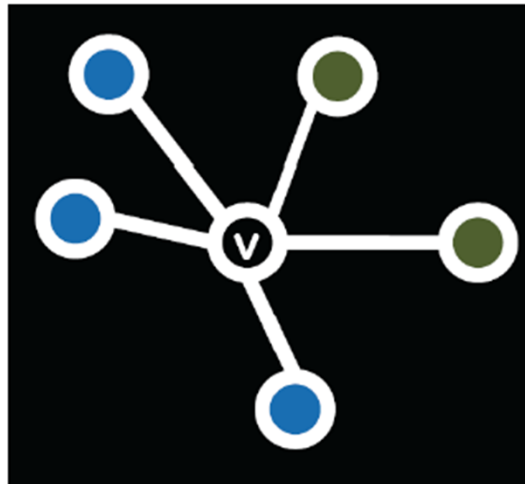


Cascading Behavior on Networks

Chapter 19

Game Theoretic Model of Cascades

- **Based on 2 player coordination game**
 - 2 players – each chooses technology A or B
 - Each person can only adopt **one** “behavior”, **A or B**
 - You gain more payoff if your friend has adopted the **same** behavior as you



Local view of the network of node **v**

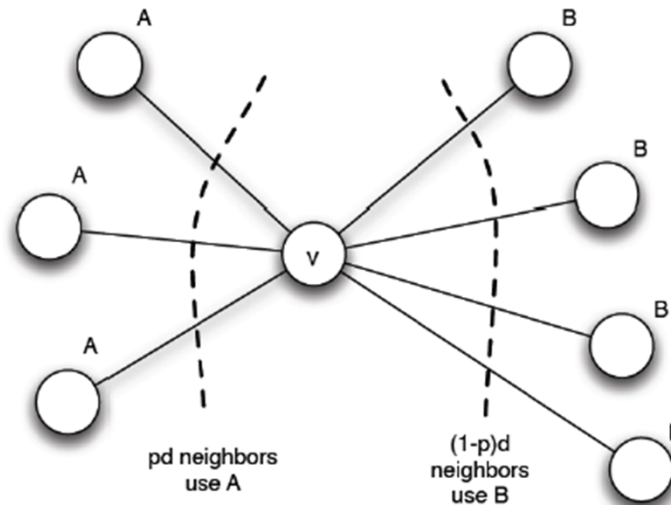
Example: VHS vs. BetaMax



Example: BlueRay vs. HD DVD



Calculation of Node v



Threshold:

v chooses **A** if

$$p > \frac{b}{a+b} = q$$

p ... frac. v 's nbrs. with A
 q ... payoff threshold

- Let v have d neighbors
- Assume fraction p of v 's neighbors adopt **A**
 - $Payoff_v = a \cdot p \cdot d$ if v chooses A
 - $= b \cdot (1-p) \cdot d$ if v chooses B
- **Thus: v chooses A if: $a \cdot p \cdot d > b \cdot (1-p) \cdot d$**

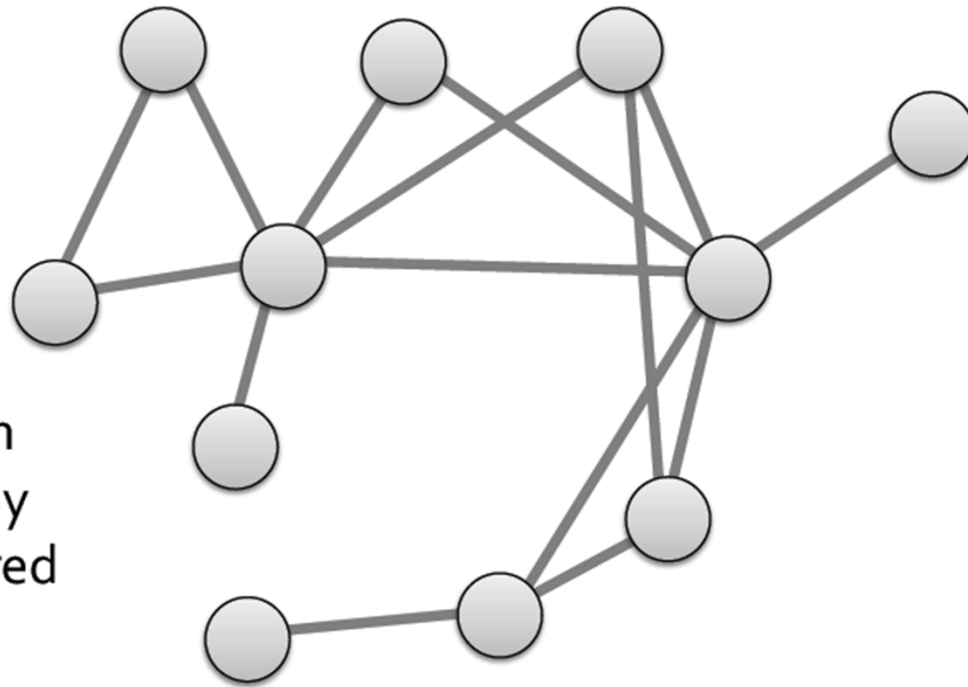
Example Scenario

Scenario:

