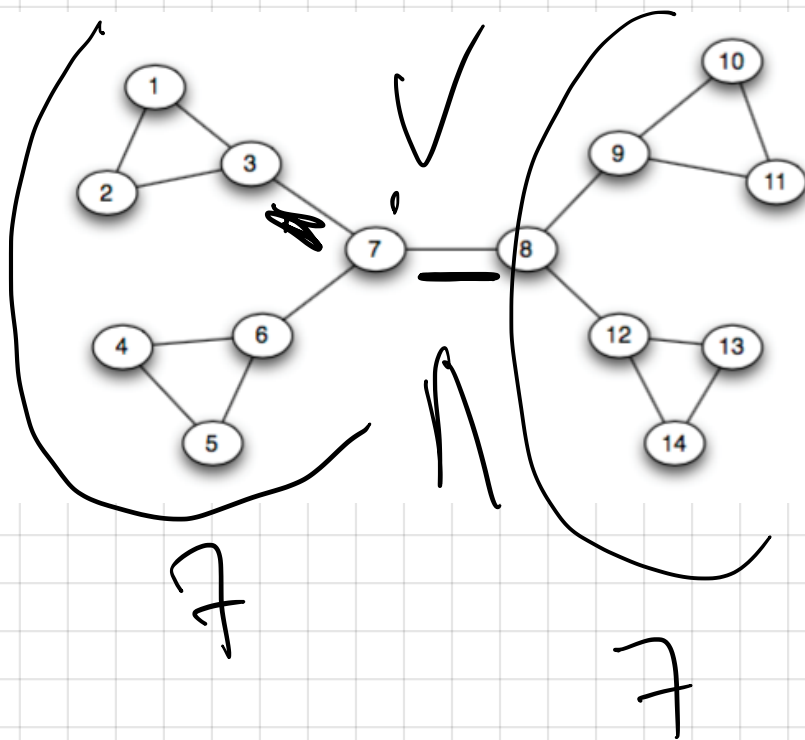


Figure 3.15: A network can display tightly-knit regions even when there are no bridges or local bridges along which to separate it.

Betweenness of edge (x, y)
 = number of shortest paths that cross through (x, y)

$$B(7, 8) = 49$$

$$B(3, 7) = 33$$



$$B(i) = \sum_{j < k} \frac{d_{jk}(i)}{d_{jk}}$$

normalized betweenness of a node

Some observations on Betweenness

local bridges

NO inversely correlated

"structural holes"?

"traffic" or "flow"

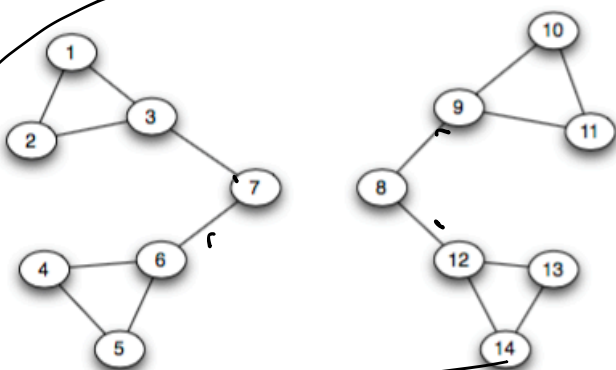
the Girvan - Newman Method

1. find edge(s) with highest betweenness

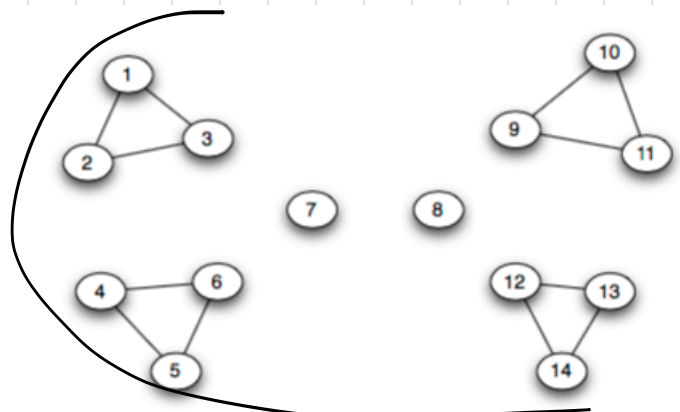
5. remove those edges

c. \rightarrow components

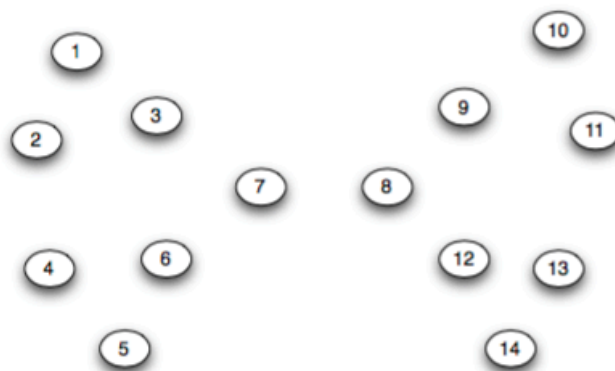
2. Re-calculate betweenness for all the edges



