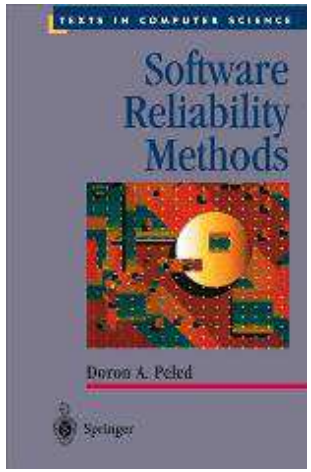
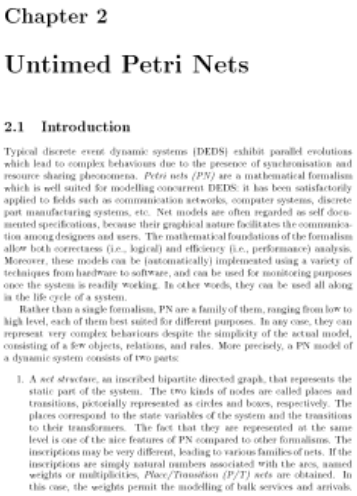
VPC 19-20
Algebre dei processi

Prof.ssa Susanna Donatelli
Universita' di Torino
www.di.unito.it
susi@di.unito.it

Reference material books:





Chapter 2

Untimed Petri Nets

2.1 Introduction

Typical discrete event dynamic systems (DEDS) exhibit parallel evolutions which lead to complex behaviours due to the presence of synchronisation and resource sharing phenomena. *Petri nets (PN)* are a mathematical formalism which is well suited for modelling concurrent DEDS: it has been satisfactorily applied to fields such as communication networks, computer systems, discrete part manufacturing systems, etc. Net models are often regarded as self documented specifications, because their graphical nature facilitates the communication among designers and users. The mathematical foundations of the formalism allow both correctness (i.e., logical) and efficiency (i.e., performance) analysis. Moreover, these models can be (automatically) implemented using a variety of techniques from hardware to software, and can be used for monitoring purposes once the system is readily working. In other words, they can be used all along in the life cycle of a system.

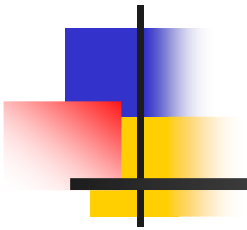
Rather than a single formalism, PN are a family of them, ranging from low to high level, each of them best suited for different purposes. In any case, they can represent very complex behaviours despite the simplicity of the actual model, consisting of a few objects, relations, and rules. More precisely, a PN model of a dynamic system consists of two parts:

1. A *net structure*, an inscribed bipartite directed graph, that represents the static part of the system. The two kinds of nodes are called places and transitions, pictorially represented as circles and boxes, respectively. The places correspond to the state variables of the system and the transitions to their transformers. The fact that they are represented at the same level is one of the nice features of PN compared to other formalisms. The inscriptions may be very different, leading to various families of nets. If the inscriptions are simply natural numbers associated with the arcs, named weights or multiplicities, *Place/Transition (P/T) nets* are obtained. In this case, the weights permit the modelling of bulk services and arrivals.

Notes of the EU-sponsored Jaca MATCH school

Prof. Doron A. Peled
(University of Warwick, UK)

Testo di Aceto et al su reactive systems:
http://www.cs.ioc.ee/yik/schools/win2007/ingolfsdottir/sv-book-part1.pdf



Acknowledgements

Transparencies adapted from the course notes and trasparencies of

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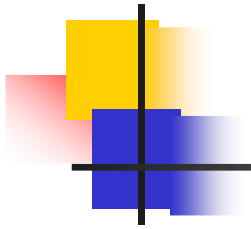
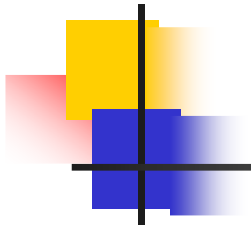


Process Algebra

(Book: Chapter 8)



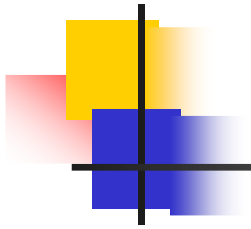
Algebra? Processi?

Def.: Sia I un insieme non vuoto. Diciamo struttura algebrica su I l'insieme

$$A = \{I, \alpha_1, \ldots\ldots \alpha_K\}$$

dove $\alpha_1, \ldots\ldots \alpha_K$ sono operazioni interne in I, rispettivamente a $n_1, \ldots, n_k$ argomenti

I è l'insieme dei processi

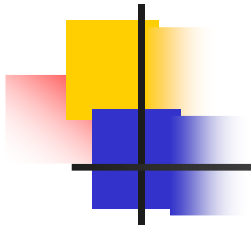


The Main Issue

Q: When are two models equivalent?

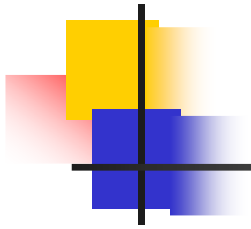
A: When they satisfy certain properties.

Q: Does this mean that the models have different executions?



What is process algebra?

- An abstract description for nondeterministic and concurrent systems.
- Focuses on the transitions observed rather than on the states reached.
- Interactions between independent processes as communication (no shared variables)
- Main correctness criterion: conformance between two models, but exhaustive state space generation and analysis is also possible
- Uses: system refinement, model checking, testing.
- *Defining algebraic laws for the process operators, which allow process expressions to be manipulated using equational reasoning*



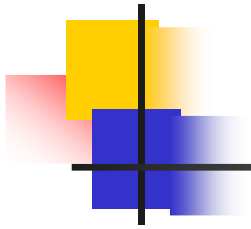
Process algebra ingredients

- *Act = {a,b,c, … }*, a set of actions
- $\forall a \in$ Act, $\exists$ co-action $\underline{a} \in$ Act
- τ-silent action, $\underline{\tau}=\tau$

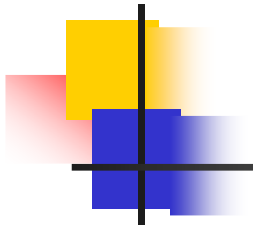
We shall define:

- a syntax for defining legal agents
- a semantics for defining how agents evolves
- equivalence rules to allow distinguishing agents

CCS - calculus of communicating systems [Milner]. Syntax



- *a*,*b*,*c*, … actions, A, B, C, …. - agents.
- $\underline{a}$,$\underline{b}$,$\underline{c}$, coactions of *a*,*b*,*c*. τ-silent action.
- 0 (o nil) - terminate.
- (E) – evaluation order
- *a*.E – prefixing. Execute a, then behave like E.
- E+E - nondeterministic choice.
- E || E - parallel composition.
- E\L - restriction: cannot use actions of L.
- E[f] - apply mapping function f between actions and agent names.



Conventions

- “.” has higher priority than “+”.
- “.0” or “.(0||0||…||0)” is omitted.

Semantics (proof rule and axioms). Structural Operational Semantics SOS

- $a.p \xrightarrow{a} p$
  - (a.p evolves with a in p)
- $p \xrightarrow{a} p' \vdash p+q \xrightarrow{a} p'$
  - (given that p evolves with a in p', then p+q evolves with a in p')
- $q \xrightarrow{a} q' \vdash p+q \xrightarrow{a} q'$
- $p \xrightarrow{a} p' \vdash p||q \xrightarrow{a} p'||q$
- $q \xrightarrow{a} q' \vdash p||q \xrightarrow{a} p||q'$
- $p \xrightarrow{a} p', q \xrightarrow{\underline{a}} q' \vdash p||q \xrightarrow{\tau} p'||q'$
- $p \xrightarrow{a} p' , a \notin R \vdash p\backslash R \xrightarrow{a} p'\backslash R$
- $p \xrightarrow{a} p' \vdash p[m] \xrightarrow{m(a)} p'[m]$

SOS rules and examples

$a.E \xrightarrow{a} E$ (Axiom)

Thus, $a.(b.(c//\underline{c})+d) \xrightarrow{a} (b.(c//\underline{c})+d)$.

Action Prefixing

$a.E \xrightarrow{a} E$ (Axiom)

Thus, $a.(\underbrace{b.(c \,|\!|\, \underline{c}) + d}) \xrightarrow{a} (\underbrace{b.(c \,|\!|\, \underline{c}) + d})$.

a. E --a--> E

Choice

$$\frac{E \xrightarrow{a} E'}{(E+F) \xrightarrow{a} E'} \qquad \frac{F \xrightarrow{a} F'}{(E+F) \xrightarrow{a} F'}$$

$$b.(c\,//\,\underline{c}) \xrightarrow{b} (c\,//\,\underline{c}).$$

Thus,

$$(b.(c\,//\,\underline{c}) + e) \xrightarrow{b} (c\,//\,\underline{c}).$$

If $E \xrightarrow{a} E'$ and $F \xrightarrow{a} F'$, then E+F has *a* nondeterministic choice.

Concurrent Composition

$$\frac{E \xrightarrow{a} E'}{E||F \xrightarrow{a} E'||F} \qquad \frac{F \xrightarrow{a} F'}{E||F \xrightarrow{a} E||F'}$$

evolutions independents

$$\frac{E \xrightarrow{a} E',\ F \xrightarrow{\underline{a}} F'}{E||F \xrightarrow{\tau} E'||F'}$$

synchronisation

$c \xrightarrow{c} nil$, $\underline{c} \xrightarrow{\underline{c}} nil$, then $c||\underline{c} \xrightarrow{\tau} nil||nil$,
$c||\underline{c} \xrightarrow{c} nil||\underline{c}$, $c||\underline{c} \xrightarrow{\underline{c}} c||nil$.

Restriction

$$\frac{E \xrightarrow{a} E',\quad a, \underline{a} \notin R}{E\backslash R \xrightarrow{a} E'\backslash R}$$

$$\frac{E \xrightarrow{\alpha} E',\ \bar{\alpha}, \alpha \notin R}{E/R \xrightarrow{\alpha} E'/R}$$

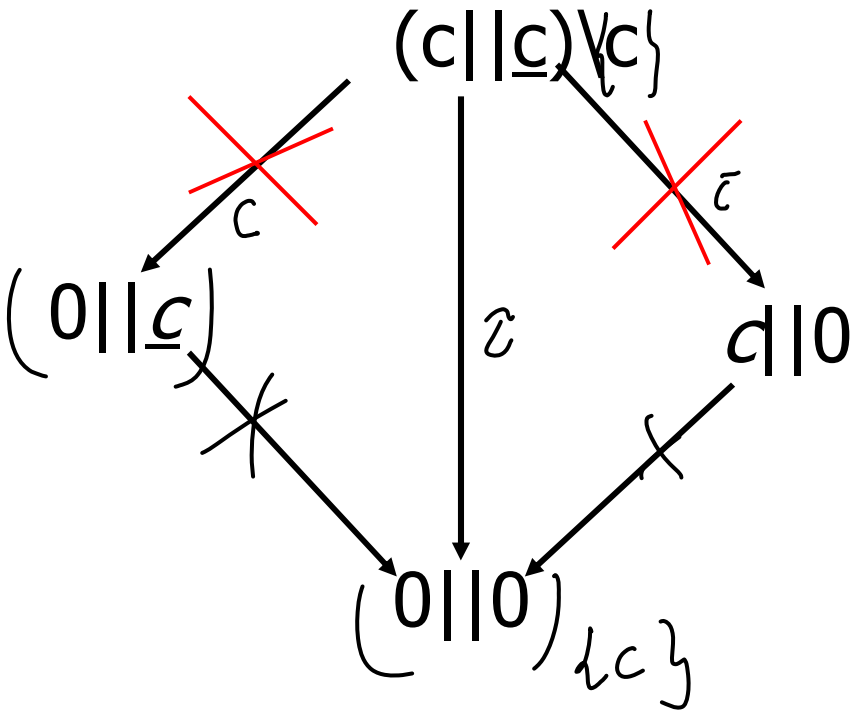
Example of interplay between parallel composition and restriction:

$c||\underline{c} \xrightarrow{\underline{c}} c||0 \qquad c||\underline{c} \xrightarrow{c} 0||\underline{c} \qquad c||\underline{c} \xrightarrow{\tau} 0||0$

then

$(c||\underline{c}) \setminus \{c\} \xrightarrow{\tau} (0||0) \setminus \{c\}$

Consequence of restriction



Relabeling

$$\frac{E \xrightarrow{a} E'}{E[m] \xrightarrow{m(a)} E'[m]}$$

$m : Act \to Act$

$a.b.c \longrightarrow b.c$

$(a.b.c)[a \to d, b \to a, c \to d] \xrightarrow{d = m(a)} b.c\,[\,]$

No axioms/rules for agent 0.

Equational definition of agents

The syntax presented only allows finite terms with finite behaviours

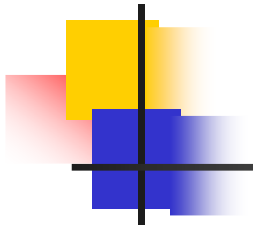
Process algebras usually allow an agent definition of the type

$P =_{def} a.(b.P)$

infinite behaviour (infinite traces)

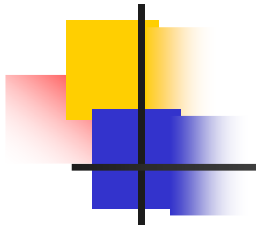
$Q =_{def} a||Q$

infinite terms and infinite behaviour



Equational Definition

$$\frac{E \xrightarrow{a} E',\ A = E}{A \xrightarrow{a} E'}$$

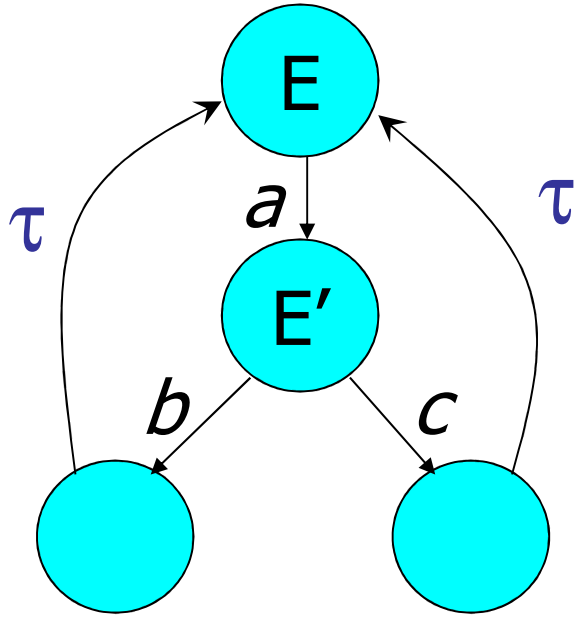


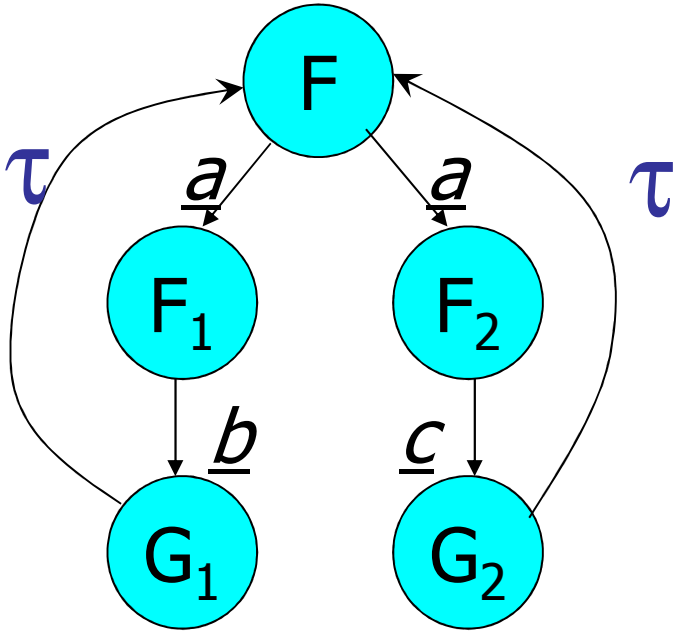
Derivation Graph

The derivatives D(E) of an agent E are the agents derived from E by applying the proof rules

The derivation graph DG(E) is the graph (N,A) where N = D(E) and (n1,n2)\in A iff n1→n2 according to the proof rules

Equational Definition and DG





E=*a*.(*b*.τ.E+c.τ.E)