- Graph where everyone starts with **B**
- Small set **S** of early adopters of **A**
 - Hard-wire **S** – they keep using **A** no matter what payoffs tell them to do
- **Assume payoffs are set in such a way that nodes say:**
If more than 50% of my friends take A
I'll also take A
This means: $a = b - \epsilon$ ($\epsilon > 0$, small positive constant)
and $q > 1/2$

Example Scenario

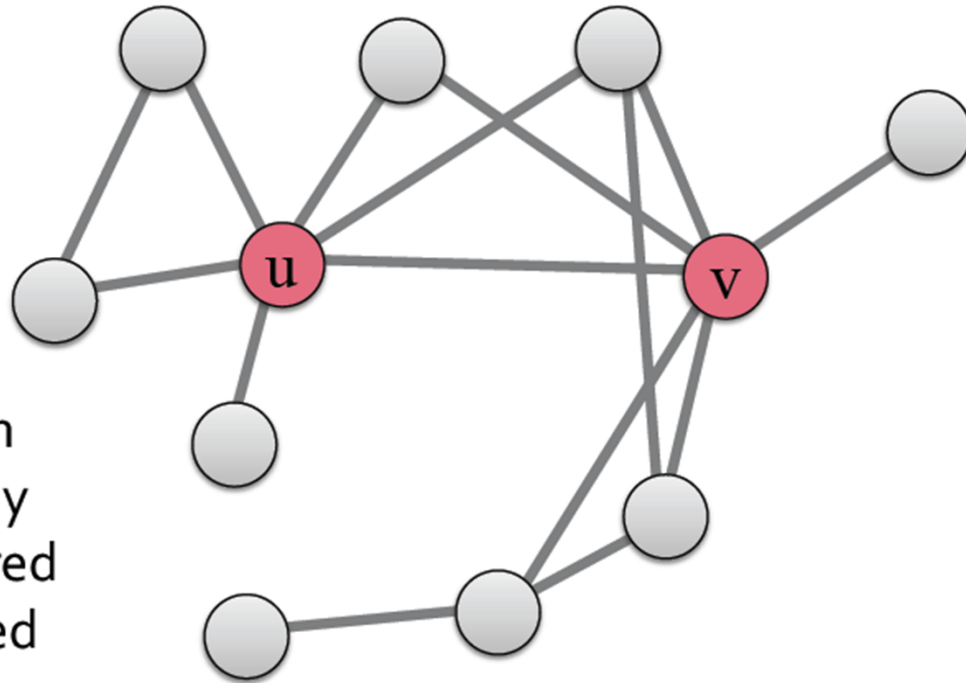
$$S = \{u, v\}$$



If **more** than
 $q=50\%$ of my
friends are red
I'll be red

Example Scenario

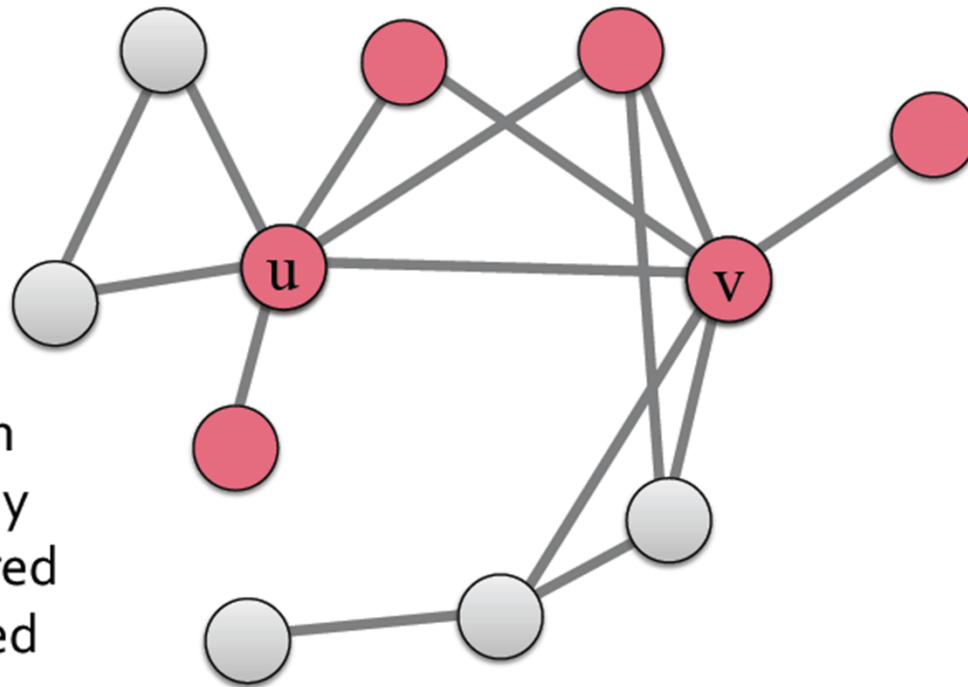
$$S = \{u, v\}$$



If **more** than
q=50% of my
friends are red
I'll also be red

Example Scenario

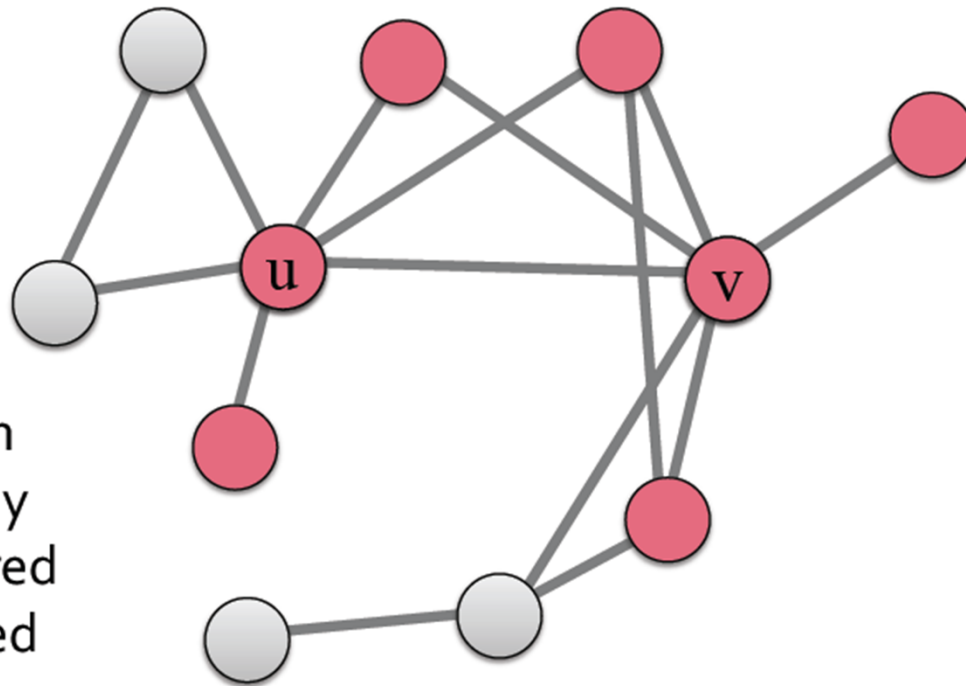
$$S = \{u, v\}$$



If **more** than
 $q=50\%$ of my
friends are red
I'll also be red

Example Scenario