(a) Step 1



(b) Step 2



(c) Step 3

Figure 3.16: The steps of the Girvan-Newman method on the network from Figure 3.14(a).

Another G-N example

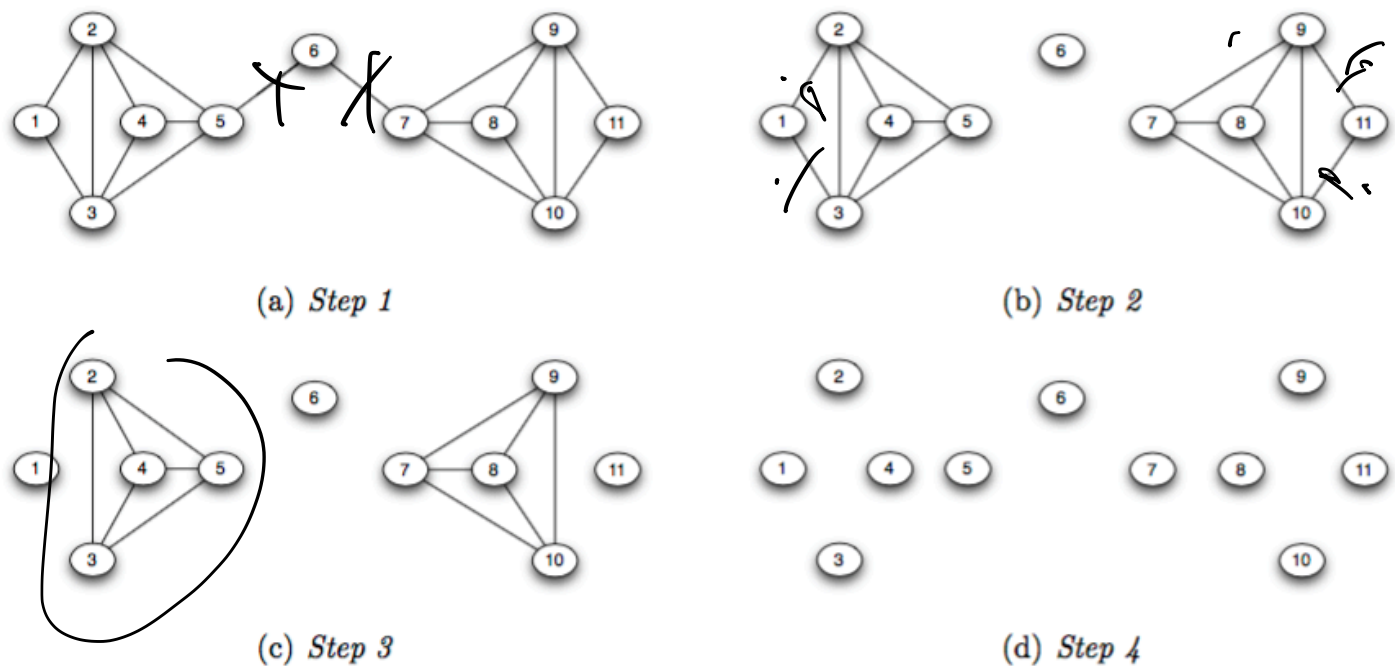
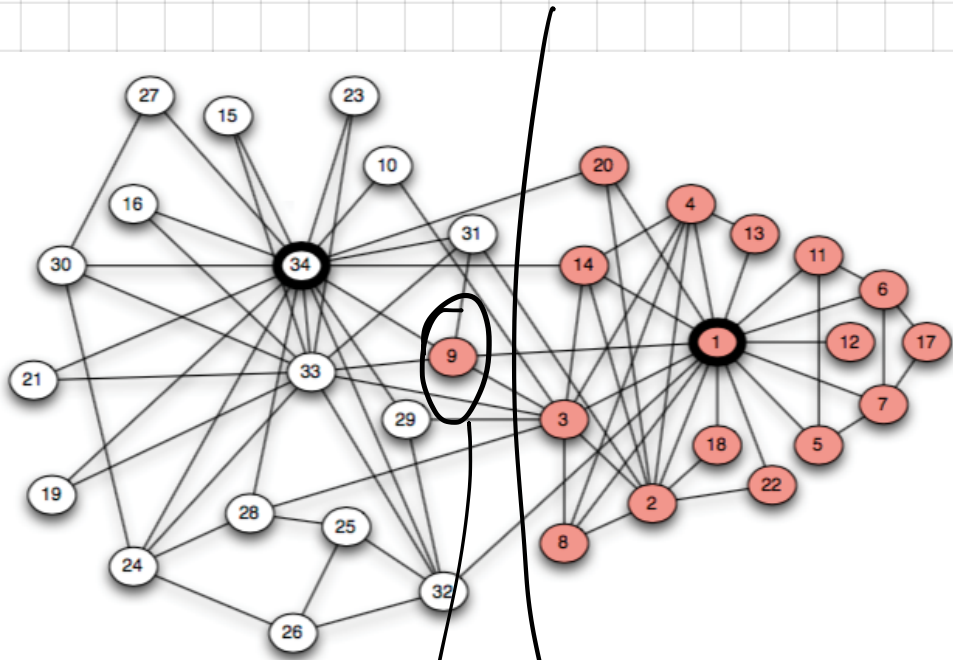