F=*a*.*b*.τ.F+*a*.*c*.τ.F

$$\frac{E \xrightarrow{a} E',\ A=E}{A \xrightarrow{a} E'}$$

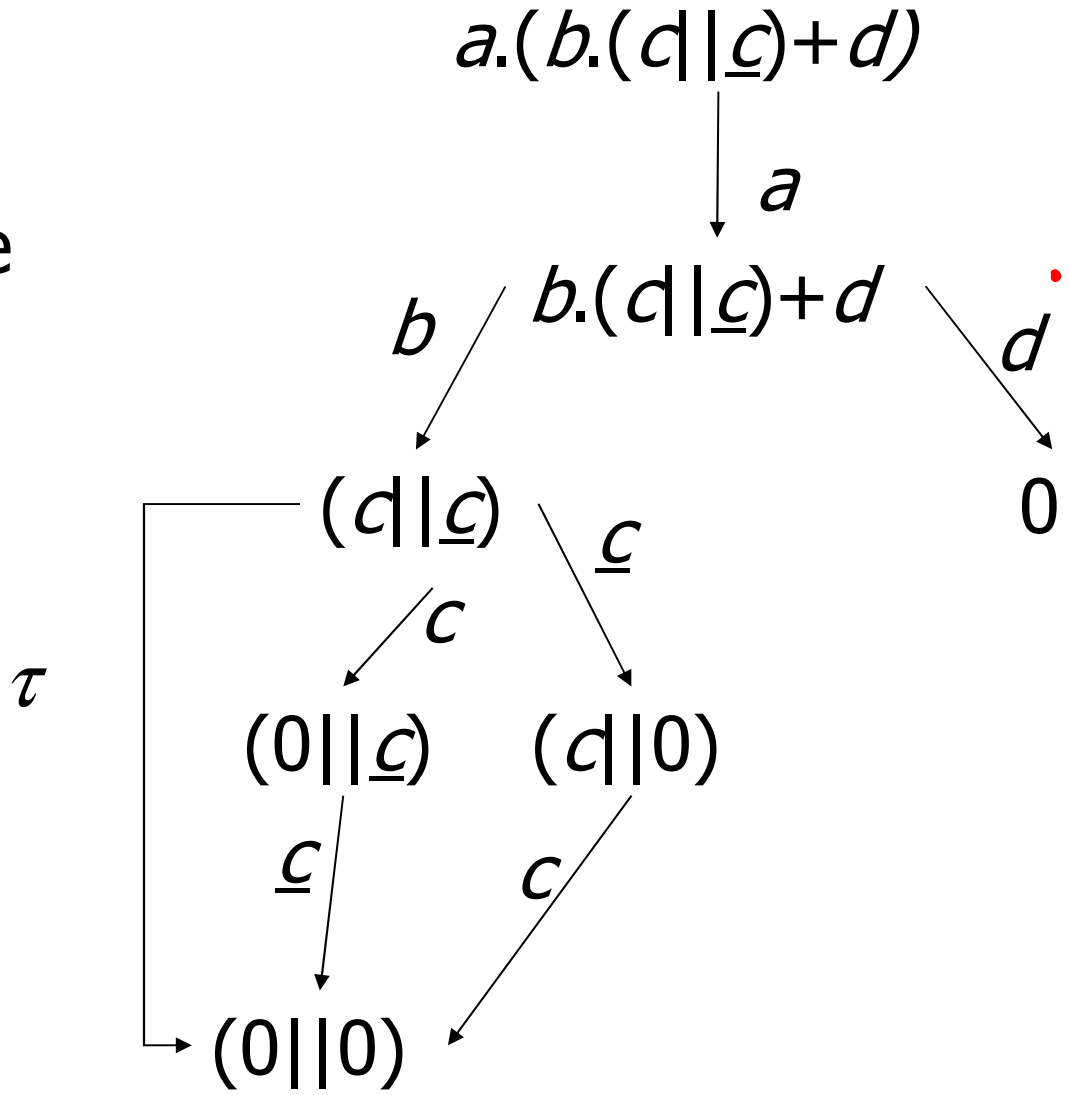
PROOF RULE

Derivations

Exercise: complete each arc with the name of the axiom/proof being applied

P=a.b.P

Q = a||Q

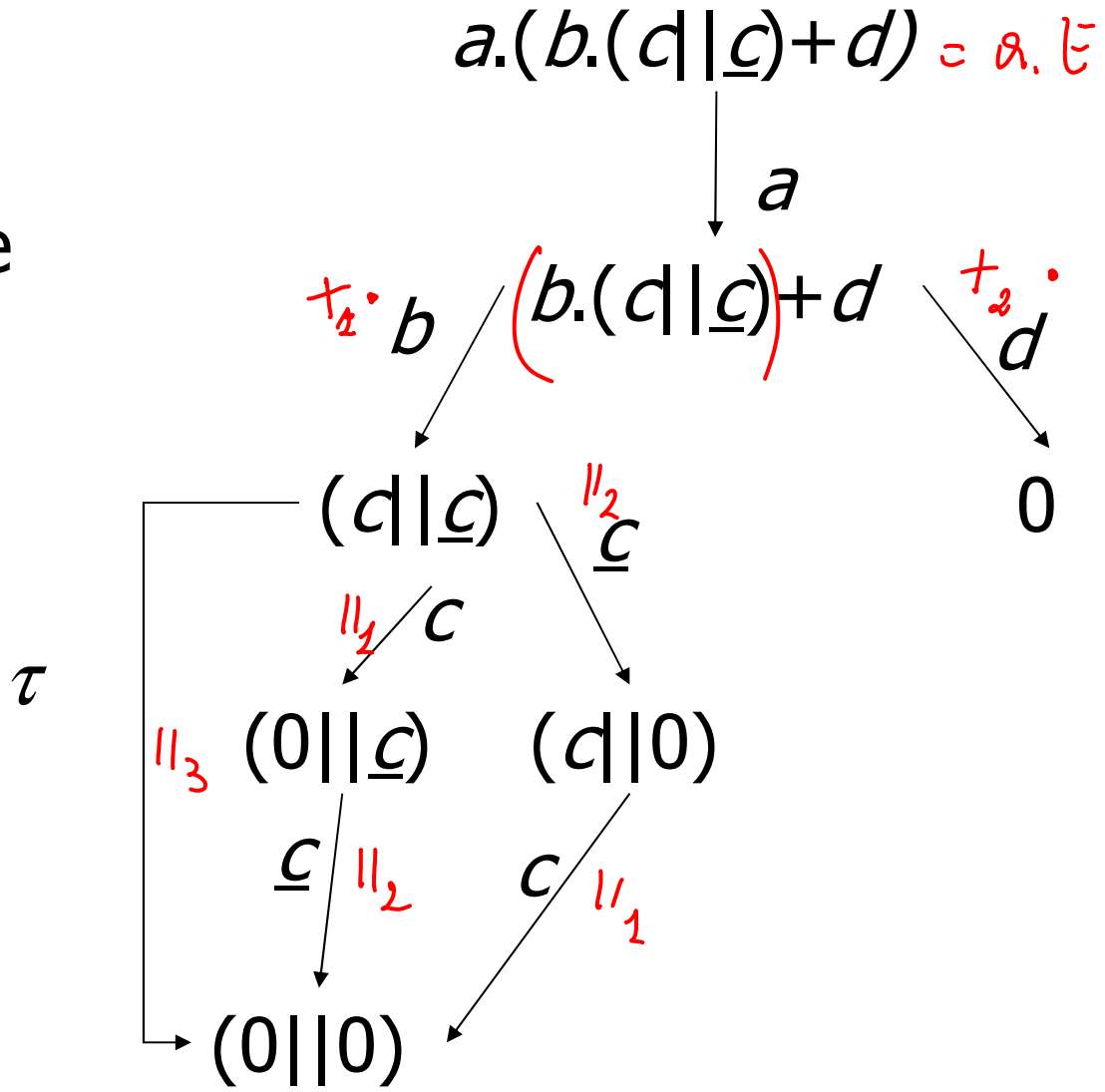


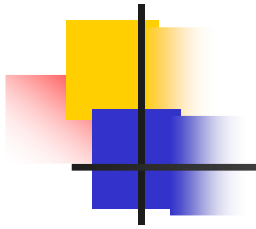
Derivations

Exercise: complete each arc with the name of the axiom/proof being applied

P=a.b.P

Q = a||Q





Example - buffer (dal Peled)

A one place buffer:

Empty1 = put.Full1

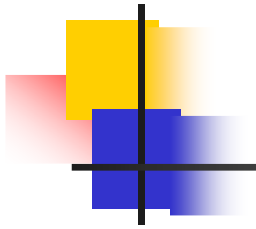
Full1 = get.Empty1

A two place buffer:

Empty2 = put.Half2

Half2 = put.Full2 + get.Empty2

Full2 = get.Half2



Example - buffer

Using the one place buffer:

Prod = loc.put.Prod Cons = get.loc.Cons

System = (Prod || Cons || Empty1) restricted over put and get

Using the two place buffer:

Prod = ….. Cons = …..

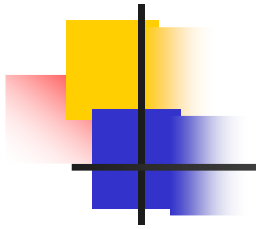
System = (Prod || Cons || Empty2) restricted over put and get

A two place buffer:

Empty2 = put.Half2

Half2 = put.Full2 + get.Empty2

Full2 = get.Half2



Example - buffer

A one place buffer:

Empty1 = put.Full1

Full1 = get.Empty1

A two place buffer using two one place buffer:

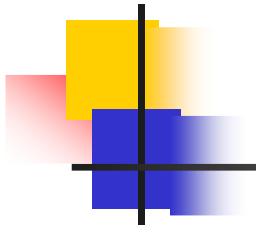
Empty2 = Empty1 || Empty1 correct?

A two place buffer:

Empty2 = put.Half2

Half2 = put.Full2 + get.Empty2

Full2 = get.Half2



Buffer

Solution and derivaton graph of

Sys = (Prod || Cons||E2)/{put, get}

con E2 = Empty1||Empty1

e Empty1 = put.Full1

Full1 = get.Empty1

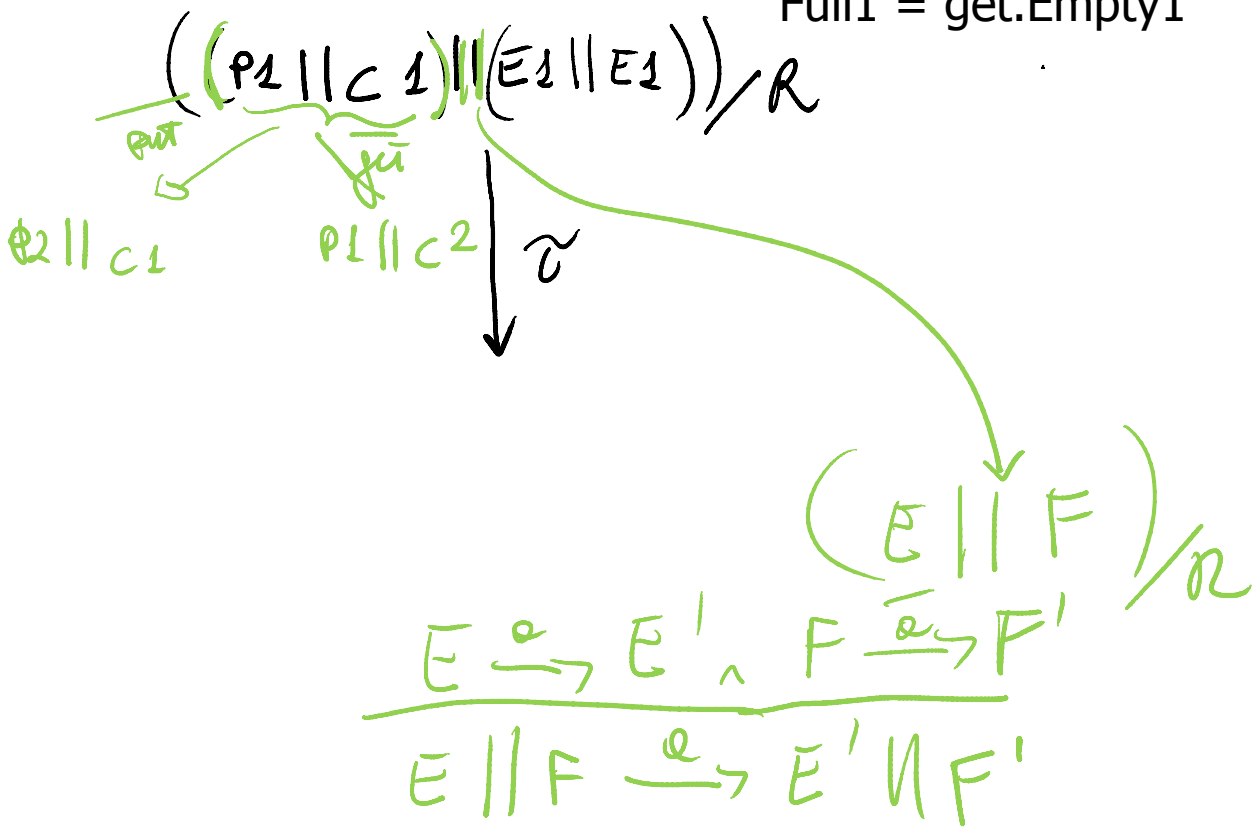
Buffer

Sys = (Prod || Cons||E2)/{put, get}

E2 = Empty1||Empty1

Empty1 = put.Full1

Full1 = get.Empty1



Nota: le annotazioni in verde descrivono le possibili evoluzioni dei termini componenti (P1||C1) da un lato e (E1||E1) dall'altro, sulla base dei quali si fa evolvere l'intero sistema (arco tau in nero)