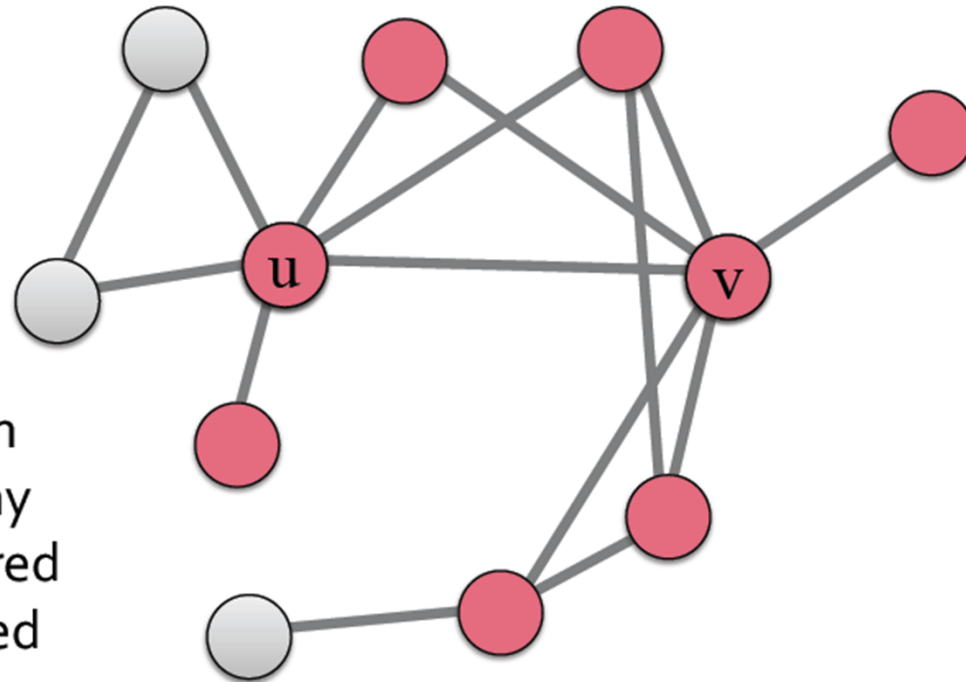
$$S = \{u, v\}$$



If **more** than
 $q=50\%$ of my
friends are red
I'll also be red

Example Scenario

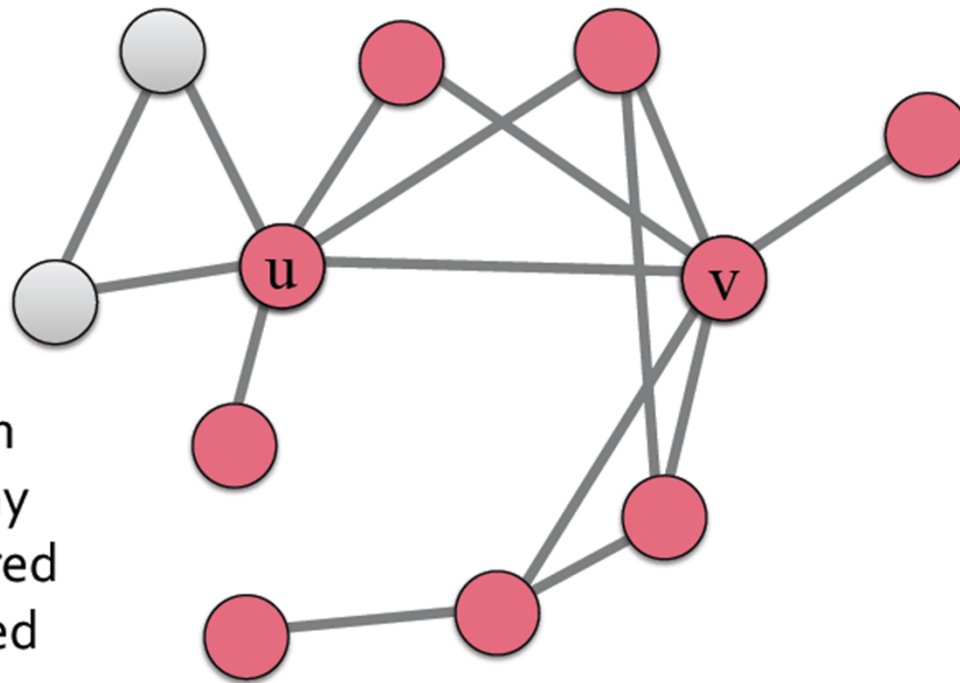
$$S = \{u, v\}$$



If **more** than
 $q=50\%$ of my
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I'll also be red

Example Scenario

$$S = \{u, v\}$$



If **more** than
 $q=50\%$ of my
friends are red
I'll also be red

Monotonic Spreading

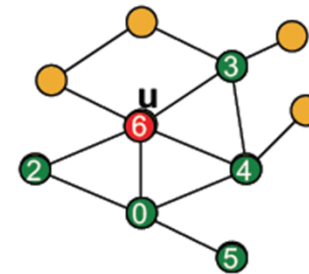
- **Observation: Use of A spreads monotonically**
(Nodes only switch $B \rightarrow A$, but never back to B)

- **Why?** Proof sketch:

- **Nodes keep switching from B to A: $B \rightarrow A$**
- Now, suppose some node switched back from $A \rightarrow B$, consider the **first** node u (not in S) to do so (say at time t)
- Earlier at some time t' ($t' < t$) the same node u switched $B \rightarrow A$