Figure 3.17: The steps of the Girvan-Newman method on the network from Figure 3.15.

We do not need
to find local bridges
anymore

⇒ betweenness is an
excellent approximation
that is calculated
through a global property of
the network

Applying G-N method to Karate Club



~minimum cut~

?

Computing Betweenness Values

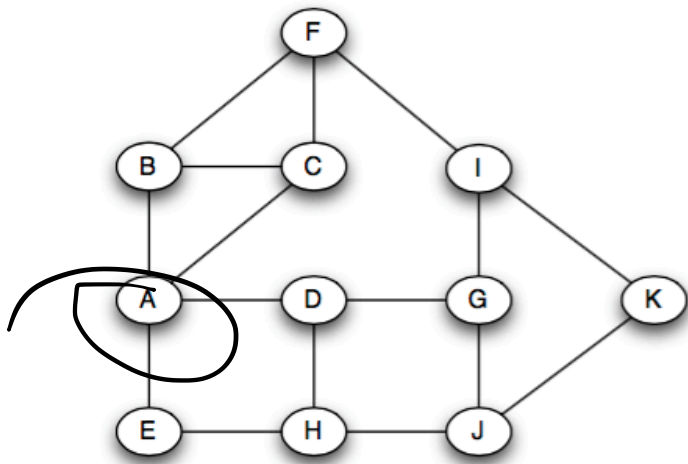
BFS

∀ node A

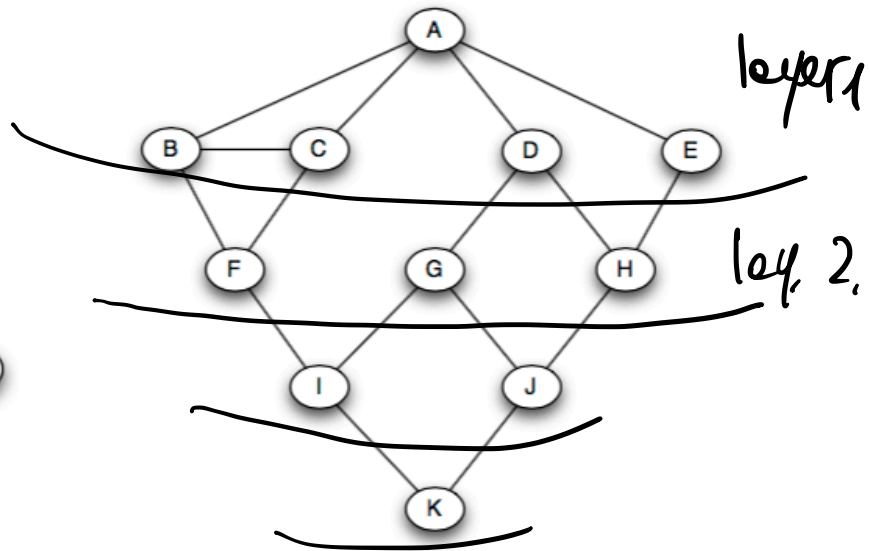
1. perform BFS
2. determine the number of shortest paths from A
3. calculate the amount of "flow" from A

First step: BFS from node A

for shown



(a) A sample network



(b) Breadth-first search starting at node A

Figure 3.18: The first step in the efficient method for computing betweenness values is to perform a breadth-first search of the network. Here the results of breadth-first from node A are shown; over the course of the method, breadth-first search is performed from each node in turn.