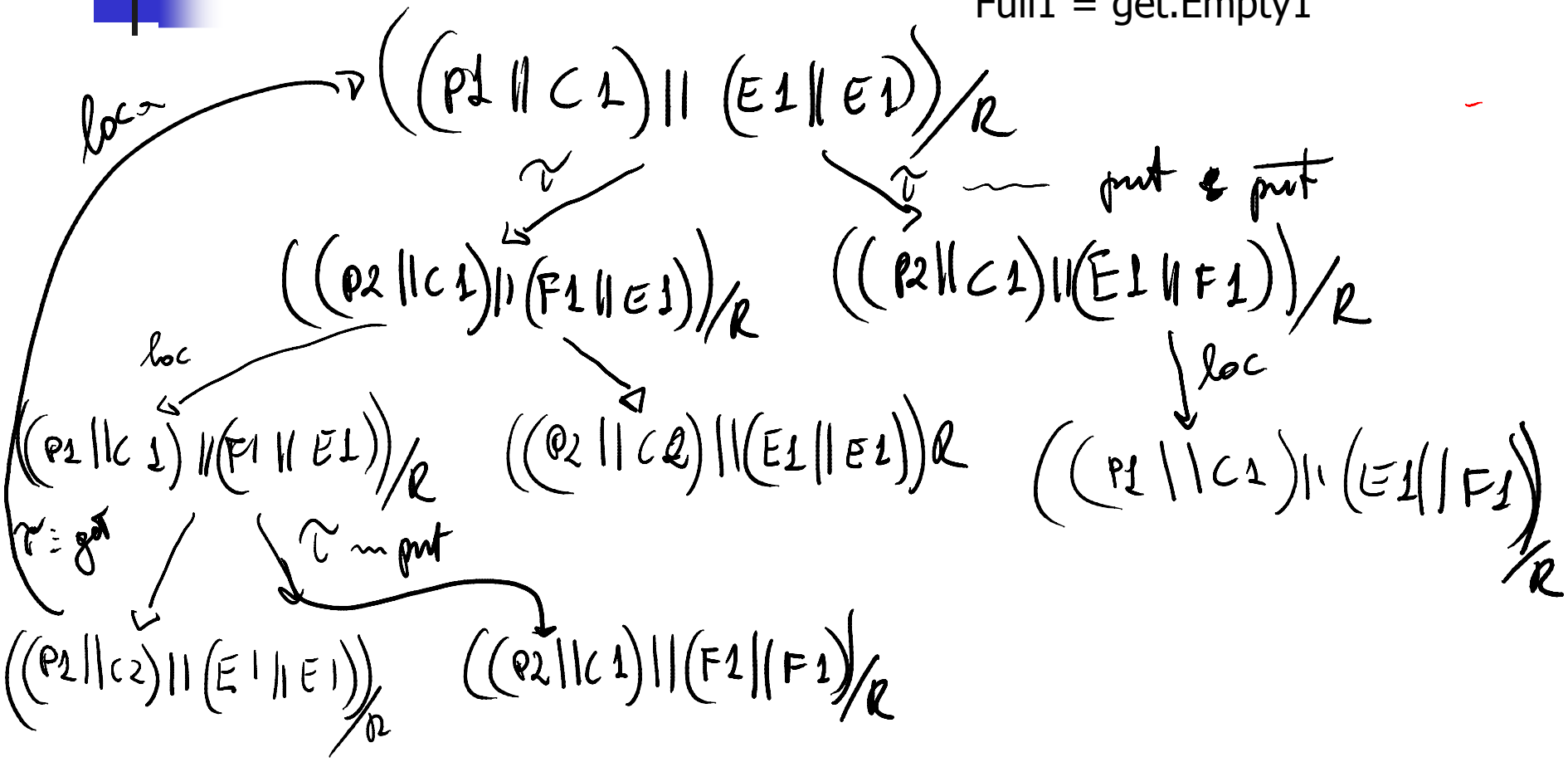
Buffer

Sys = (Prod || Cons||E2)/{put, get}

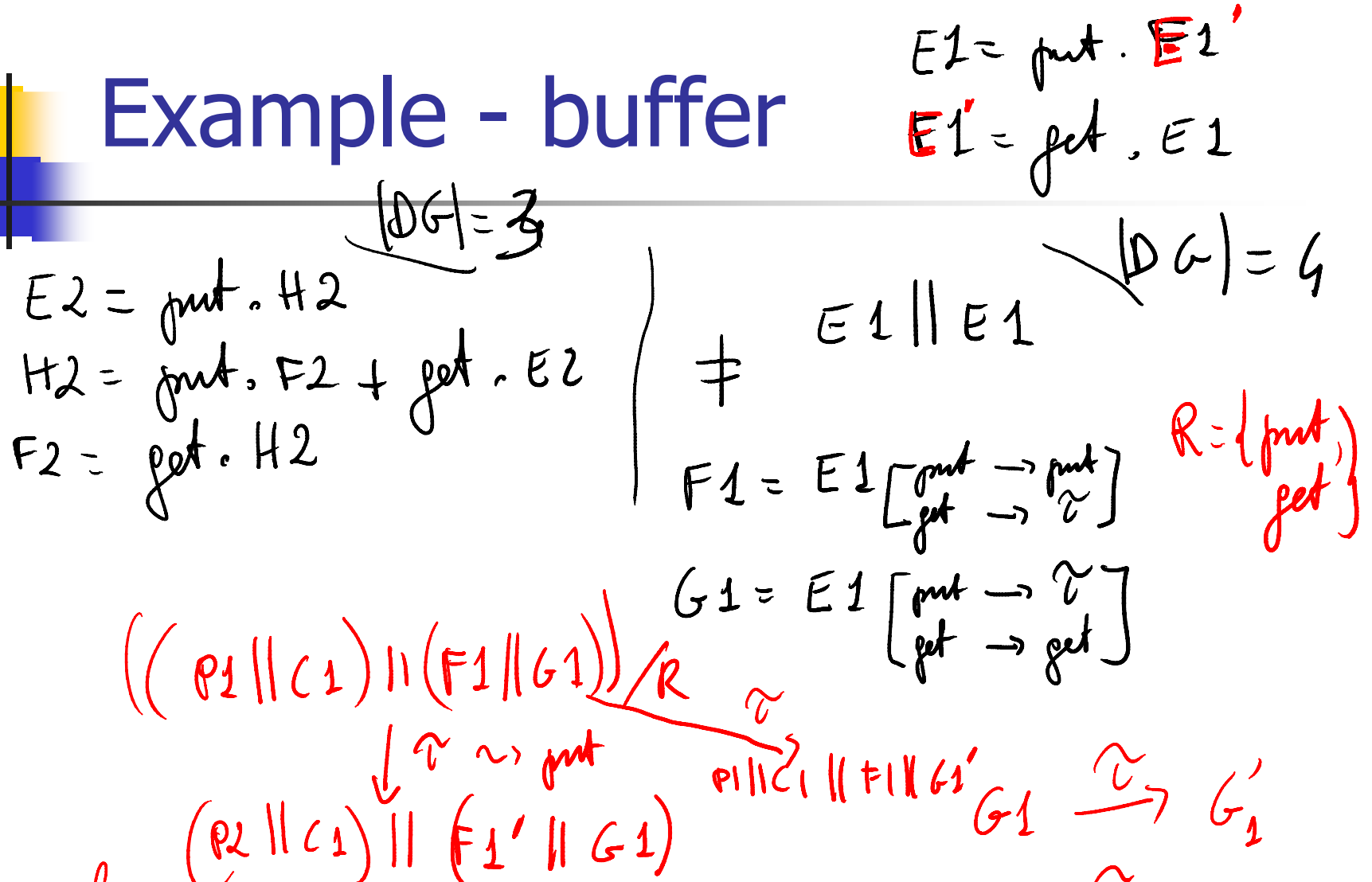
E2 = Empty1||Empty1

Empty1 = put.Full1

Full1 = get.Empty1

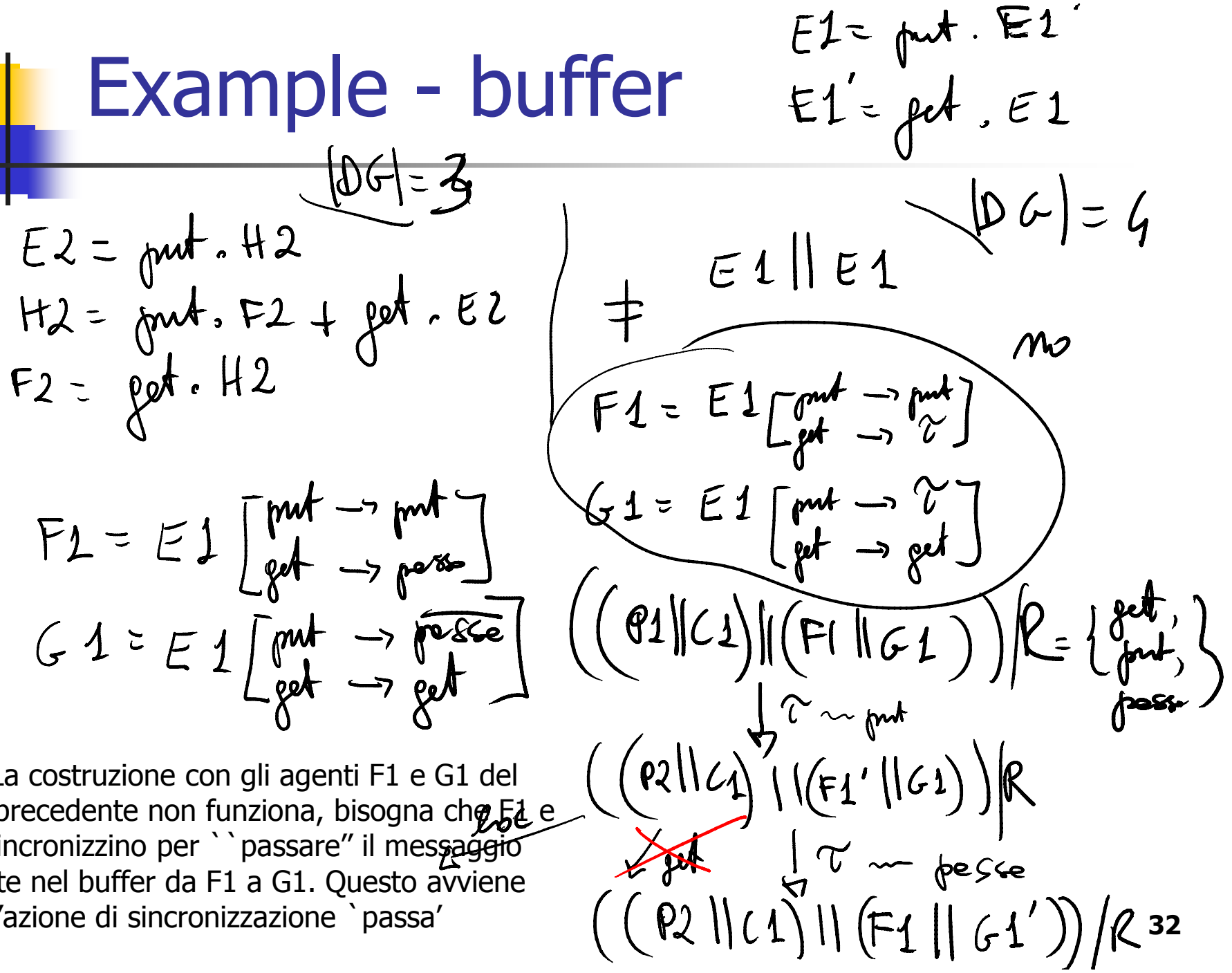


Example - buffer

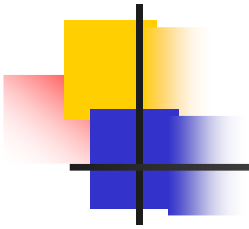


Nota: L'agente E2 e l'agente (E1||E1) hanno comportamenti simili rispetto agli agenti P1 e C1, ma il loro Derivation Graph hanno cardinalità diverse (3 e 4 rispettivamente) e quando il buffer è vuoto il primo agente offre una sola azione di put, il secondo due, quindi nella composizione con P1 e C1 avremo un DG più grande se usiamo E1||E1. F1 e G1 sono un tentativo di offrire una sola put e una sola get

Example - buffer



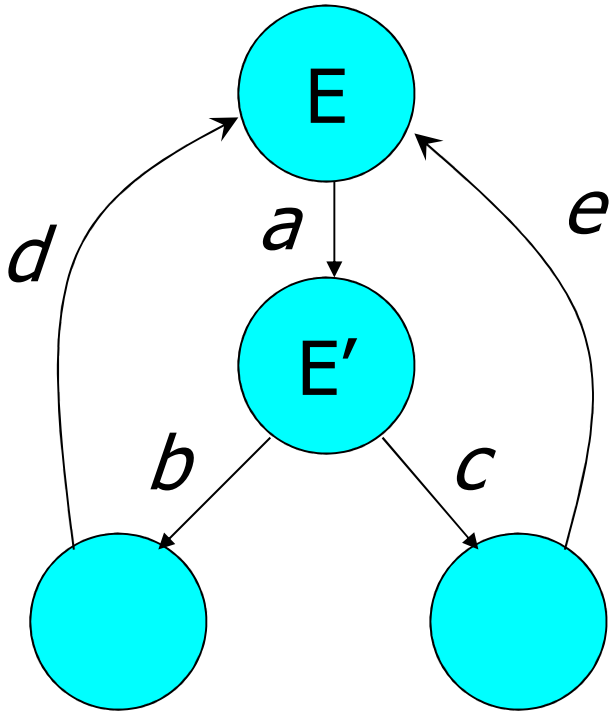
Nota: La costruzione con gli agenti F1 e G1 del lucido precedente non funziona, bisogna che F1 e G1 si sincronizzino per ``passare'' il messaggio presente nel buffer da F1 a G1. Questo avviene con un'azione di sincronizzazione `passa'

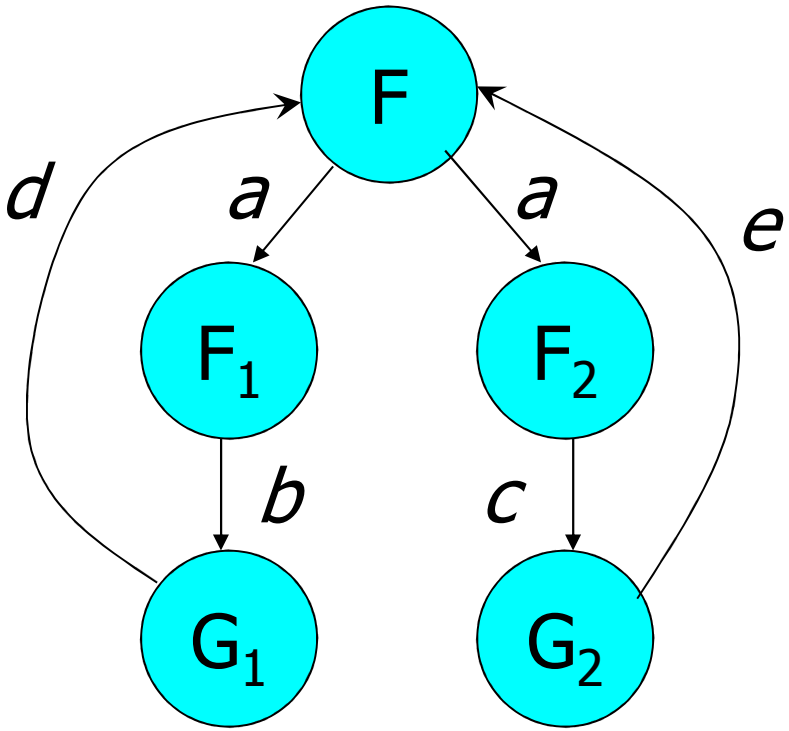


Osservabilità e sincronizzazione: different models may have the same set of executions!

Cosa vuol dire confrontare modelli? In algebra dei processi confrontiamo non solo le sequenze di azioni (quelle che sono le firing sequences in RdP), ma anche le interazioni fra agenti

Actions: Act={*a*,*b*,*c*,*d*}∪{τ}.
Agents: E, E′, F, F1, F2, G1, G2, …

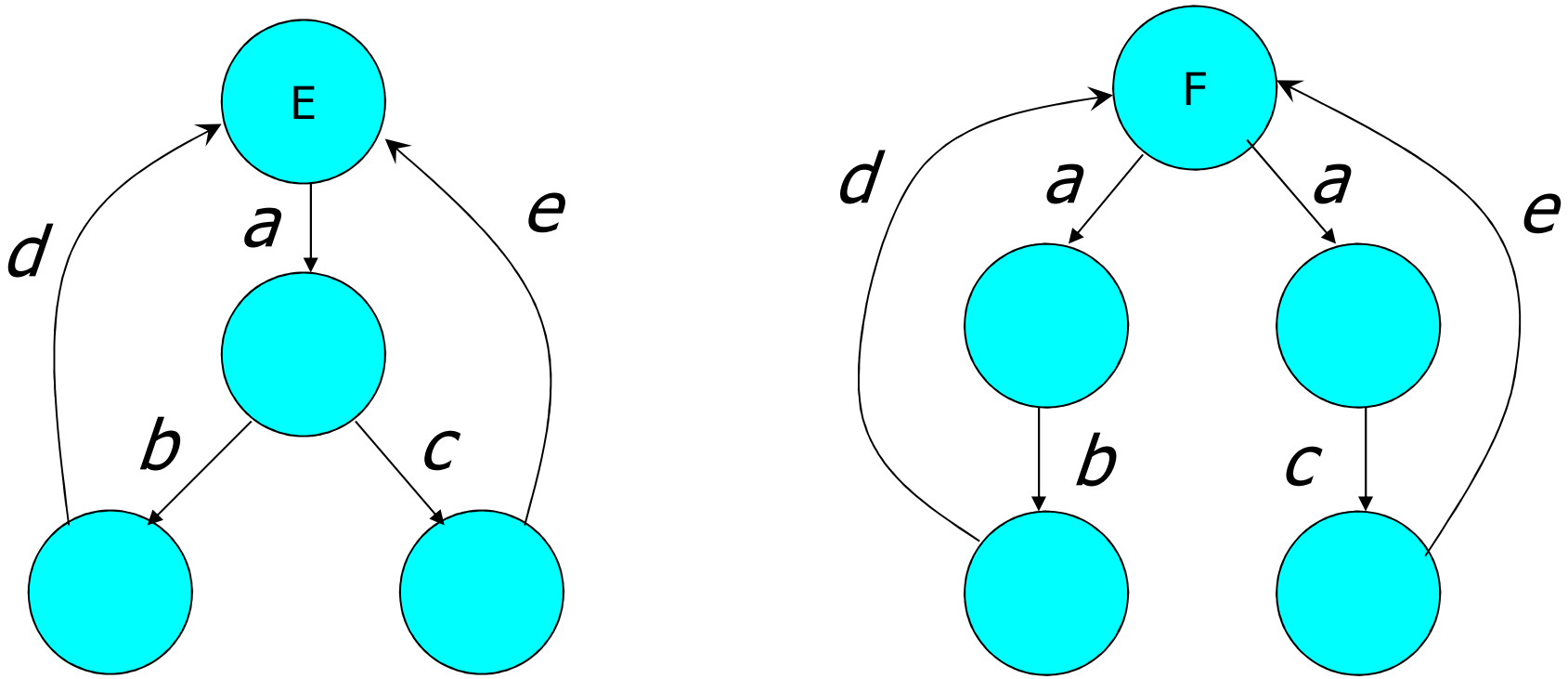




Agent E may *evolve* into agent E′.

Agent F may evolve into F_1 or F_2.

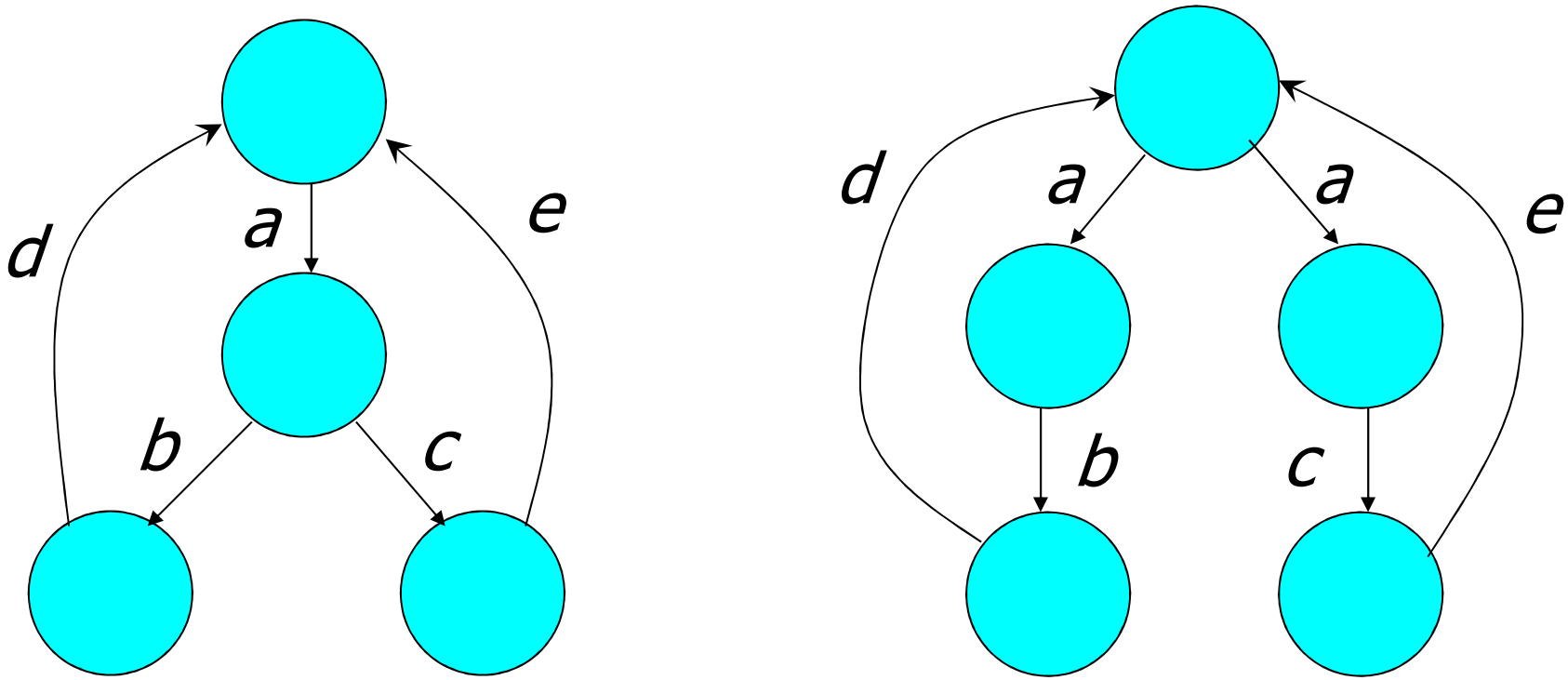
Osservabilità e sincronizzazione: different models may have the same set of executions!