- So at time t' u was above threshold for A
- But up to time t no node switched back to B , so node u could only have more neighbors who used A at time t compared to t' .

There was no reason for u to switch at the first place!

!! Contradiction !!



Infinite Graphs

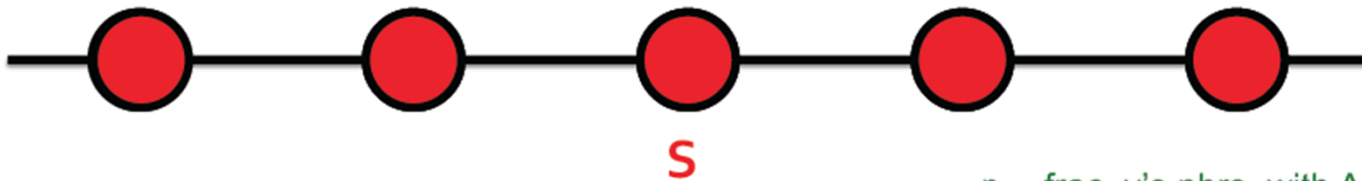
- Consider infinite graph G

- (but each node has finite number of neighbors!)

- We say that a finite set S causes a cascade in G with **threshold** q if, when S adopts A , eventually **every node in G adopts A**

- Example: **Path**

If $q < 1/2$ then cascade occurs



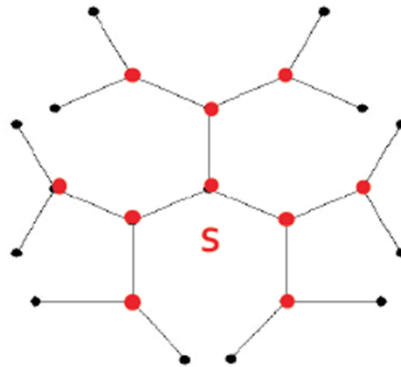
v chooses A if $p > q$

$$q = \frac{b}{a+b}$$

p ... frac. v 's nbrs. with A
 q ... payoff threshold

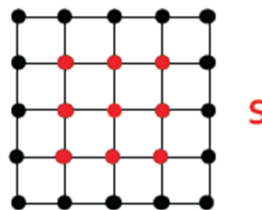
Infinite Graphs

- Infinite Tree:



If $q < 1/3$ then
cascade occurs

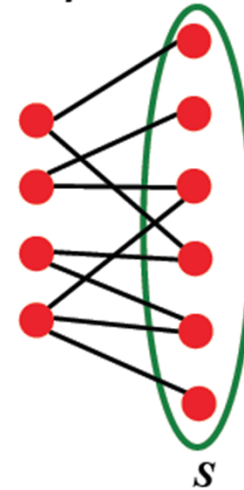
- Infinite Grid:



If $q < 1/4$ then
cascade occurs

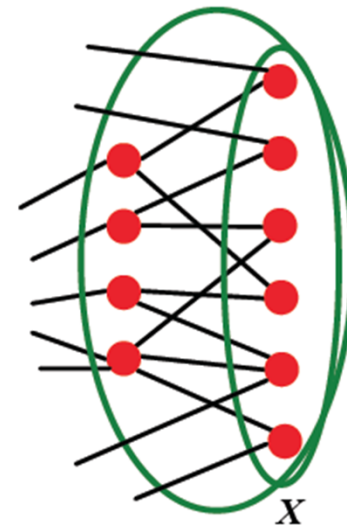
Cascade Capacity

- Def:
 - The **cascade capacity** of a graph G is the **largest q** for which some **finite set S** can cause a **cascade**
- Fact:
 - There is no (infinite) G where cascade capacity $> \frac{1}{2}$
- Proof idea:
 - Suppose such G exists: $q > \frac{1}{2}$, finite S causes cascade
 - **Show contradiction:** Argue that nodes stop switching after a finite # of steps



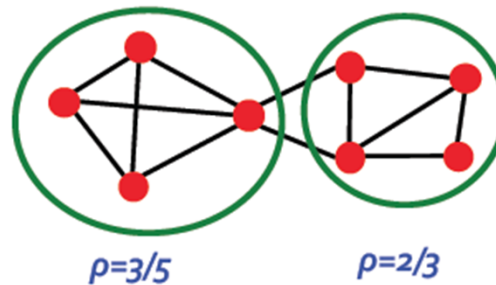
Cascade Capacity

- **Fact:** There is no G where cascade capacity $> \frac{1}{2}$
- **Proof sketch:**
 - Suppose such G exists: $q > \frac{1}{2}$, finite S causes cascade
 - **Contradiction:** Switching stops after a finite # of steps
 - Define “potential energy”
 - Argue that it starts finite (non-negative) and strictly decreases at every step
 - “Energy” := $|d^{\text{out}}(X)|$
 - $|d^{\text{out}}(X)|$:= # of outgoing edges of active set X
 - The only nodes that switch have a strict majority of its neighbors in S
 - $|d^{\text{out}}(X)|$ strictly decreases
 - It can do so only for a finite number of steps



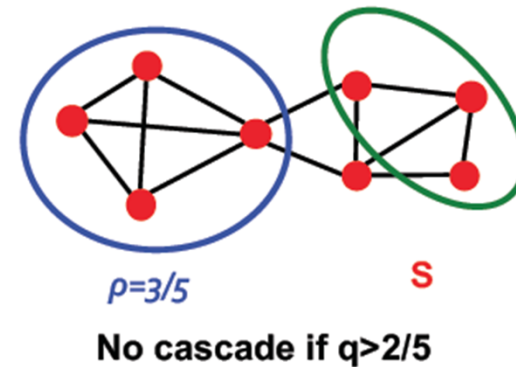
Stopping Cascades

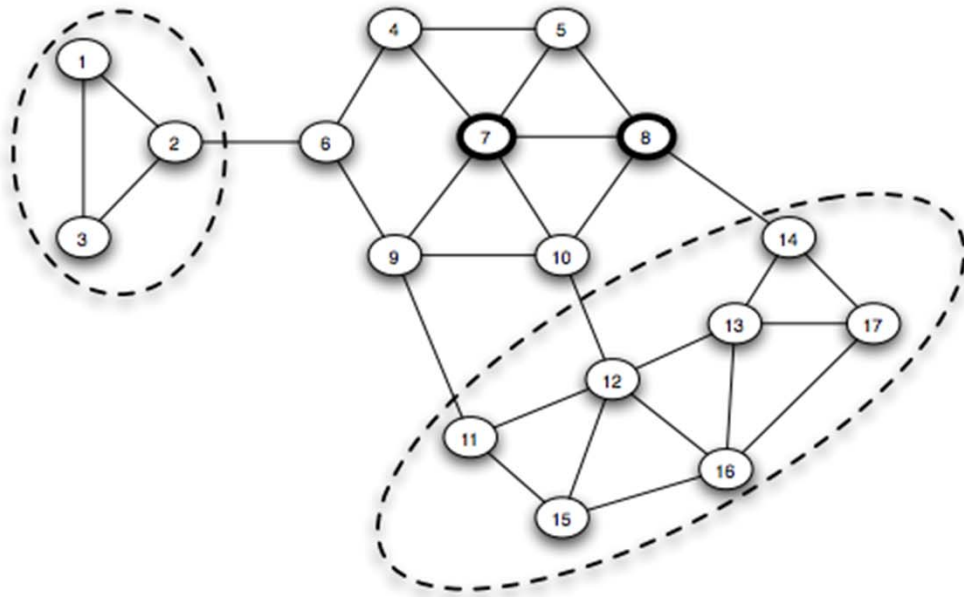
- **What prevents cascades from spreading?**
- **Def:** **Cluster of density ρ** is a **set of nodes C** where **each node** in the set has at least ρ fraction of edges in C