Second step: Counting Shortest Paths

go from - of

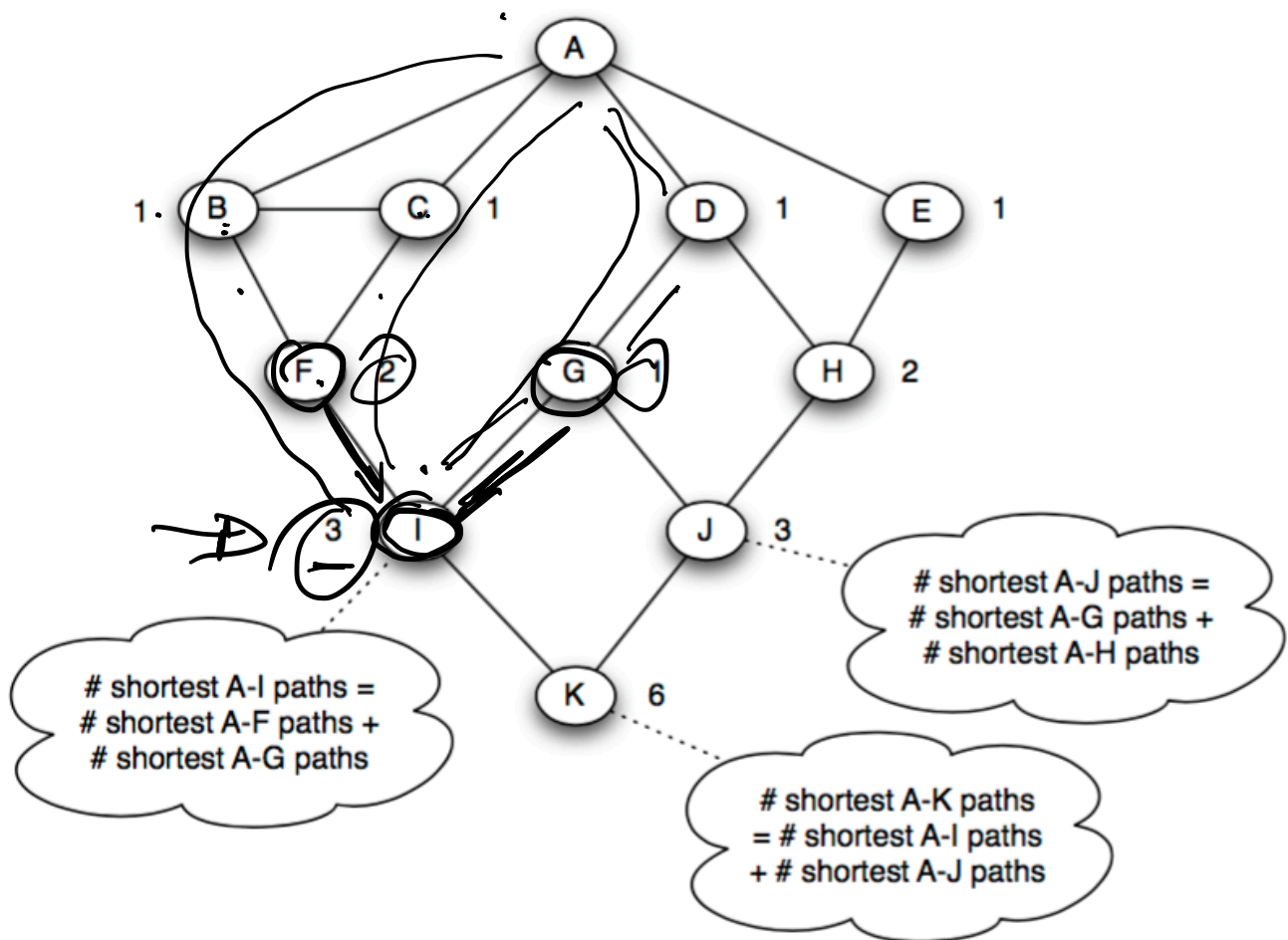


Figure 3.19: The second step in computing betweenness values is to count the number of shortest paths from a starting node A to all other nodes in the network. This can be done by adding up counts of shortest paths, moving downward through the breadth-first search structure.