a-insert coin, *b*-press coffee, *c*-press milked coffee

d-obtain coffee, *e*-obtain milked coffee

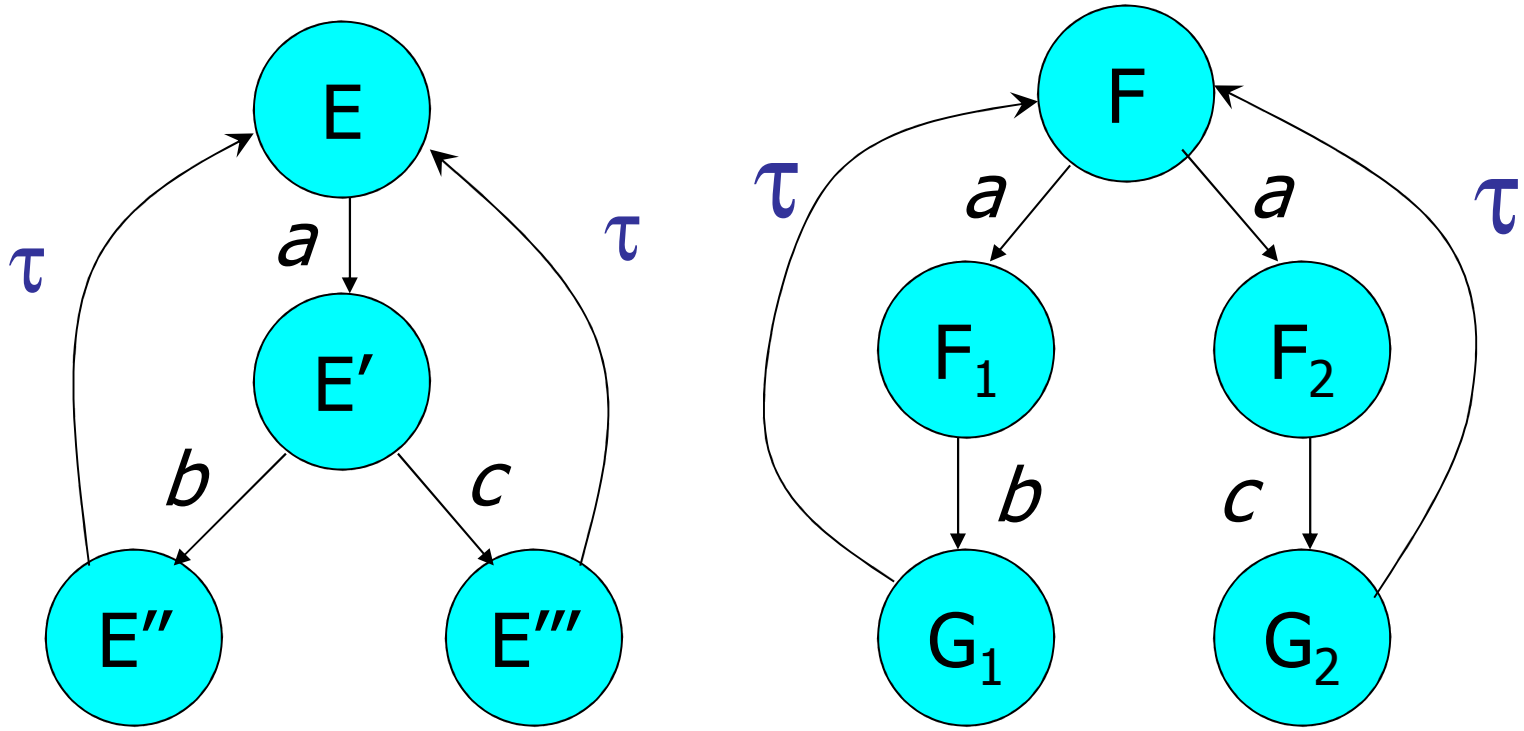
Different models may have the same set of executions, but different interactions!



a-insert coin, *b*-press coffee, *c*-press milked coffee, *d*-obtain coffee, *e*-obtain milked coffee:

it is the same to get a coffee from E or from F? Which one do you prefer?

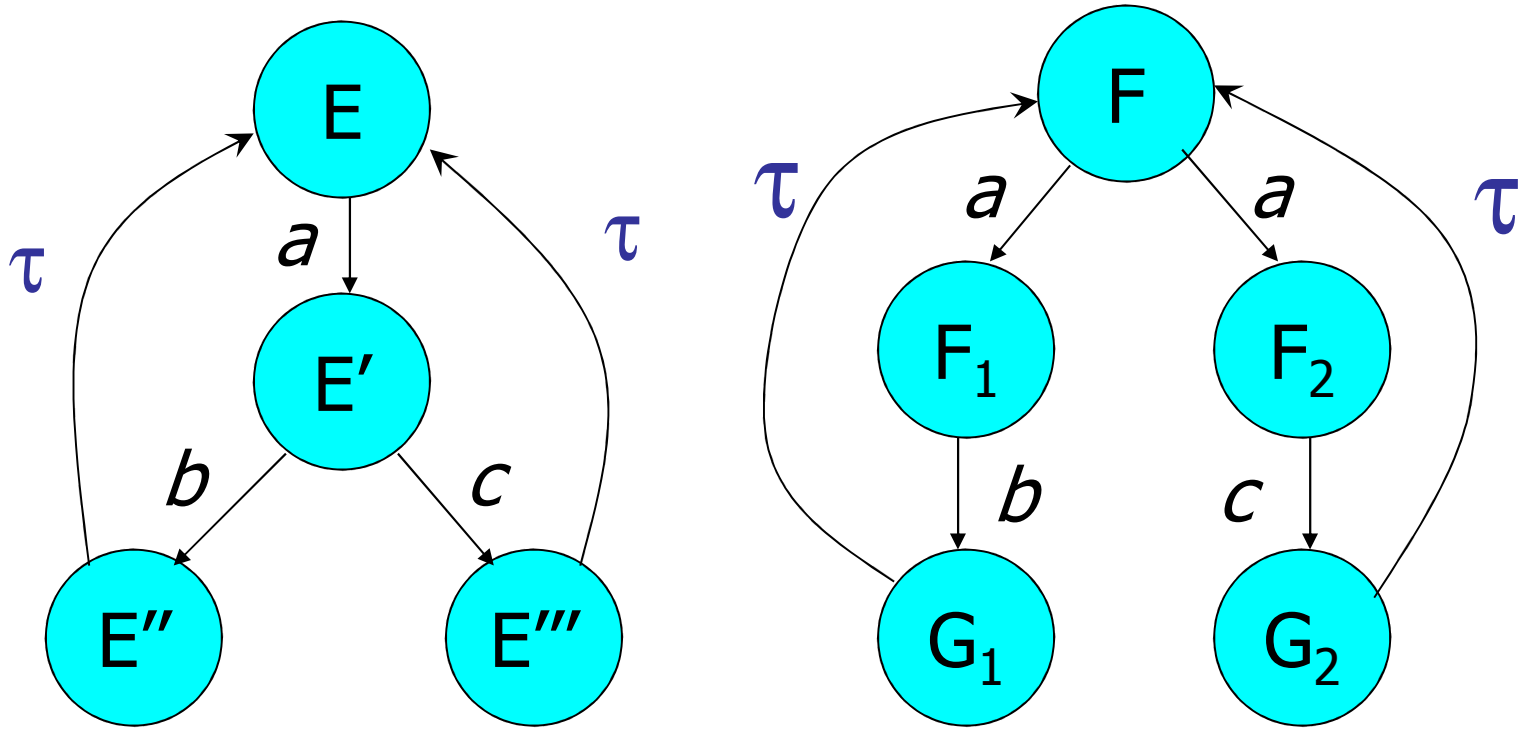
Events.



a-insert coin, *b*-get blue stamp, *c*-get red-stamp

… from a user point of view they might be equivalent

Events.

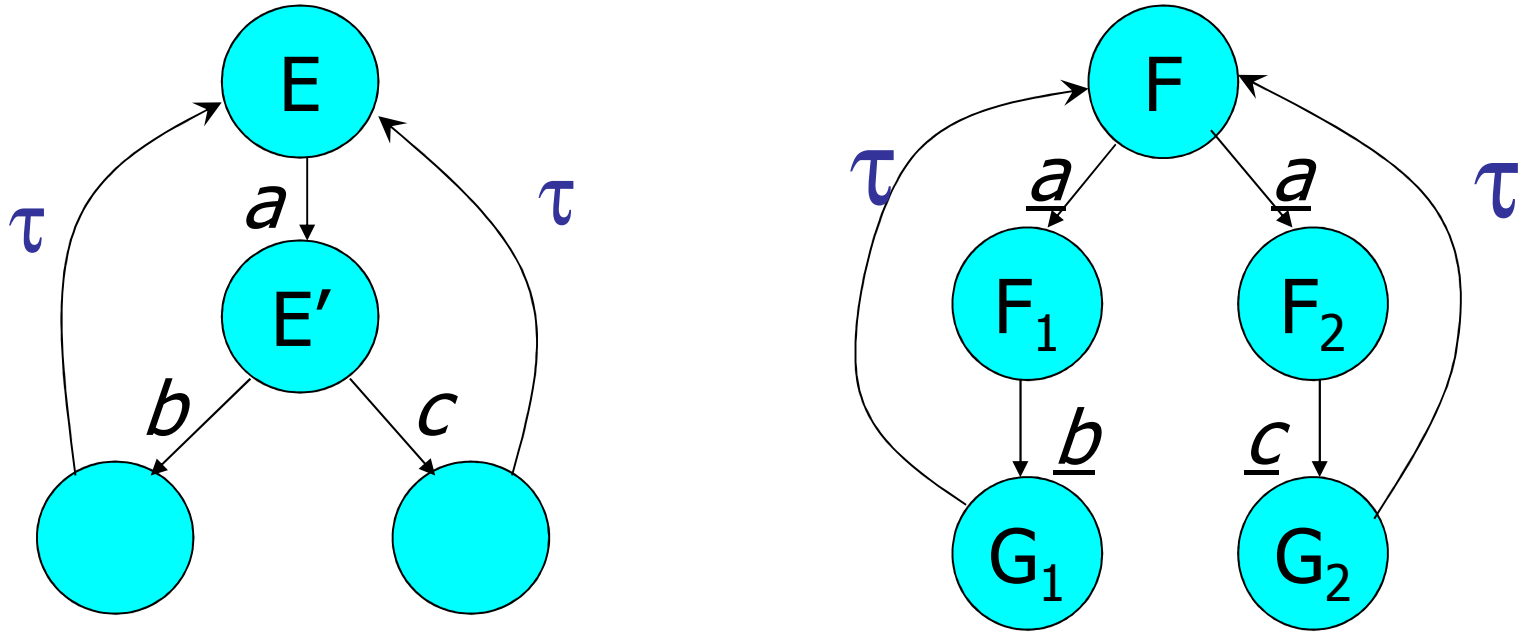


Agent evolutions:

F—*a*→F_1, F—*a*→F_2, F1—*b*→G_1, F_2—*c*→G_2, G_1—τ→F, G_2—τ→F.

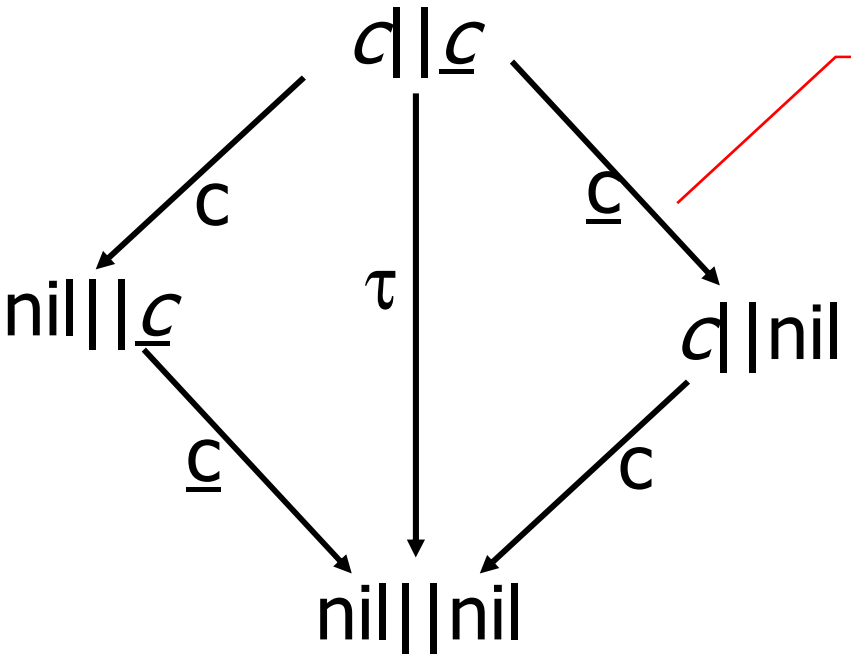
E—*a*→E′, ……

Actions and co-actions - interacting agents



If we substitute all actions of F of slide 38 with the corresponding co-actions and consider Sys = E||F , what is the behaviour of Sys?

Concurrent Composition

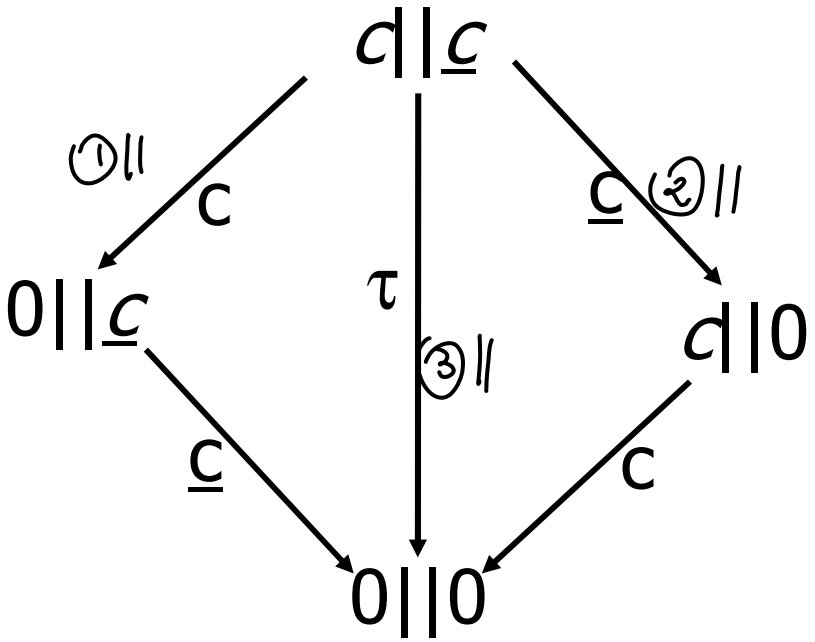


Derivation graph of an agent, D(c||c̲): (V,A), where V is the sets of all derivatives that can be derived from the agent c||c, through one or more application of the axioms or of the proof rules, and A is the set of arcs that represents the single application of a axiom or of a proof rule

Exercise: try to reproduce the same behavior with two nets, one can execute c and the other c̲.

Exercise: derivation graph of (c||c) and of (c||c̲)||c

Concurrent Composition



C||C

Derivation graph of (c||c)
and of (c||c)||c

$$\frac{E \xrightarrow{a} E'}{E \| F \xrightarrow{a} E' \| F} \qquad \frac{F \xrightarrow{a} F'}{E \| F \xrightarrow{a} E \| F'}$$

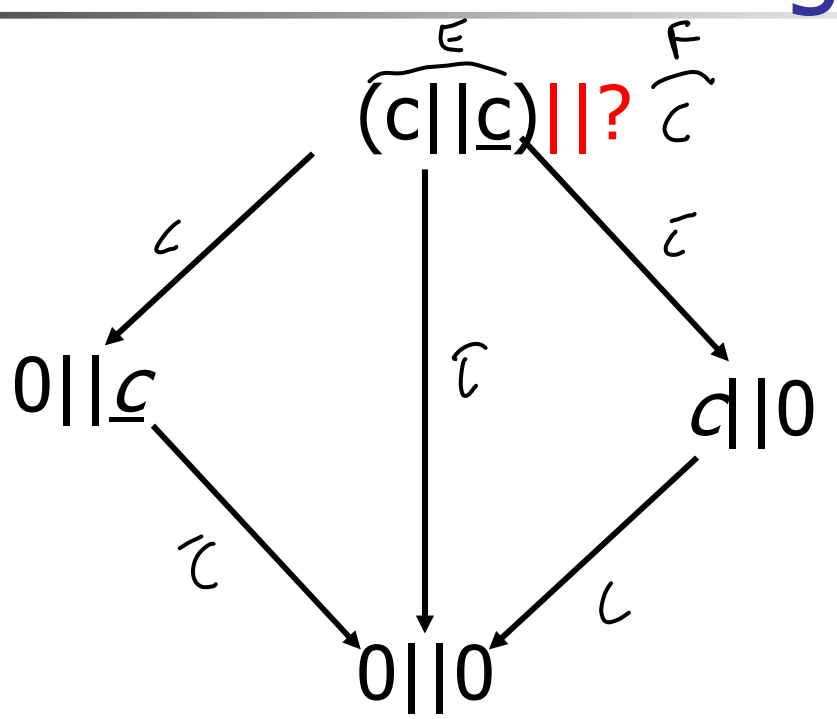
$$\frac{E \xrightarrow{a} E',\ F \xrightarrow{\underline{a}} F'}{E \| F \xrightarrow{\tau} E' \| F'}$$