Stopping Cascades

- Let S be an initial set of adopters of A
- All nodes apply threshold q to decide whether to switch to A
- **Two facts:**
 - 1) If $G \setminus S$ contains a cluster of density $>(1-q)$ then S cannot cause a cascade
 - 2) If S fails to create a cascade, then there is a cluster of density $>(1-q)$ in $G \setminus S$





Claim: Consider a set of initial adopters of behavior A, with a threshold of q for nodes in the remaining network to adopt behavior A.

Figure 19.7

- (i) If the remaining network contains a cluster of density greater than $1 - q$, then the set of initial adopters will not cause a complete cascade.*
- (ii) Moreover, whenever a set of initial adopters does not cause a complete cascade with threshold q , the remaining network must contain a cluster of density greater than $1 - q$.*

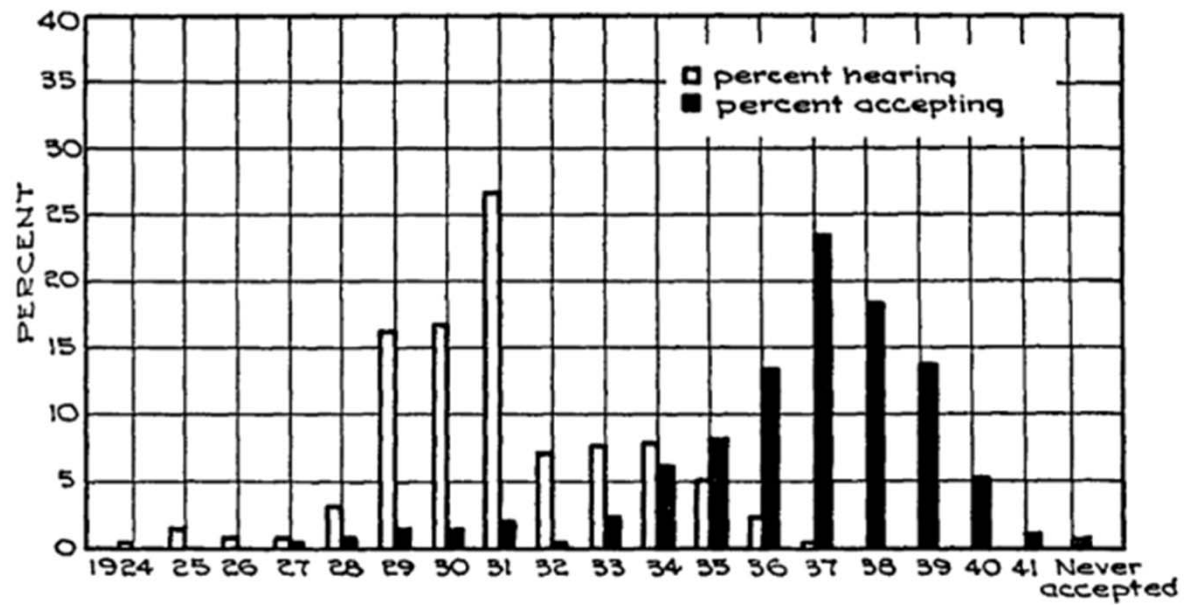


Figure 19.10: The years of first awareness and first adoption for hybrid seed corn in the Ryan-Gross study. (Image from [358].)