Final Step: Computing the "flow"

bottom-up

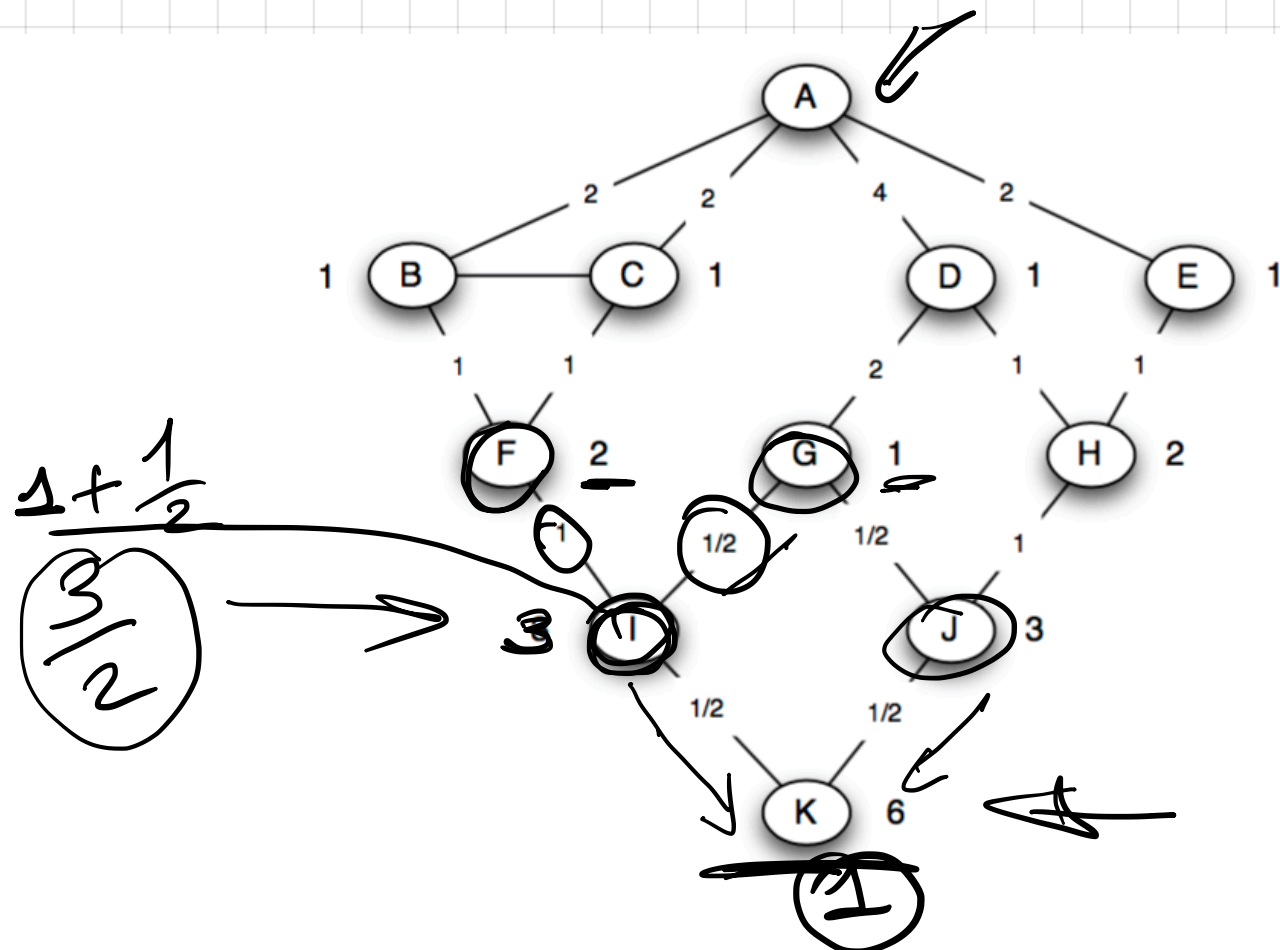


Figure 3.20: The final step in computing betweenness values is to determine the flow values from a starting node A to all other nodes in the network. This is done by working up from the lowest layers of the breadth-first search, dividing up the flow above a node in proportion to the number of shortest paths coming into it on each edge.

Output: **Betweenness** values for each edge

Repeat \forall nodes in G

Sum up all the
"flow"
values

2

↓

betweenness values

Take Home Message

Given - Newman:

$$O(N^3)$$

dense
graph

$$O(N^2)$$

sparse
graphs

"Levens"

(modularity
algorithm in
Gephi)

$$O(N)$$

sparse
graph

Open research question
how to better
define groups in
large scale networks