C||C

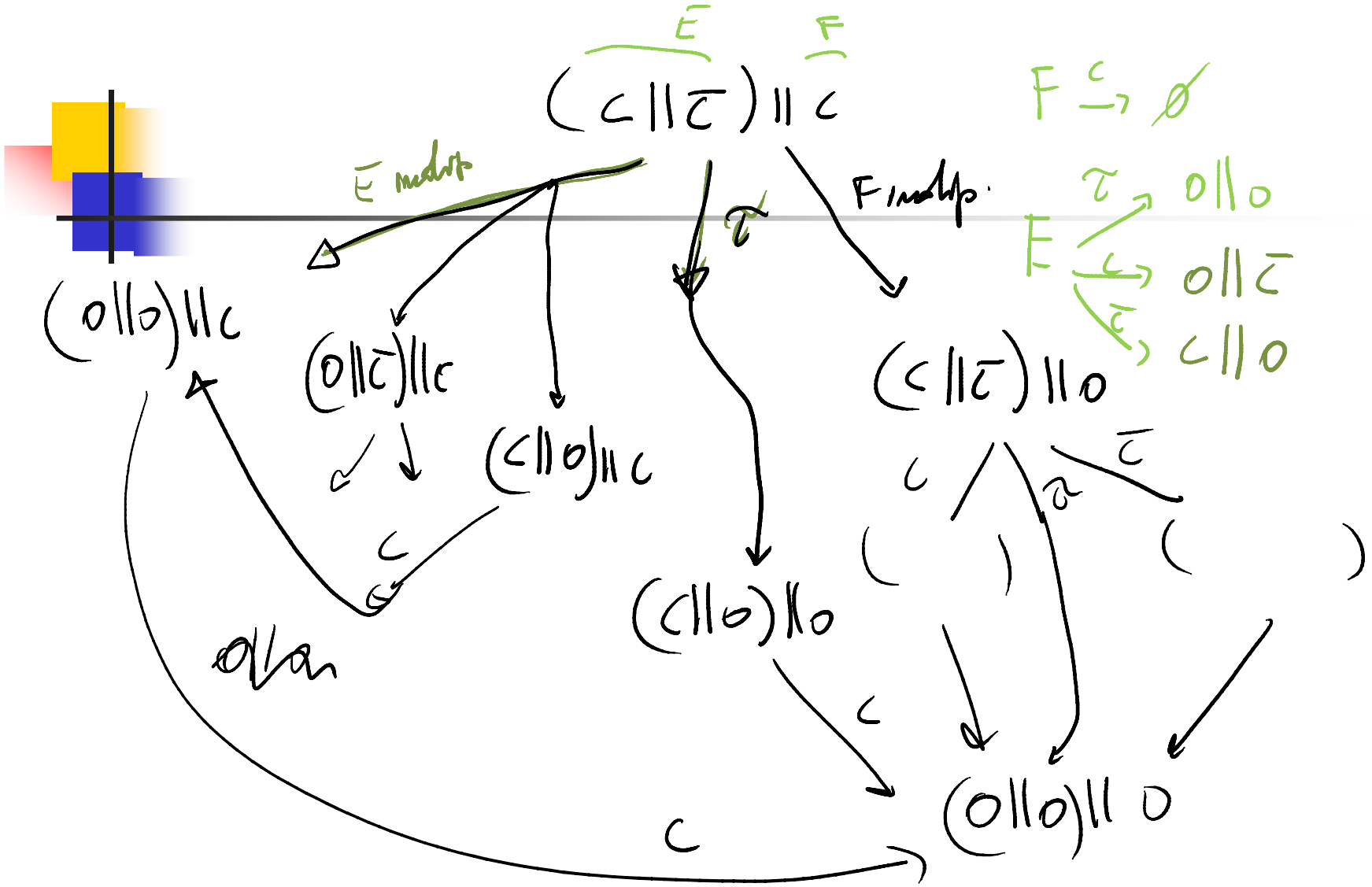
Concurrent Composition

Derivation graph of (c||c)||c

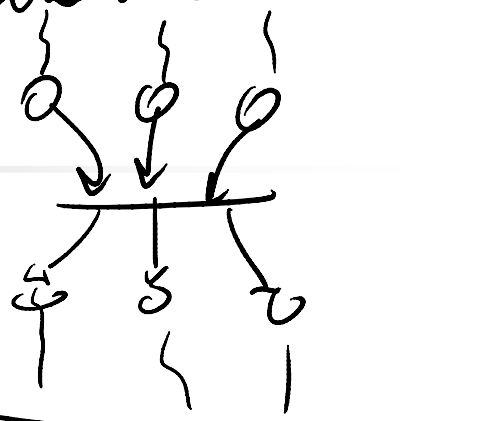
Synchronization among more than two agents



Try to have more than two agents (for example three) that synchronize on the sam action. Can I get this behaviour using CCS?


E
F
(c||c̄)||c
E matches
F matches
τ
F c↪∅
F τ↪o||o
F c↪o||c̄
F c̄↪c||o
(o||o)||c
(o||c̄)||c
(c||o)||c
(c||c̄)||o
(c||o)||o
o||o
(o||o)||o

Obiettivo : sincronizzazione simultanea di 3 processi



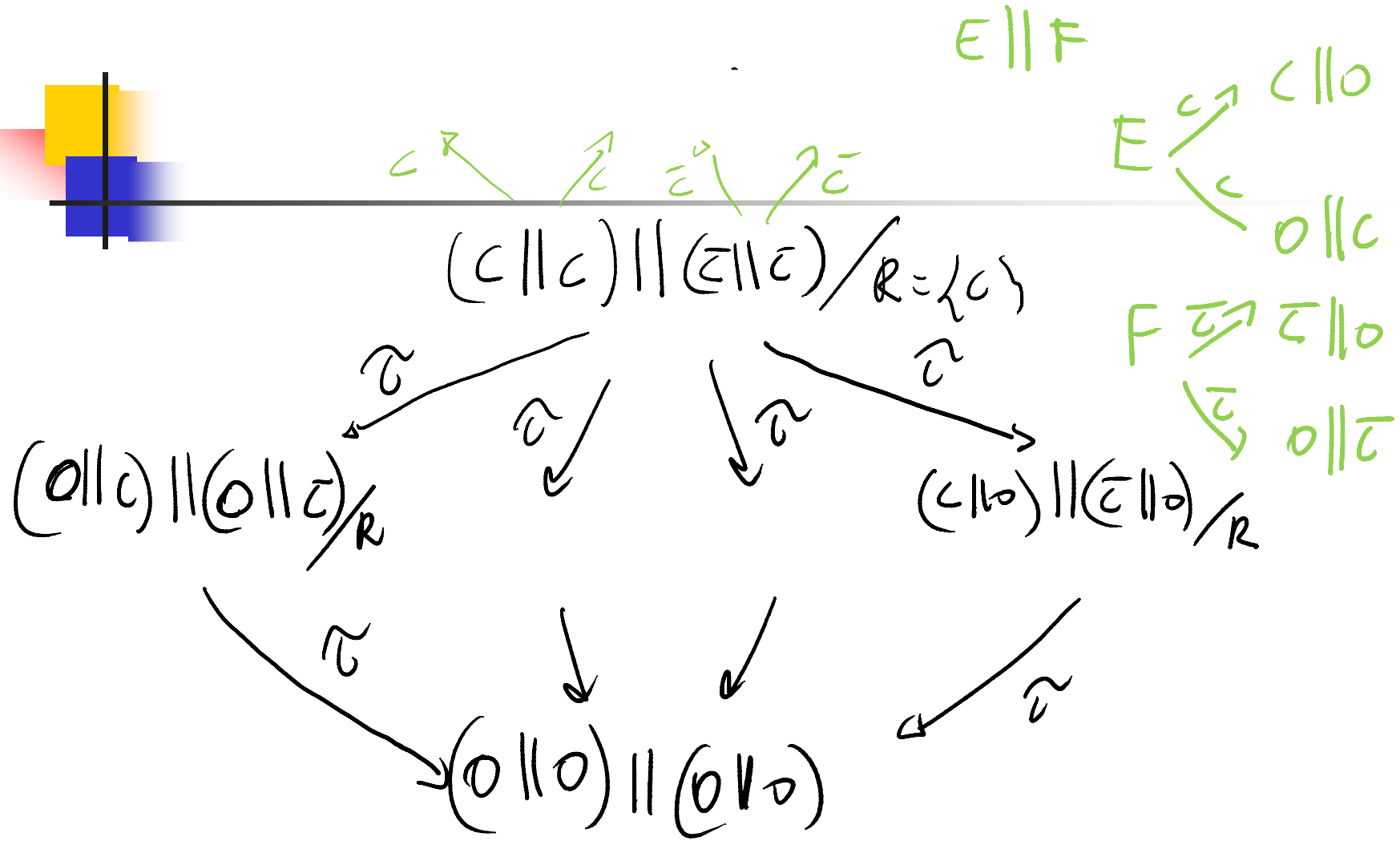
$$((c \| \bar{c}) \| c)\,/R=\{c\}$$

τ τ

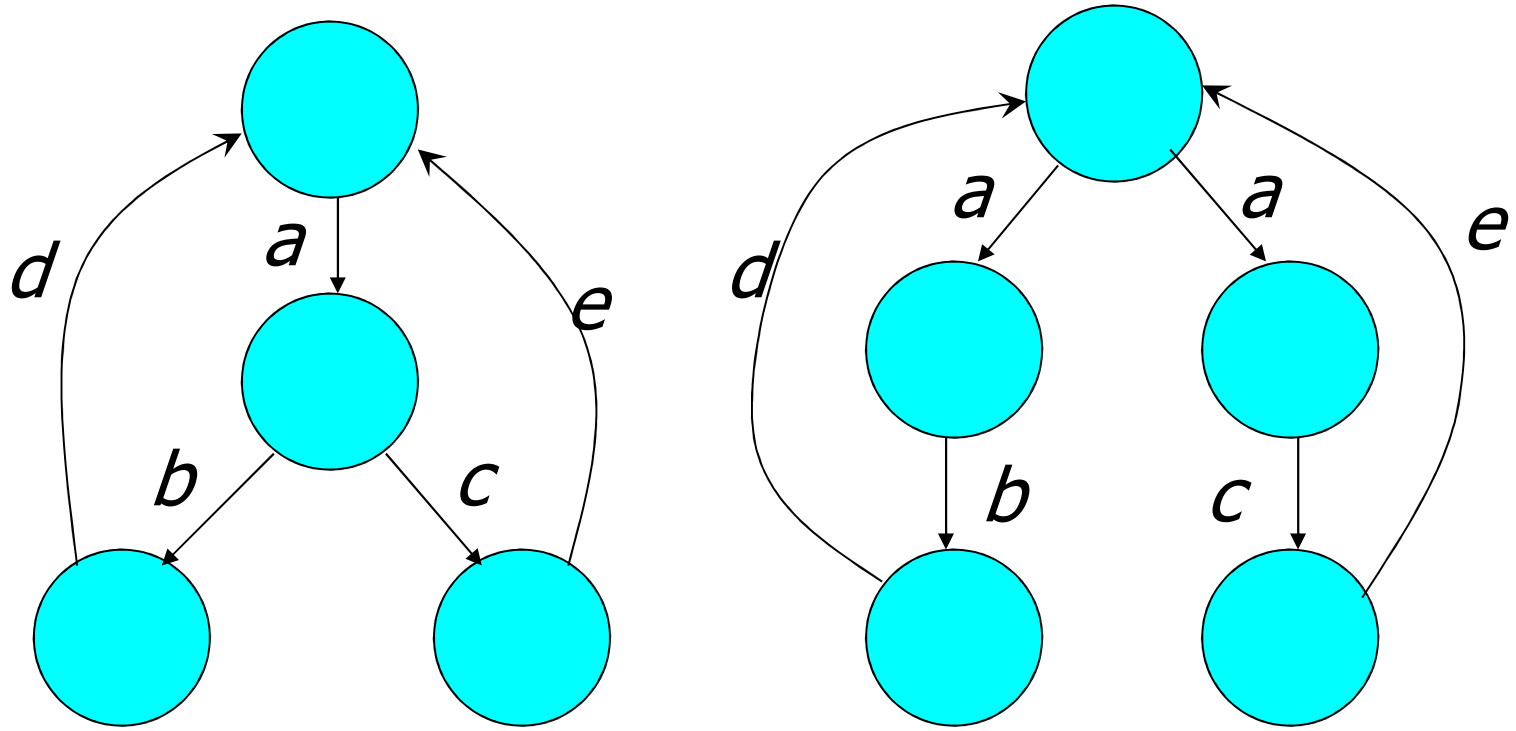
$$((o \| o) \| c)\,/R \qquad ((c \| o) \| o)\,/R$$

$$((c \| c) \| \bar{c})\,/R=\{c\}$$

$$(c \| o) \| o\,/R \qquad (o \| c) \| o\,/R$$

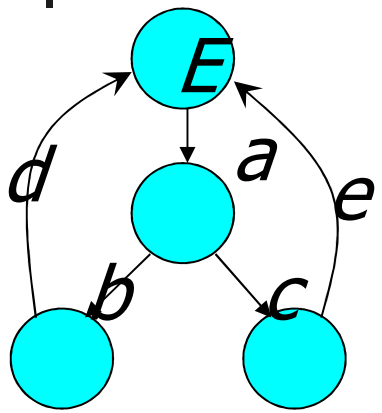

E || F
E
c
c || o
c
o || c
F
c̄
c̄ || o
c̄
o || c̄
(c || c) || (c̄ || c) / R = {c}
τ
τ
τ
τ
(o || c) || (o || c̄) / R
(c || o) || (c̄ || o) / R
τ
τ
τ
(o || o) || (o || o)

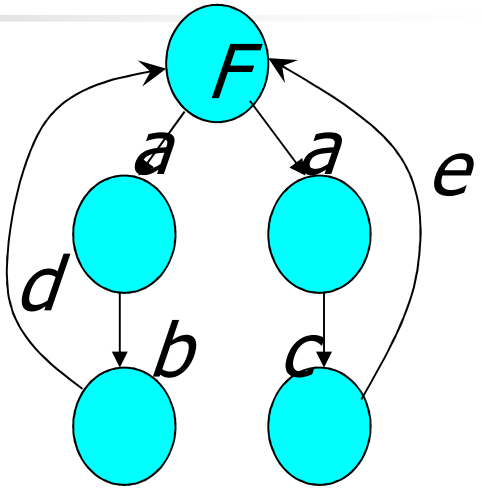
Equivalent Petri Nets models



Can we consider these as reachability graphs of two different nets? Labels are transitions names? Or transition labels?

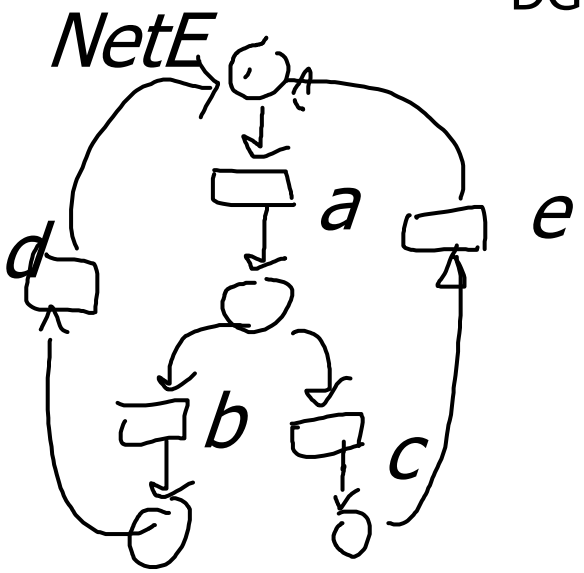
Equivalent Petri Nets models

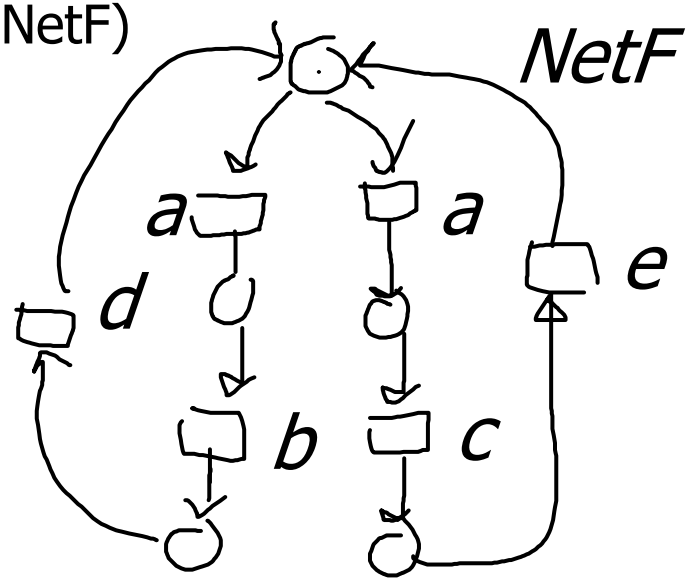




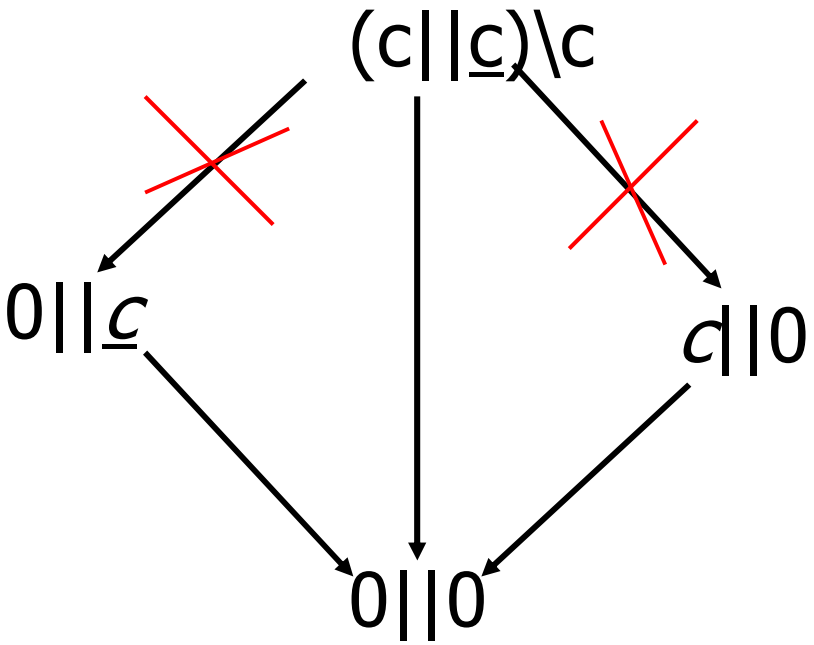
DG(E) isomorphic to RG(NetE)

DG(F) isomorphic to RG(NetF)





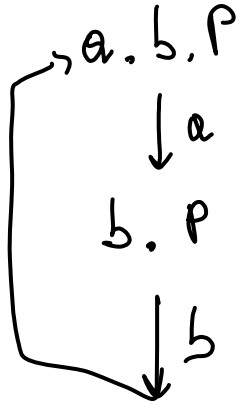
Consequence of restriction



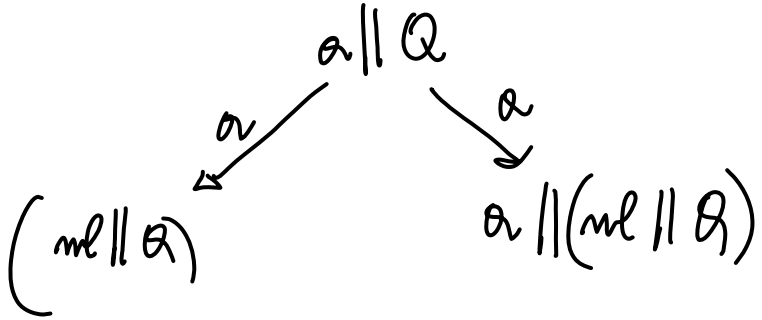
Try to reproduce the same behavior with two nets, one can execute c and the other c̲.