Randomization to Conditions

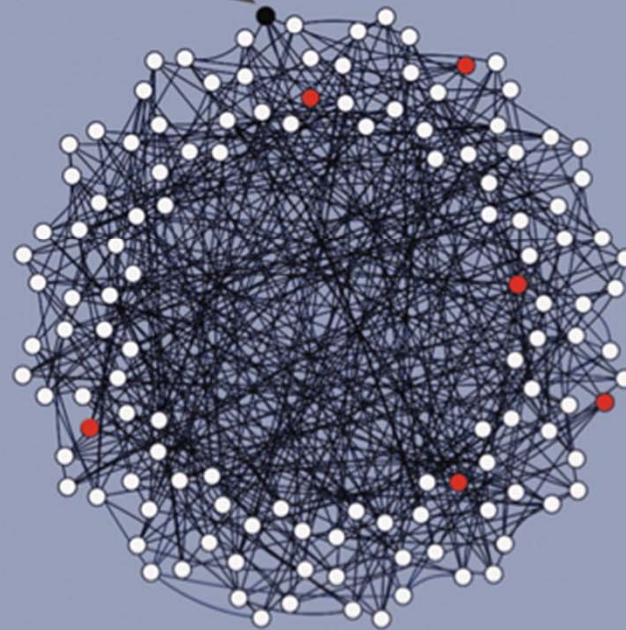
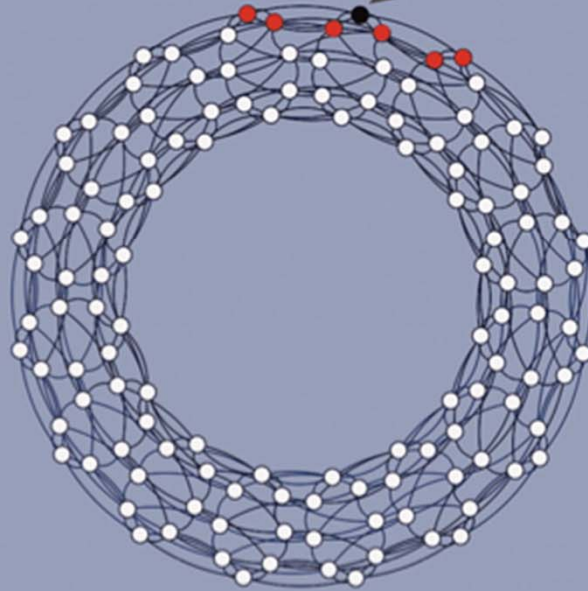
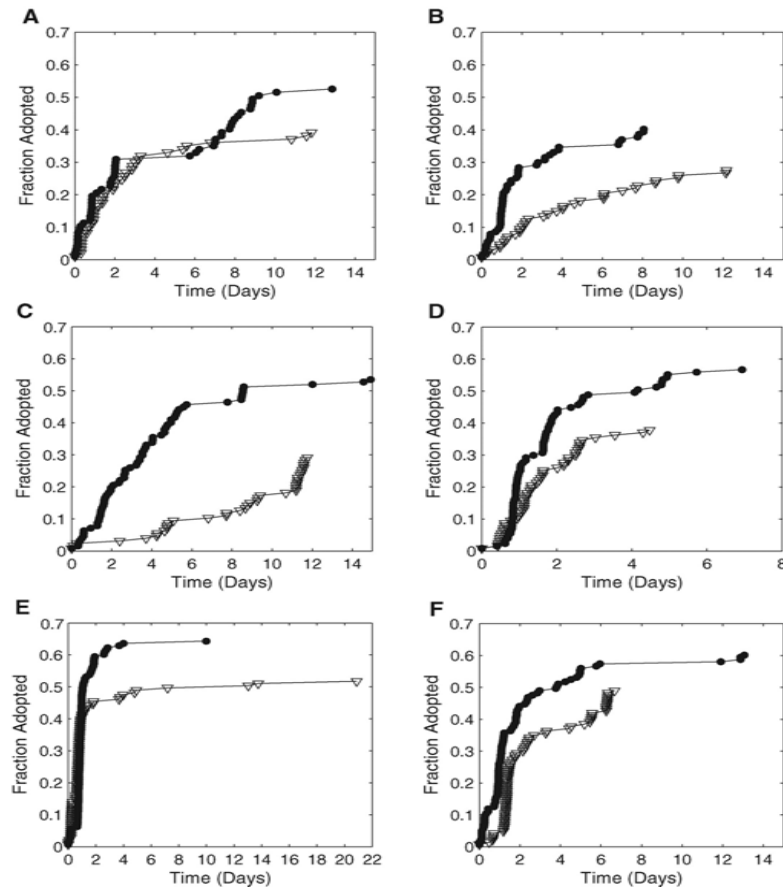
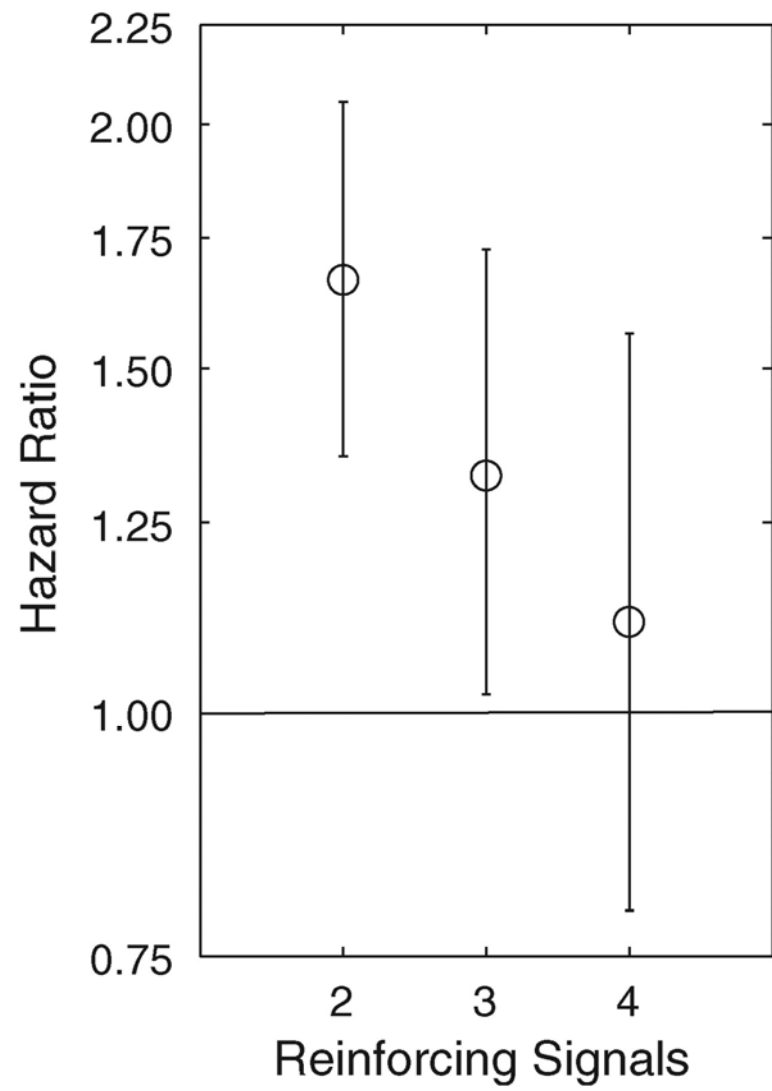
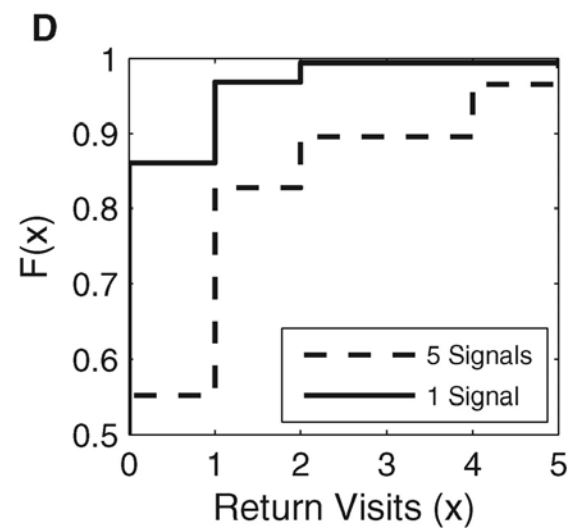
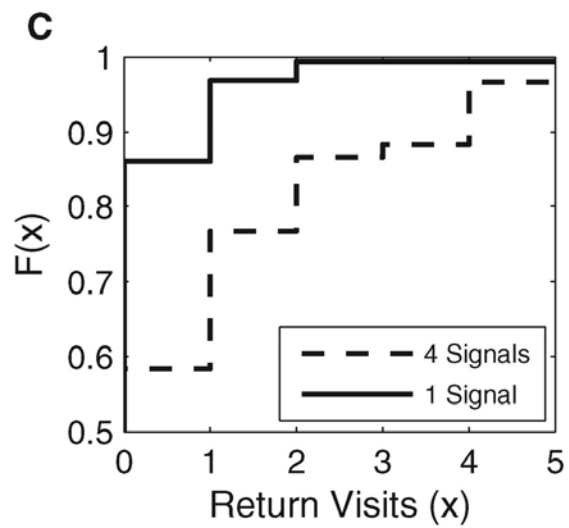
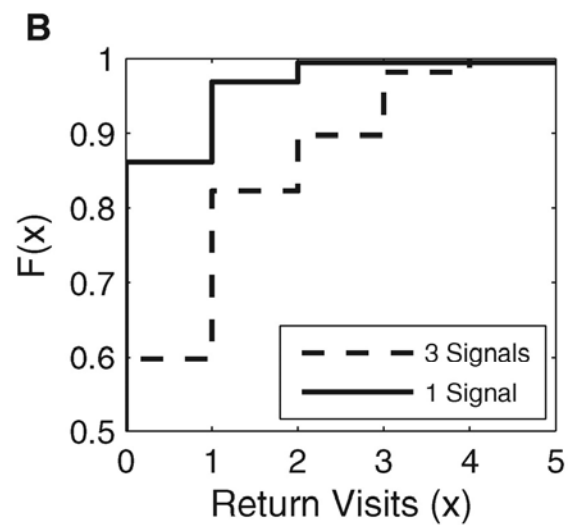
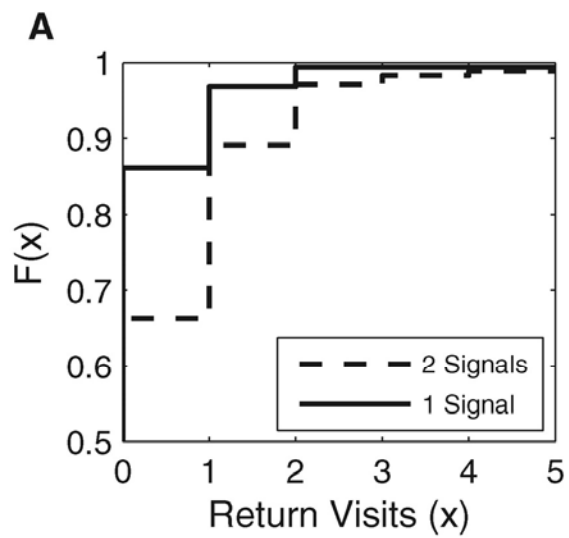


Fig. 2 Time series showing the adoption of a health behavior spreading through clustered-lattice (solid black circles) and random (open triangles) social networks.







**Extending the Model:
Allow People to Adopt A and B**