$PN_1 \|_c PN_2$

$P \triangleq a.(b.P)$



$Q \triangleq a \,||\, Q$



$P \triangleq$ soldi . (caffè + macchiato) . τ . P } syntactically not a CCS term

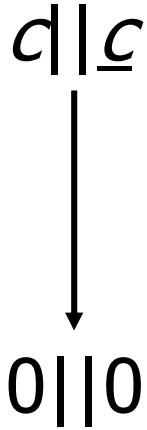
$P \triangleq$ soldi . P_1

$P_1 \triangleq$ caffè . P_2 + macchiato . P_2

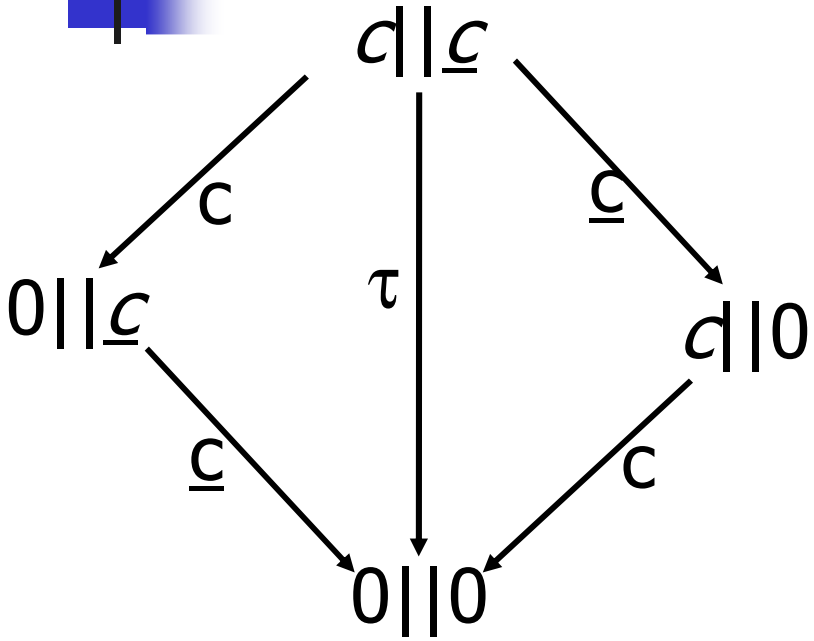
$P_2 \triangleq \tau . P$

Process algebras and Petri Nets

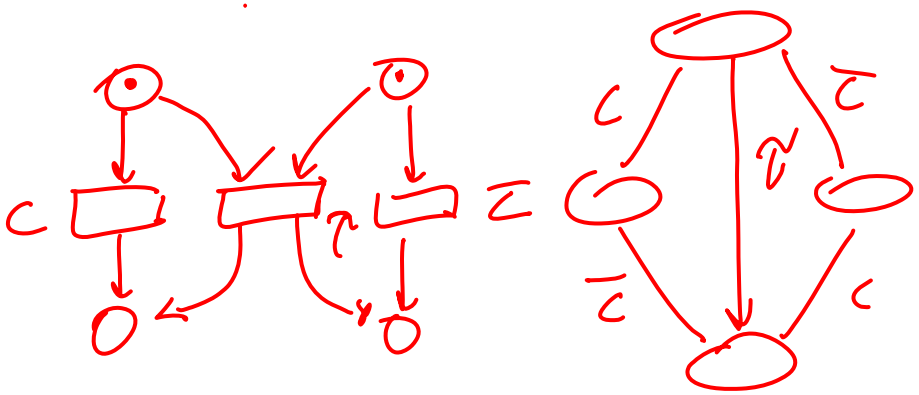
Given two nets N1 and N2 of similar behaviour (one transition each, labelled c and $\underline{c}$ respectively) produce two process algebra terms such that the RG and the derivation graph are isomorphic and equal to the below



Concurrent Composition

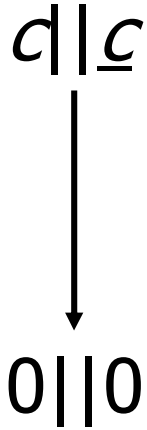


reproduce the same behavior with two nets, one can execute c and the other c̄.



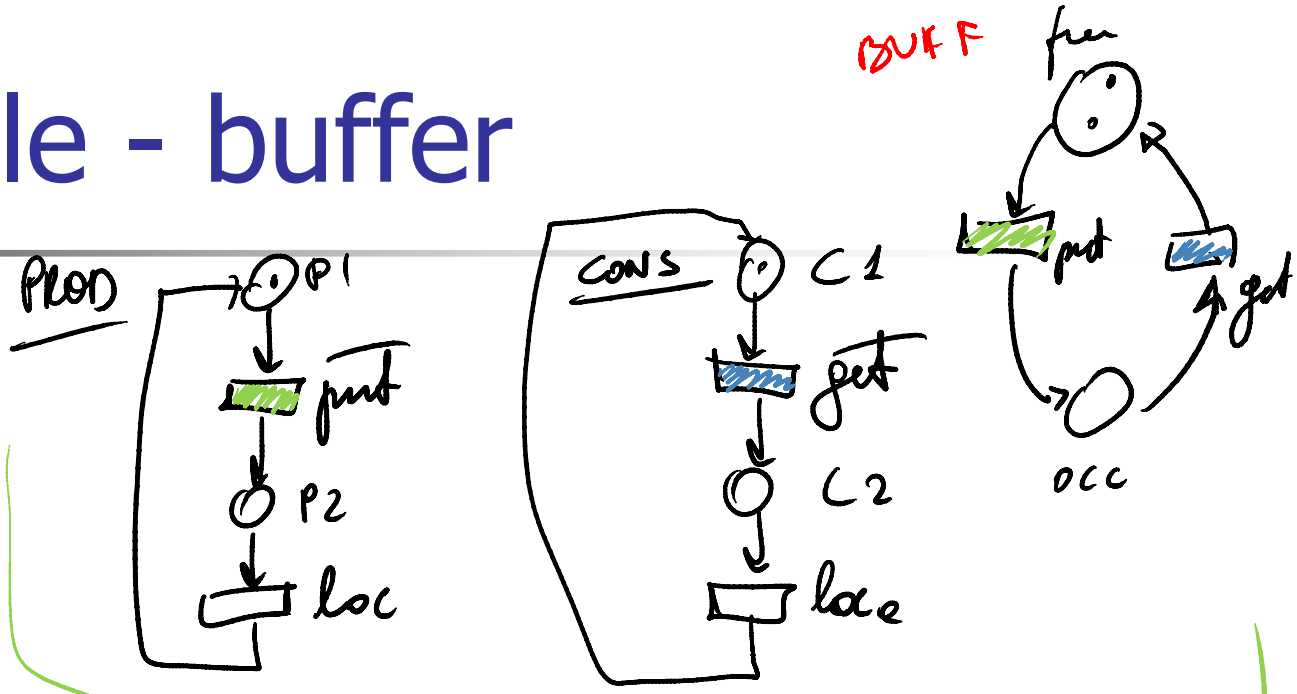
Concurrent Composition

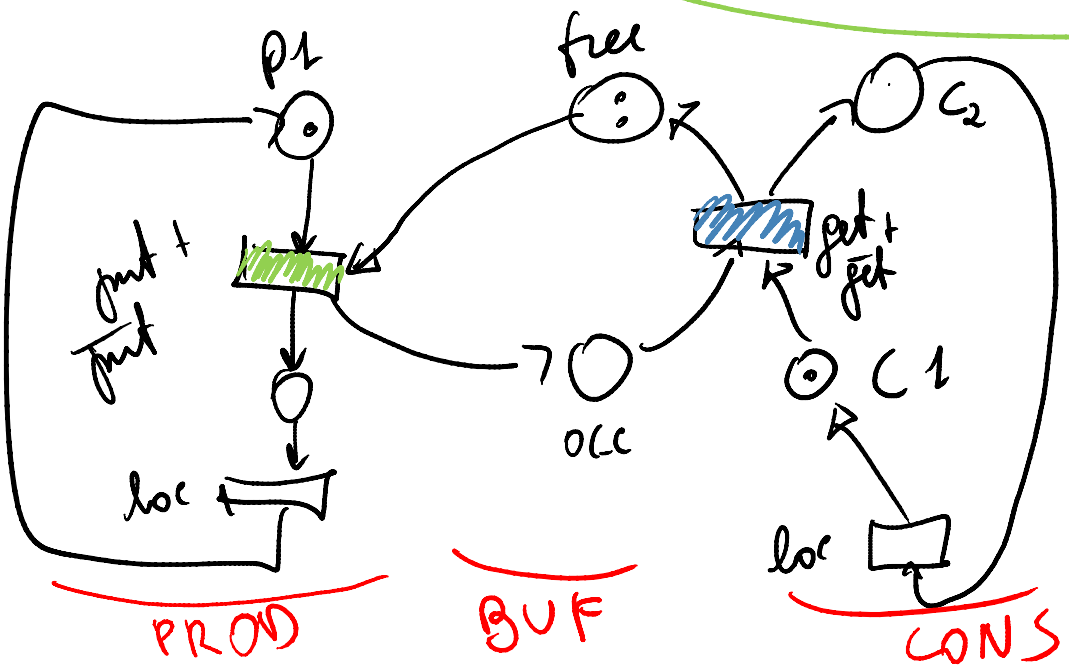
Given two nets N1 and N2 of similar behaviour (one transition each, labelled c and $\underline{c}$ respectively) produce two process algebra terms such that the RG and the derivation graph are isomorphic and equal to the below



Example - buffer

Using Petri Nets.....





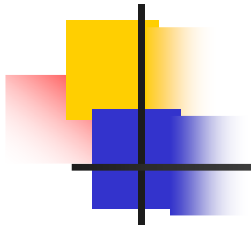
Number of positions in buffer depends on initial marking

Modellare variabili booleane

Per meglio comprendere la modellazione delle variabili booleane/binarie in process algebra consideriamo prima la loro modellazione con le reti P/T (con archi inibitori)

Operazioni previste per una variabile booleana *v*:

1. Set *v* a true
2. Set *v* a false
3. Is *v* true?
4. Is *v* false?
5. not *v* (setta la variabile al valore opposto)



Modellare variabili booleane

Due possibili scelte per l'implementazione di *v* :

1) la variabile *v* è rappresentata da un posto di nome *pv* e, in una marcatura *m*, *m*[*pv*] = 0 significa che v è falsa, mentre *m*[*pv*] = 1 significa che v è vera.

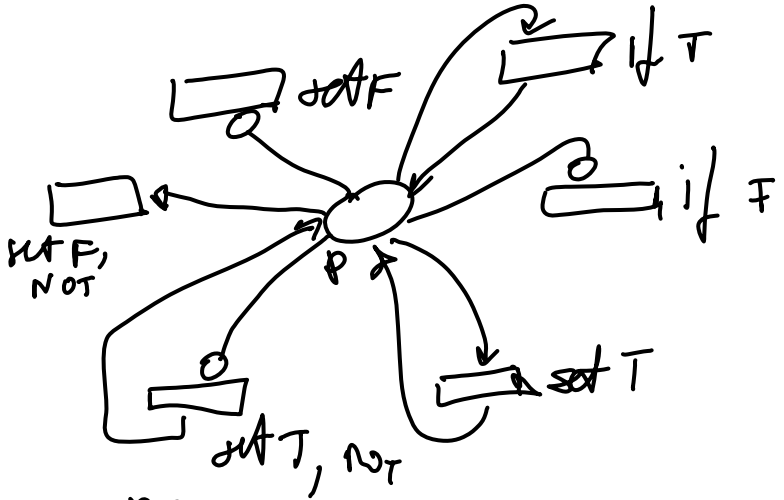
2) la variabile *v* è rappresentata da due posto di nome *pv* e, in una marcatura *m*, *m*[*pv*] = 0 significa che v è falsa, mentre *m*[*pv*] = 1 significa che v è vera.

Modellare variabili binarie con Rete P

{T, F} ← valori

set a T , if T , NOT ⇝ cambia il valore
set a F , if F

variabile è 1 solo posto

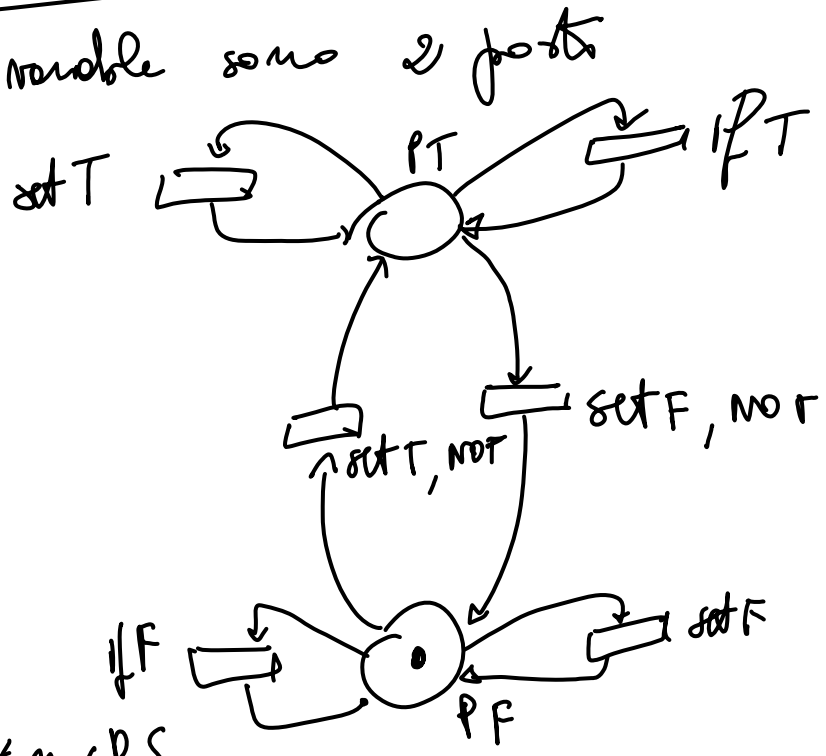


$\forall m \in RS$

$m(P) = 0 \Rightarrow F$, $m(P) = 1 \Rightarrow T$

$m(P) \leq 1$

variabile sono 2 posti

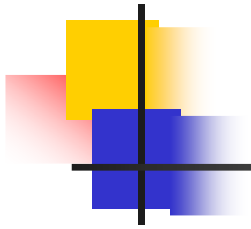


$\forall m \in RS$

$m(P_T) = 1 \Rightarrow T$, $m(P_F) = 1 \Rightarrow F$

$m(P_T) + m(P_F) = 1$

Modeling binary variable in CCS



Esempio tratto dal testo del Peled:

Variabile C binaria (valori 0 e 1). Modello con due agenti:

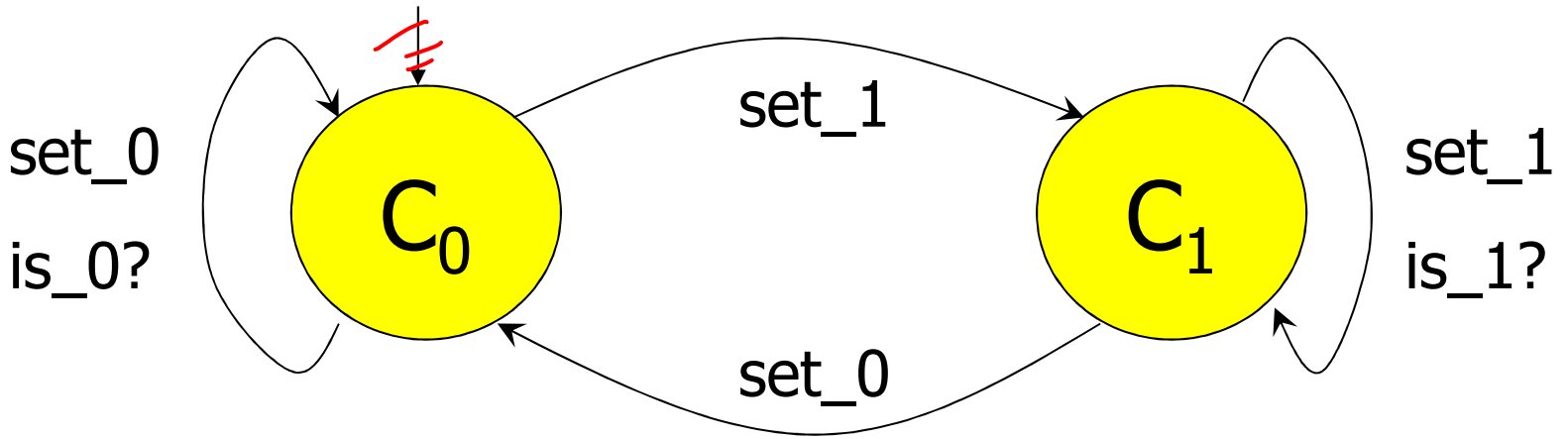
C_0 --> agente che rappresenta il comportamento di una variabile che vale zero

C_1 --> agente che rappresenta il comportamento di una variabile che vale uno

Quale agente offre quali azioni? Qui non c'è il not perché si tratta di una variabile a due valori (non necessariamente booleana)

Modeling binary variable

Act = { set_1, set_0, is_0?, is_1 }



$C_0 = is\_0? \,.\, C_0 + set\_1 \,.\, C_1 + set\_0 \,.\, C_0$

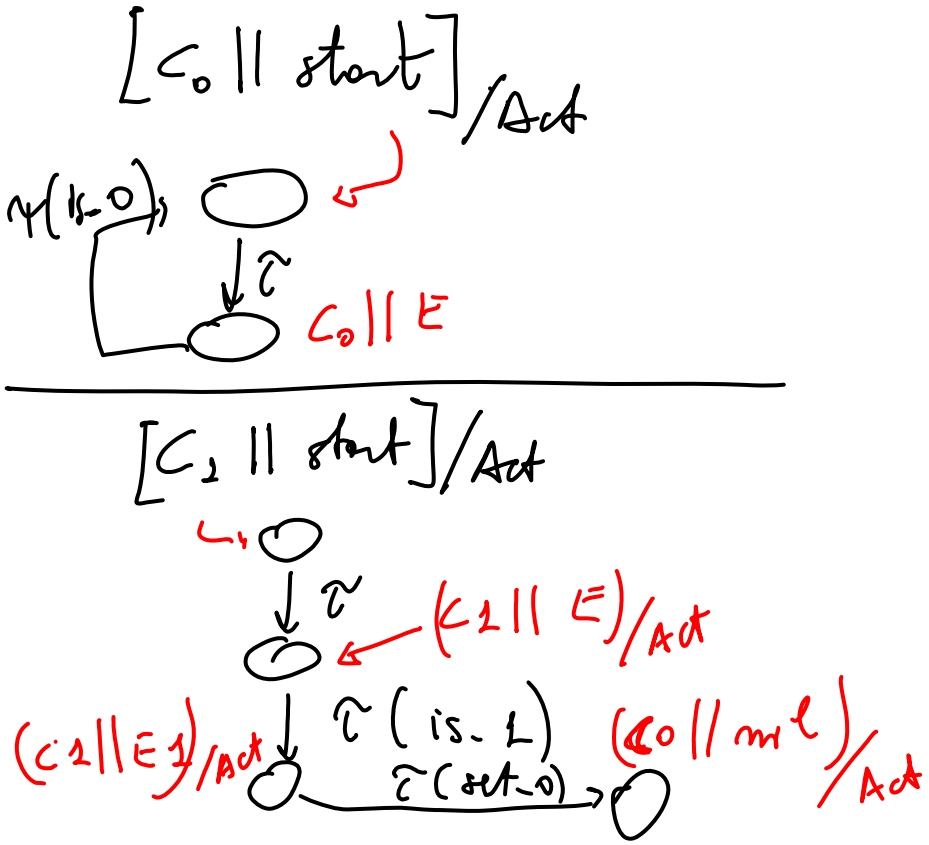
$C_1 = is\_1? \,.\, C_1 + set\_0 \,.\, C_0 + set\_1 \,.\, C_1$

$Act = \{is\_0? , is\_1?, set\_1, set\_0\}$

Modeling binary variable

$C_0 = is\_0? \,.\, C_0 + set\_1 \,.\, C_1 + set\_0 \,.\, C_0$

$C_1 = is\_1? \,.\, C_1 + set\_0 \,.\, C_0 + set\_1 \,.\, C_1$



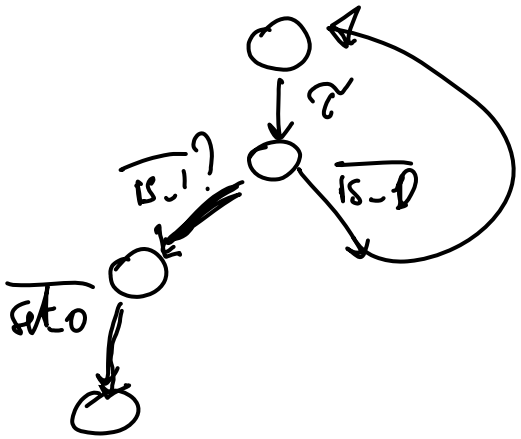
start := τ

if C == 1 then C := 0 else goto start

START ≜ ~~set_0.E~~ τ . E

E ≜ $\overline{is\_1}?$. E1 + $\overline{is\_0}?$. start

E1 ≜ $\overline{set\_0}$. nil



Dekker's algorithm

boolean c1 initially 1 – P1 does not want to enter;

boolean c2 initially 1 --

integer (1..2) turn initially 1 (ro 2);

```
P1::while true do
      begin
        non-critical section 1
        c1:=0;
        while c2=0 do
          begin
            if turn=2 then
              begin
                c1:=1;
                wait until turn=1;
              end
          end
        critical section 1
        c1:=1;
        turn:=2
      end.
```

```
P2::while true do
      begin
        non-critical section 2
        c2:=0;
        while c1=0 do
          begin
            if turn=1 then
              begin
                c2:=1;
                wait until turn=2;
              end
          end
        critical section 2
        c2:=1;
        turn:=1
      end.
```

Dekker's algorithm

```
boolean c1 initially 1;
boolean c2 initially 1;
integer (1..2) turn initially 1;
```

```
P1- p::while true do
      begin
        non-critical section 1
        c1:=0; wantp:= true
        while c2=0 wantq do
          begin
            if turn=2 then
              begin
                c1:=1; wantp:= false
                wait until turn=1;
                wantp:= true
              end
          end
        critical section 1
        c1:=1; turn:=2
        turn:=2; wantp:= false
      end.
```

Semantica variabili:

wantp == true → p vuole accedere alla critical section (CS)

wantq == true → q vuole accedere alla critical section (CS)

turn: chi ha il turno, e' una variabile che serve a redimere i casi in cui ambedue i processi facciano una richiesta (cioe' quando sia wantp che wantq sono a true). turn==1 → processo p

turn==2 → processo q

Logica dell'algoritmo: quando ambedue i processi vogliono accedere (wantp e wantq a true) entrano ambedue nel while, ma quello che ha turn in suo favore (supponiamo sia q) "passa" per primo, mentre l'altro aspetta sulla wait. Per permettere a q di uscire dal while p deve mettere a false la sua richiesta di accesso (wantp a false appena prima della wait). Quando poi turn volge a suo favore si sblocca la wait e p rimette wantp a true, in modo da poter eventualmente accedere alla CS (se nel frattempo wantq e' andato a false) avendo la variabile wantp correttamente settata a true (senno' q potrebbe, mentre p e' in regione critica, ritornare ad eseguire le istruzioni non CS e poi settare wantq a zero e, trovando wantp false, uscire immediatamente dal while e accedere a sua volta alla CS, che e' il problema che abbiamo riscontrato in classe).

Non penso che l'ordine delle ultime due istruzioni sia rilevante, ma non ho controllato

Dekker's algorithm

```
boolean c1 initially 1;
boolean c2 initially 1;
integer (1..2) turn initially 1;
```

```
P1- p::while true do
      begin
        non-critical section 1
        c1:=0; wantp:= true
        while c2=0 wantq do
          begin
            if turn=2 then
              begin
                c1:=1; wantp:= false
                wait until turn=1;
                wantp:= true
              end
          end
        critical section 1
        c1:=1; turn:=2
        turn:=2; wantp:= false
      end.
```

```
P2- q::while true do
      begin
        non-critical section 2
        c2:=0; wantq:= true
        while c1=0 wantp do
          begin
            if turn=1 then
              begin
                c2:=1; wantq:= false
                wait until turn=2;
                wantq:= true
              end
          end
        critical section 2
        c2:=1; ; turn:=1
        turn:=1; wantq:= false
      end.
```

Dekker's algorithm

```
P1::while true do
      begin
        non-critical section 1
        c1:=0;
        while c2=0 do
          begin
            if turn=2 then
              begin
                c1:=1;
                wait until turn=1;
              end
          end
        critical section 1
        c1:=1;
        turn:=2
      end.
```

Translate into process algebra and into Petri nets

Build the derivation graph of the process algebra term and the RG of the Petri net and compare

Modeling binary variable

Modellare la variabile turn e C2 partendo dagli agenti per C1

C_0=is_0? . C_0 + set_1 . C_1 + set_0 . C_0

C_1=is_1? . C_1 + set_0 . C_0 + set_1 . C_1

C1_0 = set_0_c1 . C1_0 + set_1_c1 . C1_1 + if_c1_0 . C1_0

C1_1 = set_0_c1 . C1_0 + set_1_c1 . C1_1 + if_c1_1 . C1_1

non usato da Dekker

C2_0 = C1_0 /{set_0_c1 → set_0_c2 ⋮ }

C2_1

relabelling

TURN_1 = C1_0 /{set_0_c1 → set_1_TURN ⋮ }

TURN_2 =

Dekker's algorithm

```
boolean c1 initially 1;
boolean c2 initially 1;
integer (1..2) turn initially 1;
```

```
P1::while true do
      begin
        non-critical section 1
        c1:=0;
        while c2=0 do
           begin
              if turn=2 then
                begin
                   c1:=1;
                   wait until turn=1;
                end
           end
        critical section 1
        c1:=1;
        turn:=2
      end.
```

```
P2::while true do
      begin
        non-critical section 2
        c2:=0;
        while c1=0 do
           begin
              if turn=1 then
                begin
                   c2:=1;
                   wait until turn=2;
                end
           end
        critical section 2
        c2:=1;
        turn:=1
      end.
```

Dekker's algorithm

boolean c1 initially 1;

boolean c2 initially 1;

integer (1..2) turn initially 1;

```
P1::while true do
      begin
p1.e  non-critical section 1
p1.b  c1:=0;
p1.c  while c2=0 do
        begin
p1.d      if turn=2 then
            begin
              c1:=1;
              wait until turn=1;
            end
        end
CS1   critical section 1
      c1:=1;
      turn:=2
    end. P1
```

CSP-like process algebra

CSP defined by Hoare in '83

Similar to CCS but no notion of action and co-action

$$P ::= Nil \mid a.P \mid P + P \mid P \|_S P \mid P/L \mid A$$

Parallel Composition

$$\frac{P \xrightarrow{a} P'}{P \|_S Q \xrightarrow{a} P' \|_S Q} \; a \notin S \qquad \frac{Q \xrightarrow{a} Q'}{P \|_S Q \xrightarrow{a} P \|_S Q'} \; a \notin S$$

$$\frac{P \xrightarrow{a} P' \;\; Q \xrightarrow{a} Q'}{P \|_S Q \xrightarrow{a} P' \|_S Q'} \; a \in S$$

CSP-like process algebra

Consequences of the new proof rule for the parallel composition :

$(E \|_c F) \|_c G$

$(c.0 \|_{\{c\}} c.0) \|_{\{c\}} c.0$

M $\quad$ N

$M' \quad \xrightarrow{c} \quad N'$

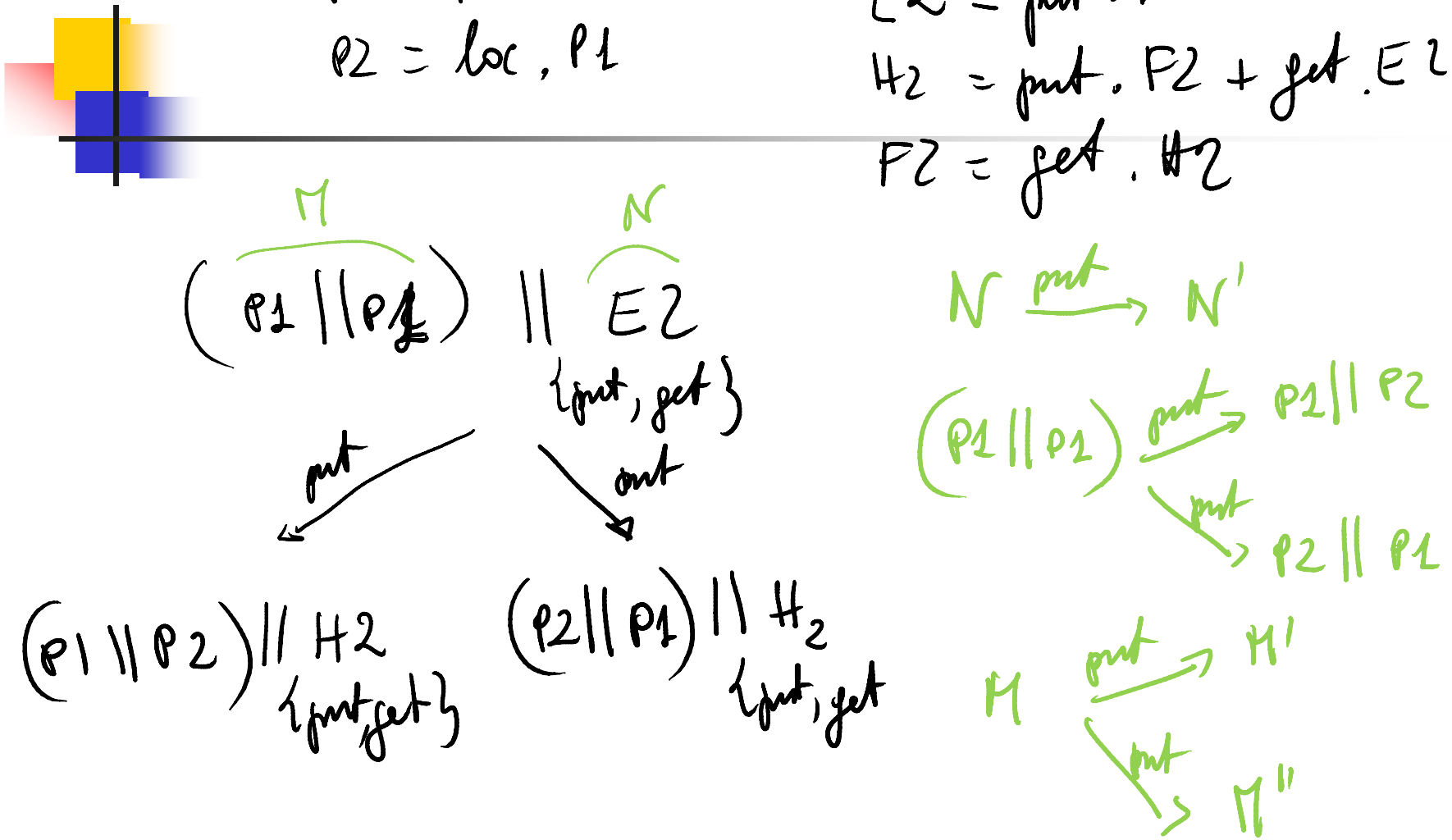
$(0 \|_c 0) \|_c 0$

$M \xrightarrow{c} M'$

$(c \|_c c) \xrightarrow{c} 0 \|_c 0$

$N : c.0 \xrightarrow{c} 0$

$N \xrightarrow{c} N'$



È importante che in $(P1||_{\{.\}}P1)||_{\{put,get\}}E2$ l'insieme di sincronizzazione del primo parallelo sia vuoto, sennò I due processi P1 si sincronizzerebbero fra loro e con il buffer

CSP-like process algebra

Consequences of the new proof rule for the parallel composition :

$(E \| F) \|_C G$

$(c.0 \| c.0) \|_{\{c\}} c.0$

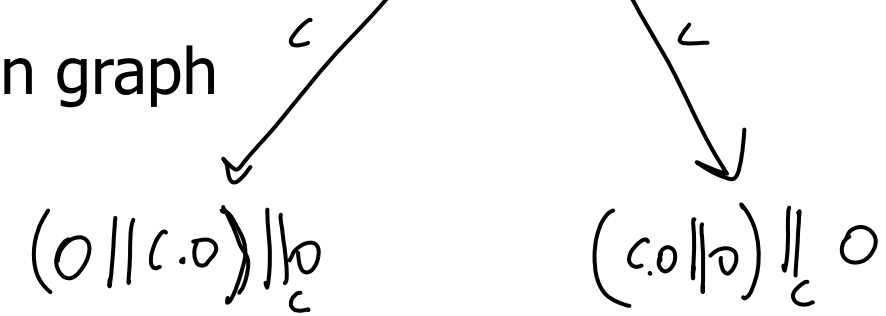
$E \| F \xrightarrow{c} 0 \| c.0$

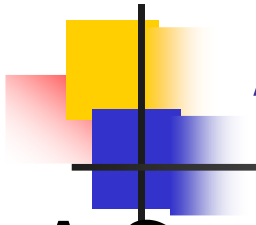
$E \| F \xrightarrow{c} c.0 \| 0$

$G \xrightarrow{c} 0$

$E \| F \not\xrightarrow{} 0 \| 0$

Derivation graph





Algebra

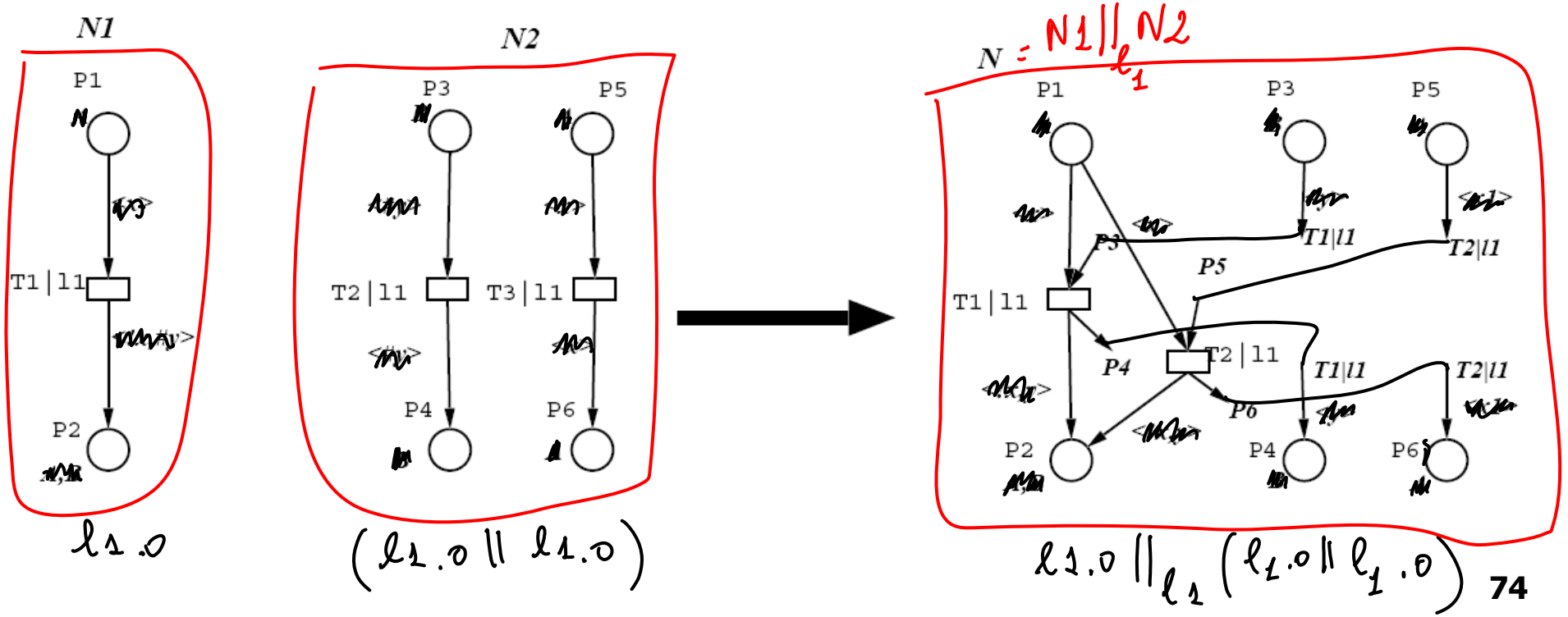
A C program that takes a set of LGSPN and LSWN nets and it superposes them on transition/place (subset of) labels

It allows non-injective multilabelling (cross product of trans. and places and n-process synchronization)

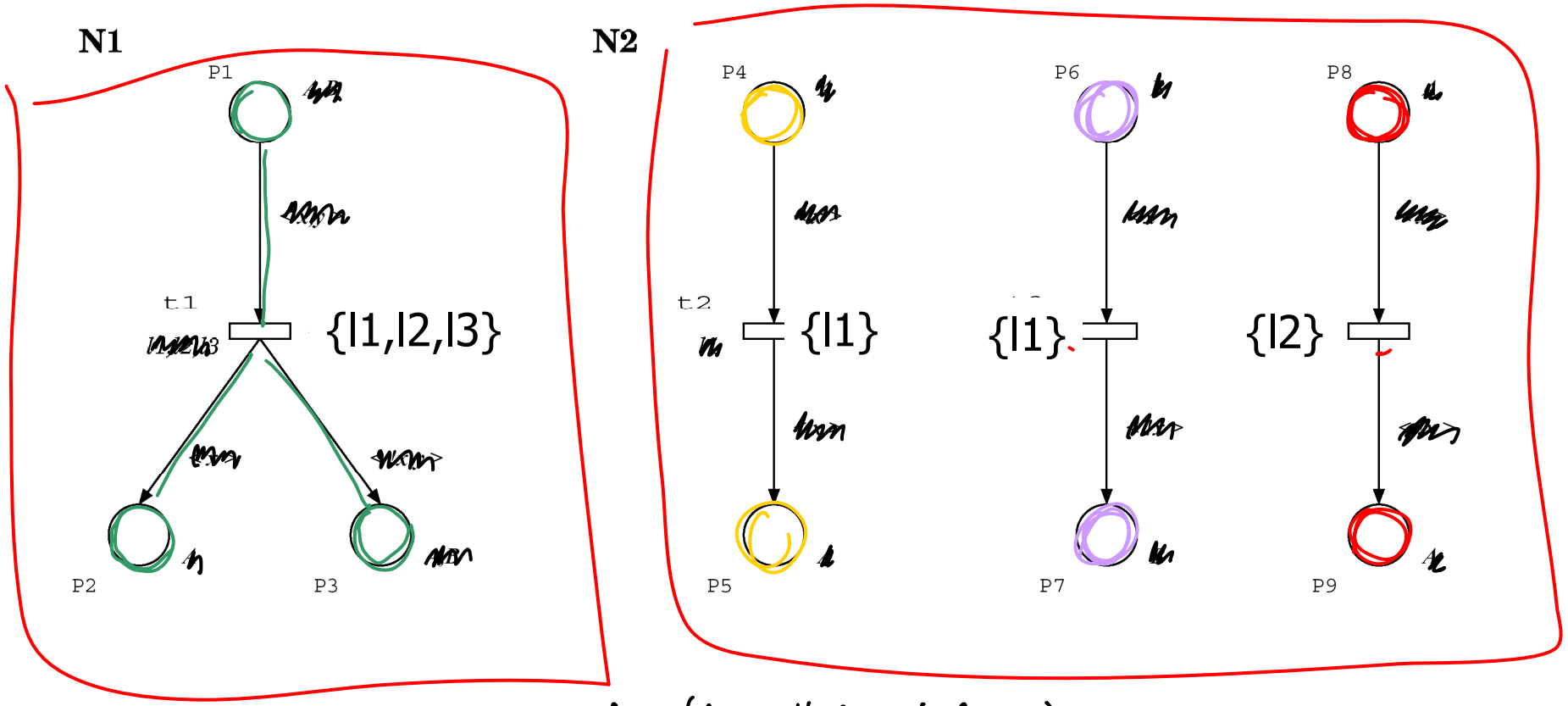
Algebra

A C program that takes a set of LGSPN and LSWN nets and it superposes them on transition/place (subset of) labels

$$N1 \;||_{\{l1\}}\; N2$$

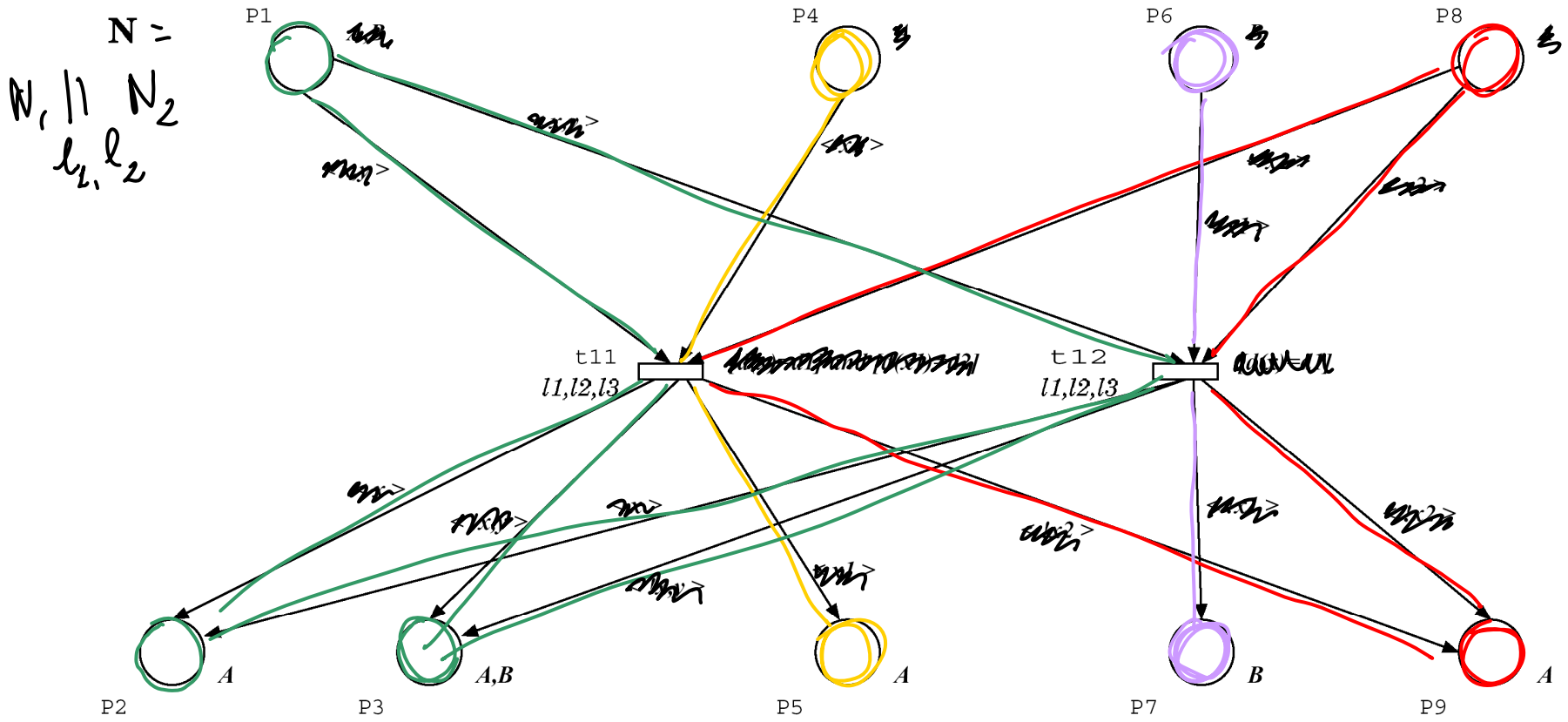


Algebra: transition superposition

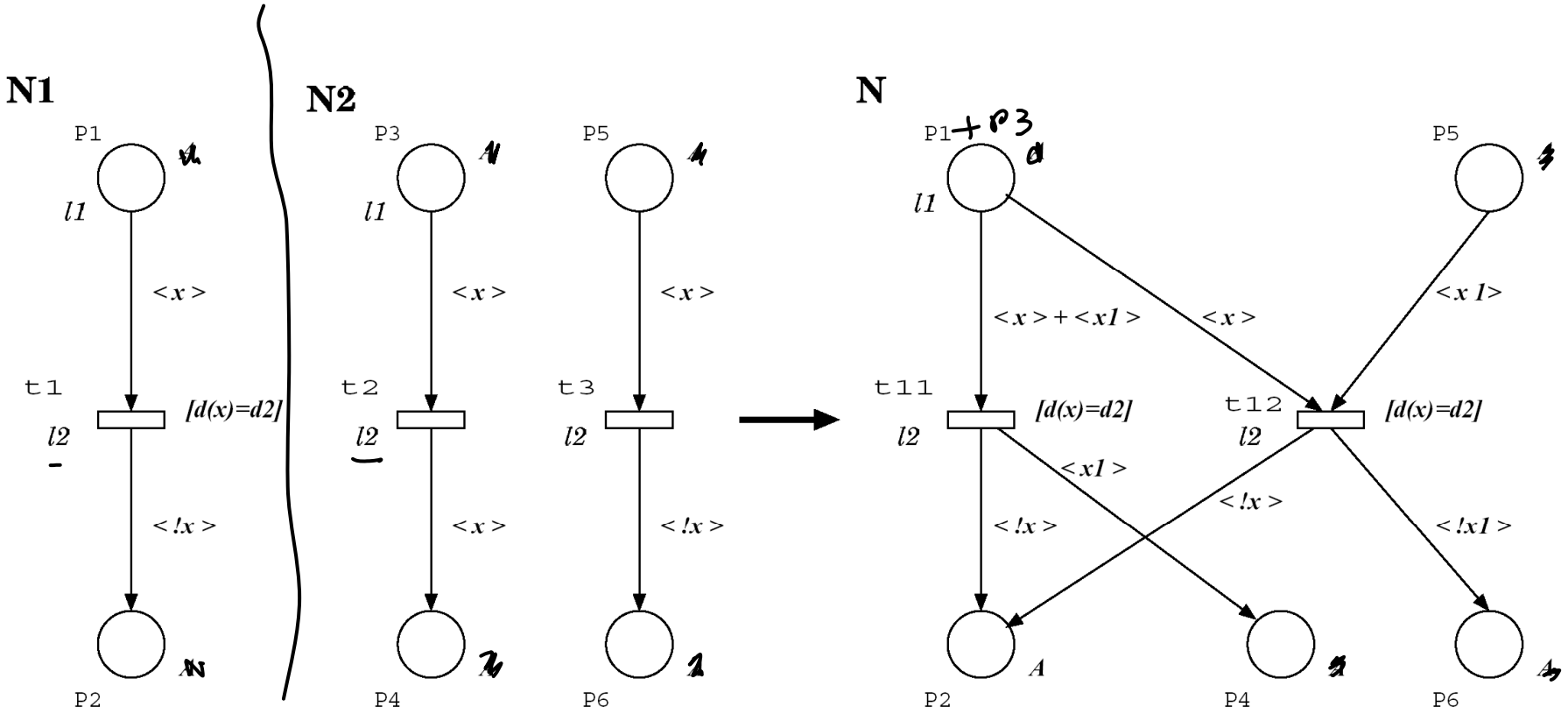


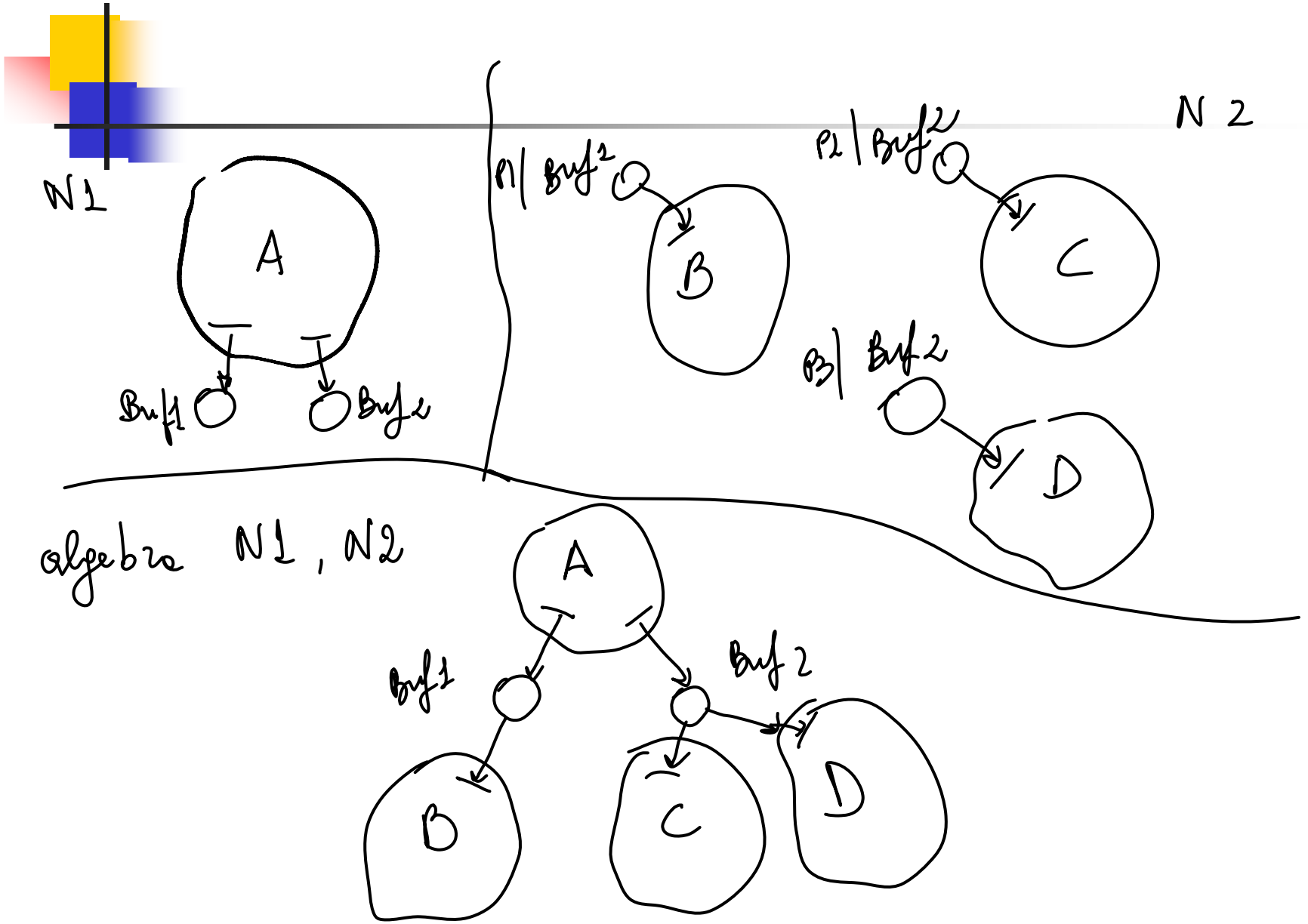
$N_2 = (l_{1.0} \,||\, l_{1.0} \,||\, l_{2.0})$

Algebra: transition superposition



Algebra: place and transition superposition

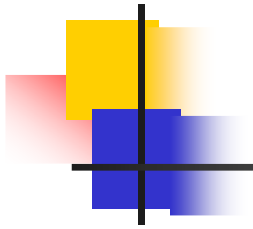



N 2
N1
A
P1 | Buf2
B
P2 | Buf2
C
B1 | Buf2
D
Buf1
Buf2
algebra N1, N2
A
Buf1
Buf 2
B
C
D

Algebra script

```
>> algebra net1 net2 op labels net
                        [plac. Dx Dy]
```

- `net1 net2` input nets
- `op = t` or p for type of superposition
- `labels` to be considered by `op`
- `net` resulting net
- `[plac. Dx Dy]` placement of `net1 net2` in `net`



Algebra

Additional remove function to eliminate

- # in tags (used to avoid renaming of variables)
- labels (solvers crash if tags have a "|" char)